

Subtyping for Access Control

MyThS REVIEW MEETING

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Group Types for Mobility

Aim: *Resource Access Control*

- Detect and prevent unwanted access to resources.
- We focus on *static* approaches based on enforcing type disciplines.

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Groups: Sets of processes with common access rights. (Cardelli-Ghelli-Gordon)

Constraints like $k : \text{CanEnter}(n)$ are modelled as:

n belongs to group G

k may cross the border of any ambient of group G .

For instance, the system:

$$k[\text{in } n \mid l[\text{out } k]] \mid n[-]$$

is *well-typed* under assumptions of the form:

$$k : \text{amb}[K, \text{cross}(N)]$$
$$l : \text{amb}[L, \text{cross}(K)]$$
$$n : \text{amb}[N, \dots]$$

Indirect Border Crossing in MA

Trojan Horses: The system

$\text{Odysseus}[\text{in Horse.out Horse.Destroy}] \mid \text{Horse}[\text{in Troy}] \mid \text{Troy}[\text{Trojans}]$

is well-typed under assumptions:

$\text{Odysseus} : \text{amb}[\text{Achaean}, \text{cross}(\text{Toy})]$

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$\text{Troy} : \text{amb}[\text{City}, -]$

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$\text{Troy} : \text{amb}[\text{City}, -]$

However, the system may evolve to

$\text{Troy}[\text{Trojans} \mid \text{Horse}[-] \mid \text{Odysseus}[\text{Destroy}]]$

where Odysseus got inside **Troy's Walls** taking by surprise the *Trojans*.

Typed Boxed Ambients: Syntax

$\eta ::= n$ names
| \uparrow enclosing amb
| \star local

$P ::= \mathbf{0}$ nil process
| $P_1 | P_2$ composition
| $(\nu n:A)P$ restriction
| $!P$ replication
| $V[P]$ ambient
| $V.P$ prefixing
| $(x : W)^\eta.P$ input
| $\langle V \rangle^\eta.P$ output
| $(\nu G)P$ group creation

$V, U ::= n$ name
| $\mathbf{in} V$ may enter V
| $\mathbf{out} V$ may exit V
| $V_1.V_2$ path

Reduction Semantics

Mobility:

$$n[\mathbf{in} m.P|Q] \mid m[R] \rightarrow m[n[P \mid Q] \mid R] \quad (\text{In})$$

$$m[n[\mathbf{out} m.P \mid Q] \mid R] \rightarrow n[P \mid Q] \mid m[R] \quad (\text{Out})$$

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$$m[n[\mathbf{out} m.P|Q] \mid R] \rightarrow n[P|Q] \mid m[R] \quad (\text{Out})$$

Communication:

$$(\mathbf{x}).P \mid \langle \mathbf{V} \rangle.Q \rightarrow P\{\mathbf{V}/\mathbf{x}\} \mid Q \quad (\text{Comm Local})$$

$$(\mathbf{x})^n.P \mid n[\langle \mathbf{V} \rangle.Q \mid R] \rightarrow P\{\mathbf{V}/\mathbf{x}\} \mid n[Q \mid R] \quad (\text{Comm Input } n)$$

$$(\mathbf{x}).P \mid n[\langle \mathbf{V} \rangle^\uparrow.Q \mid R] \rightarrow P\{\mathbf{V}/\mathbf{x}\} \mid n[Q|R] \quad (\text{Comm Output } \uparrow)$$

$$\langle \mathbf{V} \rangle^n.P \mid n[(\mathbf{x}).Q \mid R] \rightarrow P \mid n[Q\{\mathbf{V}/\mathbf{x}\} \mid R] \quad (\text{Comm Output } n)$$

$$\langle \mathbf{V} \rangle.P \mid n[(\mathbf{x})^\uparrow.Q \mid R] \rightarrow P \mid n[Q\{\mathbf{V}/\mathbf{x}\} \mid R] \quad (\text{Comm Input } \uparrow)$$

Types

Groups: G, H, \dots

Sets of groups: $\mathcal{G}, \mathcal{D}, \mathcal{I}, \dots$ \mathcal{U} The universal set of groups

Ambients types:

$A ::= \text{amb}_{\chi}[G, M, C]$ amb of group G , good for actions $\chi \subseteq \{i, o, c, r, w\}$,
with mobility type M , and communication type C

Process types:

$\Pi ::= \text{proc}[G, M, C]$ process that can be enclosed in an ambient of group G ,
may drive to ambients whose groups are in M ,
and communicates as described by type C

Capability types:

$K ::= \text{cap}[G, M, F]$ capability that can appear in an ambient of group G ,
may drive it to ambients whose groups are in M ,
with exchange type F for local communication

Types (cont.)

