

School of Electronics and Computer Science

Data provenance in a distributed calculus

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work in progress with A. Francalanza, J. Rathke, and I. Souilah

Motivation

- ♦ (Meta)data is almost entirely neglected in the process calculi literature
- ♣ Track data provenance both for its important applications and as an challenging exercise in modelling (meta)data. We aim at simplicity:
 - data annotations representing provenance
 - * structure, interpretation and management of provenance information
 - provenance tracking
- ♦ Provenance-based security (aspects: trust + data confidentiality and privacy)
 - Example: photography competition
- ♦ The overall ambition is to underpin practical development, like trust-policy languages and protocols, and provenance-middleware

Model features

- ♦ Two the central features of the basic model:
 - 1. values are annotated with their provenance
 - 2. provenance is kept up-to-date as computation proceeds
- ♦ Focus on one particular kind of provenance information:
 - The principals that influenced a value, and how they did it



Annotated data

Annotated value

 $\begin{bmatrix} v : K \end{bmatrix}$

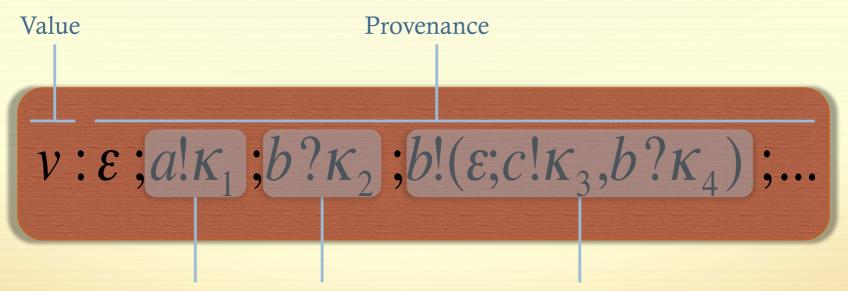
Value

Actual data

Provenance

Meta information describing the origin of the value

Structure and interpretation of provenance



"Operations" that were performed on the value. They record the principals that "influenced" the value and how.

Structure and interpretation of provenance

 ε (empty provenance) denotes value ν originated here



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 $v: \varepsilon; a! \kappa_1$

It was sent by a on a channel with provenance κ_1

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 $v: \varepsilon; a! \kappa_1; b? \kappa_2$

It was sent by a on a channel with provenance κ_1

Was then received by b on a channel with provenance κ_2

Structure and interpretation of provenance

 ε (empty provenance) denotes value ν originated here

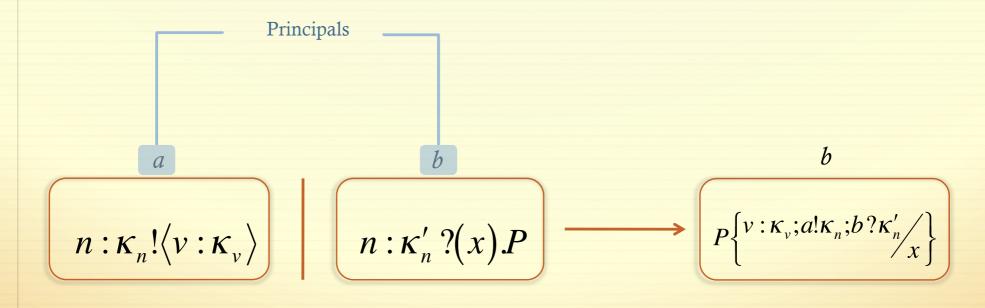
And then sent by b on a channel that b received from c...

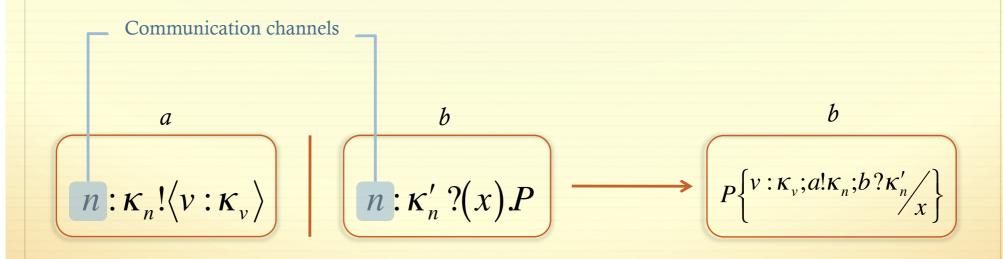
 $v: \varepsilon; a! \kappa_1; b? \kappa_2; b! (\varepsilon; c! \kappa_3, b? \kappa_4); \dots$

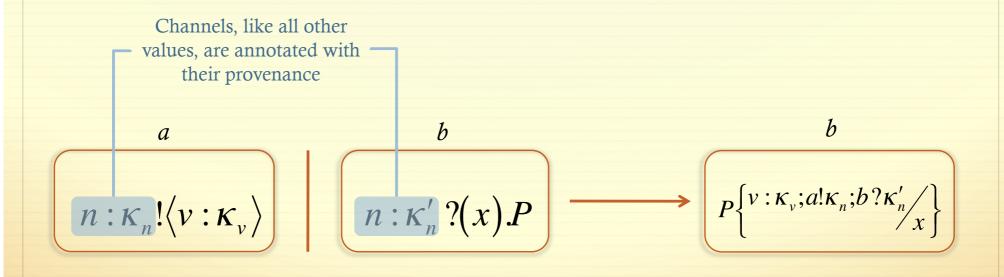
It was sent by a on a channel with provenance κ_1