Mobility types:

$M ::= \text{mob}[\mathcal{G}]$

mobility specs

Communication types:

$C ::= \text{com}[E, F]$

E for local and F for upward exchange

Exchange types:

$E, F ::= \text{rw}[I, O]$

read/write values (valid if $O \prec I$)

Message types:

$I, O ::= \perp \mid W_1 \times \dots \times W_k \mid \top$

bottom, tuple, top

Value types:

$W, Y ::= A$
 $\mid K$

ambient name

capability

Subtyping

(sAmb)

$$\frac{\chi_1 \subseteq \chi_0 \subseteq \{\mathbf{i}, \mathbf{o}, \mathbf{c}, \mathbf{r}, \mathbf{w}\}}{\text{amb}_{\chi_0}[\mathbf{G}, M, C] \prec \text{amb}_{\chi_1}[\mathbf{G}, M, C]}$$

(sProc)

$$\frac{M_0 \prec M_1; \quad C_0 \prec C_1}{\text{proc}[\mathbf{G}, M_0, C_0] \prec \text{proc}[\mathbf{G}, M_1, C_1]}$$

(sCap)

$$\frac{M_0 \prec M_1; \quad F_0 \prec F_1}{\text{cap}[\mathbf{G}, M_0, F_0] \prec \text{cap}[\mathbf{G}, M_1, F_1]}$$

(sMob)

$$\frac{\mathcal{G}_0 \subseteq \mathcal{G}_1}{\text{mob}[\mathcal{G}_0] \prec \text{mob}[\mathcal{G}_1]}$$

(sCom)

$$\frac{E_0 \prec E_1; \quad F_0 \prec F_1}{\text{com}[E_0, F_0] \prec \text{com}[E_1, F_1]}$$

(sExc)

$$\frac{I_1 \prec I_0; \quad O_0 \prec O_1}{\text{rw}[I_0, O_0] \prec \text{rw}[I_1, O_1]}$$

(sMsg)

$$\frac{-}{\perp \prec W_1 \times \dots \times W_k \prec \top}$$

(sTuple)

$$\frac{W_i \prec T_i; \quad i \in 1..k}{W_1 \times \dots \times W_k \prec T_1 \times \dots \times T_k}$$

Good Values

$$\frac{\text{(Val n)} \quad \Gamma, n : W, \Gamma' \vdash \diamond}{\Gamma, n : W, \Gamma' \vdash n : W}$$

$$\frac{\text{(Val pfx)} \quad \Gamma \vdash V_0 : K; \quad \Gamma \vdash V_1 : K}{\Gamma \vdash V_0.V_1 : K}$$

$$\frac{\text{(Val in)} \quad \Gamma \vdash V : \text{amb}_i[\mathbf{G}, M, \text{com}[E, F]]}{\Gamma \vdash \text{in } V : \text{cap}[\mathbf{H}, \text{mob}[\{\mathbf{G}\}], E]} \quad \mathbf{H} \in \Gamma$$

$$\frac{\text{(Val sub)} \quad \Gamma \vdash V : W; \quad W \prec W'}{\Gamma \vdash V : W'}$$

$$\frac{\text{(Val out)} \quad \Gamma \vdash V : \text{amb}_o[\mathbf{G}, M, \text{com}[E, F]]}{\Gamma \vdash \text{out } V : \text{cap}[\mathbf{H}, M, F]} \quad \mathbf{H} \in \Gamma$$

Good Processes – Mobility

(Pro pfx)

$$\frac{\Gamma \vdash V:\text{cap}[\mathbf{G}, M, F]; \quad \Gamma \vdash P:\text{proc}[\mathbf{G}, M, \text{com}[E, F]]}{\Gamma \vdash V.P:\text{proc}[\mathbf{G}, M, \text{com}[E, F]]}$$

(Pro amb)

$$\frac{\Gamma \vdash V:\text{amb}_c[\mathbf{H}, \text{mob}[\mathcal{S}], \text{com}[E, F]]; \quad \Gamma \vdash P:\text{proc}[\mathbf{H}, \text{mob}[\mathcal{S}], \text{com}[E, F]]}{\Gamma \vdash V[P]:\text{proc}[\mathbf{G}, \text{mob}[\emptyset], \text{com}[F, \text{zero}]]} \quad \mathbf{G} \in \mathcal{S}$$

(Pro res)

$$\frac{\Gamma, n : A \vdash P : \Pi}{\Gamma \vdash (\nu n : A)P : \Pi}$$

(Pro gres)

$$\frac{\Gamma, \mathbf{G} \vdash P : \Pi}{\Gamma \vdash (\nu \mathbf{G})P : \Pi} \quad \mathbf{G} \notin \text{fg}(\Pi)$$

(Pro 0)

$$\frac{\mathbf{G} \in \text{dom}(\Gamma)}{\Gamma \vdash \mathbf{0} : \text{proc}[\mathbf{G}, \text{mob}[\emptyset], \text{com}[\text{zero}, \text{zero}]]}$$

(Pro par)

$$\frac{\Gamma \vdash P : \Pi; \quad \Gamma \vdash Q : \Pi}{\Gamma \vdash P \mid Q : \Pi}$$

(Pro rep)

$$\frac{\Gamma \vdash P : \Pi}{\Gamma \vdash !P : \Pi}$$

(Pro sub)

$$\frac{\Gamma \vdash P : \Pi \quad \Pi \prec \Pi'}{\Gamma \vdash P : \Pi'}$$

Good Processes – Communication

(inp \star)

$$\frac{\Gamma, \mathbf{x} : \mathbf{W} \vdash P : \text{proc}[\mathbf{G}, M, \text{com}[\text{rw}[I, O], F]]}{\Gamma \vdash (\mathbf{x} : \mathbf{W}).P : \text{proc}[\mathbf{G}, M, \text{com}[\text{rw}[I, O], F]]} I \prec \mathbf{W}$$

(out \star)

$$\frac{\Gamma \vdash \mathbf{V} : \mathbf{W}; \quad \Gamma \vdash P : \text{proc}[\mathbf{G}, M, \text{com}[\text{rw}[I, \mathbf{W}], F]]}{\Gamma \vdash \langle \mathbf{V} \rangle.P : \text{proc}[\mathbf{G}, M, \text{com}[\text{rw}[I, \mathbf{W}], F]]}$$

(inp \uparrow)

$$\frac{\Gamma, \mathbf{x} : \mathbf{W} \vdash P : \text{proc}[\mathbf{G}, M, \text{com}[E, \text{rw}[I, O]]]}{\Gamma \vdash (\mathbf{x} : \mathbf{W}_k)^\uparrow.P : \text{proc}[\mathbf{G}, M, \text{com}[E, \text{rw}[I, O]]]} I \prec \mathbf{W}$$

(output \uparrow)

$$\frac{\Gamma \vdash \mathbf{V} : \mathbf{W}; \quad \Gamma \vdash P : \text{proc}[\mathbf{G}, M, \text{com}[E, \text{rw}[I, \mathbf{W}]]]}{\Gamma \vdash \langle \mathbf{V} \rangle^\uparrow.P : \text{proc}[\mathbf{G}, M, \text{com}[E, \text{rw}[I, \mathbf{W}]]]}$$

...and so on analogously

Properties

Communication properties:

➤ If $\Gamma \vdash (x : W).P \mid \langle V \rangle.Q : \Pi$ then $\Gamma \vdash V:Y$ with $Y \prec W$.

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Mobility properties:

- If $\Gamma \vdash n[\text{in } m.P \mid Q] \mid m[R] : \Pi$, then

$$\Gamma \vdash m : \text{amb}_{\chi_0}[\mathbf{M}, -, -] \quad \text{and} \quad \Gamma \vdash n : \text{amb}_{\chi_1}[-, \text{mob}[\mathcal{S}], -]$$

with $\mathbf{M} \in \mathcal{S}$, $i, c \in \chi_0$ and $c \in \chi_1$.

- If $\Gamma \vdash m[n[\text{out } m.P \mid Q] \mid R] : \Pi$, then

$$\Gamma \vdash m : \text{amb}_{\chi_0}[\mathbf{M}, \text{mob}[\mathcal{S}_m], -] \quad \text{and} \quad \Gamma \vdash n : \text{amb}_{\chi_1}[\mathbf{N}, \text{mob}[\mathcal{S}_n], -]$$

with $o, c \in \chi_0$, $c \in \chi_1$, $\mathbf{M} \in \mathcal{S}_n$ and $\mathcal{S}_m \subseteq \mathcal{S}_n$.

Properties

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with $\mathbf{M} \in \mathcal{S}$, $i, c \in \chi_0$ and $c \in \chi_1$.

- If $\Gamma \vdash m[n[\text{out } m.P \mid Q] \mid R] : \Pi$, then

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with $o, c \in \chi_0$, $c \in \chi_1$, $\mathbf{M} \in \mathcal{S}_n$ and $\mathcal{S}_m \subseteq \mathcal{S}_n$.

Subject reduction:

- If $\Gamma \vdash P : \Pi$ and $P \equiv Q$ or $P \rightarrow Q$, then there exist groups G_0, \dots, G_k such that $G_0, \dots, G_k, \Gamma \vdash Q : \Pi$.

Detecting Odysseus' intentions

Now, in order to assign a type to

`Odysseus[in Horse.out Horse.Destroy] | Horse[in Troy] | Troy[Trojans]`

we need assumptions of the form:

`Odysseus : ambc[Achaean, mob[{Ground, Toy, City}], -]`

`Horse : ambioc[Toy, mob[{Ground, City}], -]`

`Troy : ambioc[City, -, -]`

representing that `Odysseus` is an `Achaean` intended to move into a `City`!

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$\text{Odysseus}[\text{in Horse.out Horse.Destroy}] \mid \text{Horse}[\text{in Troy}] \mid \text{Troy}[\text{Trojans}]$

we need assumptions of the form:

$\text{Odysseus} : \text{amb}_c[\text{Achaean}, \text{mob}[\{\text{Ground}, \text{Toy}, \text{City}\}], -]$

$\text{Horse} : \text{amb}_{ioc}[\text{Toy}, \text{mob}[\{\text{Ground}, \text{City}\}], -]$

$\text{Troy} : \text{amb}_{ioc}[\text{City}, -, -]$

representing that *Odysseus* is an *Achaean* intended to move into a *City*!
On the other hand, under assumptions of the form

$\text{Odysseus} : \text{amb}_c[\text{Achaean}, \text{mob}[\{\text{Ground}, \text{Toy}\}], -]$

the *Trojans* should not fear any attack from *Odysseus*.

But what if *Odysseus* is *lying* about his intentions (i.e. type)?



BSA: Adding co-capabilities

Reduction Semantics:

$$n[\mathbf{in} m.P \mid Q] \mid m[\overline{\mathbf{in}} \alpha.R \mid S] \rightarrow m[n[P \mid Q] \mid R \mid S] \quad \text{for } \alpha \in \{\star, n\}$$

$$m[n[\mathbf{out} m.P \mid Q] \mid R] \mid \overline{\mathbf{out}} \alpha.S \rightarrow n[P \mid Q] \mid m[R] \mid S \quad \text{for } \alpha \in \{\star, n\}$$

Mobility Types: (extended: \mathcal{C} tells which processes are allowed.)

$$M ::= \text{mob}[\mathcal{I}, \mathcal{C}]$$

Subtyping Relation: (extended)

$$\frac{\text{(sMob)} \quad \mathcal{G}_0 \subseteq \mathcal{G}_1, \quad \mathcal{C}_0 \subseteq \mathcal{C}_1}{\text{mob}[\mathcal{G}_0, \mathcal{C}_0] \prec \text{mob}[\mathcal{G}_1, \mathcal{C}_1]}$$

Good Values in BSA

(Val in)

$$\frac{\Gamma \vdash V:\text{amb}_i[\mathbf{G}, M, \text{com}[E, F]]}{\Gamma \vdash \text{in } V:\text{cap}[\mathbf{H}, \text{mob}[\{G\}, \emptyset], E]}$$

(Val out)

$$\frac{\Gamma \vdash V:\text{amb}_o[\mathbf{G}, \text{mob}[\mathcal{S}, \mathcal{C}], \text{com}[E, F]]}{\Gamma \vdash \text{out } V:\text{cap}[\mathbf{H}, \text{mob}[\mathcal{S}, \emptyset], F]}$$

(Val coin)

$$\frac{\Gamma \vdash V:\text{amb}_i[\mathbf{G}, M, \text{com}[E, F]]}{\Gamma \vdash \overline{\text{in}} V:\text{cap}[\mathbf{H}, \text{mob}[\emptyset, \{\mathbf{G}\}], F]}$$

(Val coout)

$$\frac{\Gamma \vdash V:\text{amb}_o[\mathbf{G}, M, \text{com}[E, F]]}{\Gamma \vdash \overline{\text{out}} V:\text{cap}[\mathbf{H}, \text{mob}[\emptyset, \{\mathbf{G}\}], F]}$$

(Val coin \star)

$$\frac{\mathbf{G} \in \text{dom}(\Gamma)}{\Gamma \vdash \overline{\text{in}} \star:\text{cap}[\mathbf{G}, \text{mob}[\emptyset, \mathcal{U}], \text{zero}]}$$

(Val coout \star)

$$\frac{\mathbf{G} \in \text{dom}(\Gamma)}{\Gamma \vdash \overline{\text{out}} \star:\text{cap}[\mathbf{G}, \text{mob}[\emptyset, \mathcal{U}], \text{zero}]}$$

In (Val in), (Val out), (Val coin), (Val coin), assume $\mathbf{H} \in \Gamma$

Good Processes – Mobility in BSA

(Pro 0)

$$\frac{\mathbf{G} \in \text{dom}(\Gamma)}{\Gamma \vdash \mathbf{0}:\text{proc}[\mathbf{G}, \text{mob}[\emptyset, \emptyset], \text{com}[\text{zero}, \text{zero}]]}$$

(Pro pfx)

$$\frac{\Gamma \vdash V:\text{cap}[\mathbf{G}, M, F]; \quad \Gamma \vdash P:\text{proc}[\mathbf{G}, M, \text{com}[E, F]]}{\Gamma \vdash V.P:\text{proc}[\mathbf{G}, M, \text{com}[E, F]]}$$

(Pro amb)

$$\frac{\begin{array}{l} \Gamma \vdash V:\text{amb}_c[\mathbf{H}, \text{mob}[\mathcal{S}, \mathcal{C}], \text{com}[E, F]] \\ \Gamma \vdash P:\text{proc}[\mathbf{H}, \text{mob}[\mathcal{S}, \mathcal{C}], \text{com}[E, F]] \end{array}}{\Gamma \vdash V[P]:\text{proc}[\mathbf{G}, \text{mob}[\emptyset, \{\mathbf{H}\}], \text{com}[F, \text{zero}]]} \quad \mathbf{G} \in \mathcal{S}$$

Control Properties in BSA

Access Control Theorem:

Whenever

$$\Gamma \vdash m[\overline{\text{in}}\alpha.P \mid Q] : \Pi \quad \text{or} \quad \Gamma \vdash m[\overline{\text{out}}\alpha.P \mid Q] : \Pi,$$

with $\alpha \in \{\star, n\}$, then

- $\Gamma \vdash m : \text{amb}_{\chi_0}[-, \text{mob}[-, \mathcal{C}], -]$, and
 - either $\alpha = \star$ and $\mathcal{C} = \mathcal{U}$,
 - or $\alpha = n$ with $\Gamma \vdash n : \text{amb}_{\chi_1}[\mathbf{N}, -, -]$ and $\mathbf{N} \in \mathcal{C}$.

Using co-capabilities to defend Troy

Our running example in BSA:

The Trojan War \triangleq `Odysseus[in Horse.out Horse.Destroy]`

| `Horse[$\overline{\text{in}}$ * .in Troy]`

| `Troy[$\overline{\text{in}}$ Horse.Trojans | $\overline{\text{out}}$ Odysseus.Sinon]`

which can be *well-typed* only if

$\Gamma \vdash \text{Troy} : \text{amb}_{\text{ioc}}[\text{City}, \text{mob}[-, \{\text{Toy}, \text{Achaean}\}], -]$

That is if Troy (in suicidal mood) allowed *Achaean* in.

Using co-capabilities to defend Troy (ctd)

Consider now the system:

$$\begin{aligned} \text{The Trojan Trap} &\triangleq \text{Odysseus}[\text{in Horse.out Horse.Destroy}] \\ &| \text{Horse}[\overline{\text{in}} \star .\text{in Troy}] \\ &| \text{Troy}[\overline{\text{in}} \text{Horse.Trojans}] \end{aligned}$$

This situation would be perfectly safe for Troy (but dangerous for Odysseus!) provided we can type it under the assumptions of the form

$$\text{Odysseus} : \text{amb}_c[\text{Achaean}, -, -]$$
$$\text{Horse} : \text{amb}_{ioc}[\text{Toy}, -, \text{com}[E, \mathbf{0}]]$$
$$\text{Troy} : \text{amb}_{ioc}[\text{City}, \text{mob}[\emptyset, \mathcal{C}], -]$$

with $\text{Achaean} \notin \mathcal{C}$.