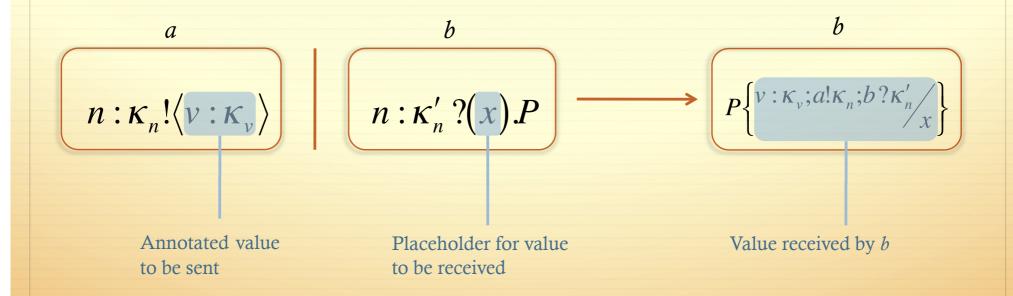
Was then received by b on a channel with provenance κ_2

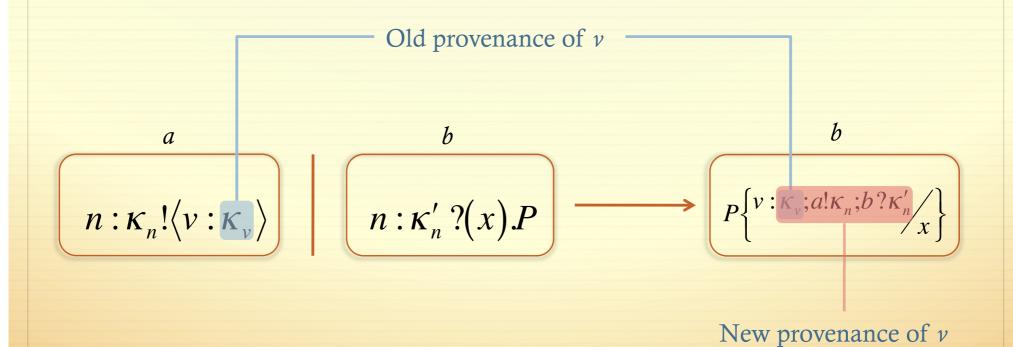
Provenance tracking











Provenance tracking

Old provenance of ν

a

$$n: \kappa_n! \langle v: \kappa_v \rangle$$

h

$$n:\kappa'_n?(x)P$$

ľ

$$P\left\{v: \kappa_{v}; a! \kappa_{n}; b? \kappa_{n}'\right\}$$

Sender and provenance of channel used

Receiver and provenance of channel used

Confidentiality in provenance systems

- ♦ Data may be public, yet its provenance confidential, or vice versa
- Principals who may access data are not necessarily the same as those who may access its provenance
- ♣ In general, fine grained access control over provenance "histories" is needed as different parts of it have different sensitivity

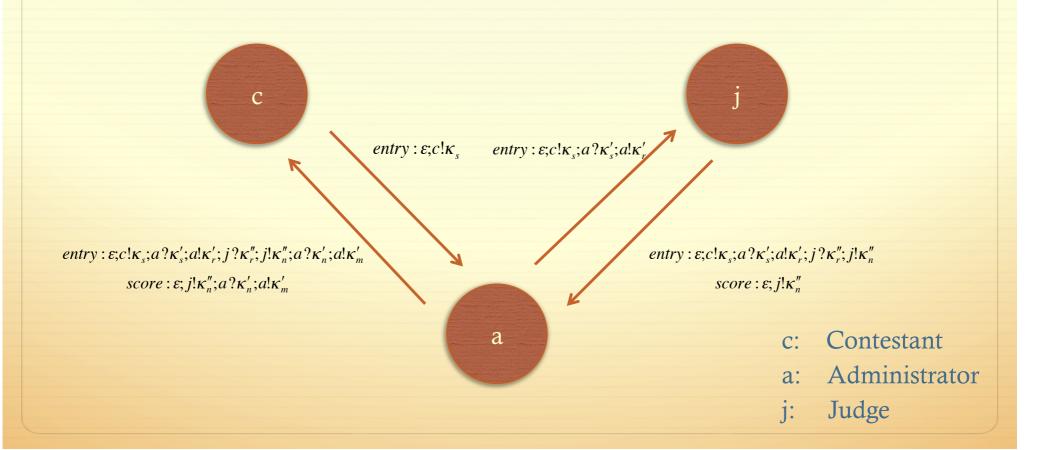
Security requirements of data



Security requirements of its provenance

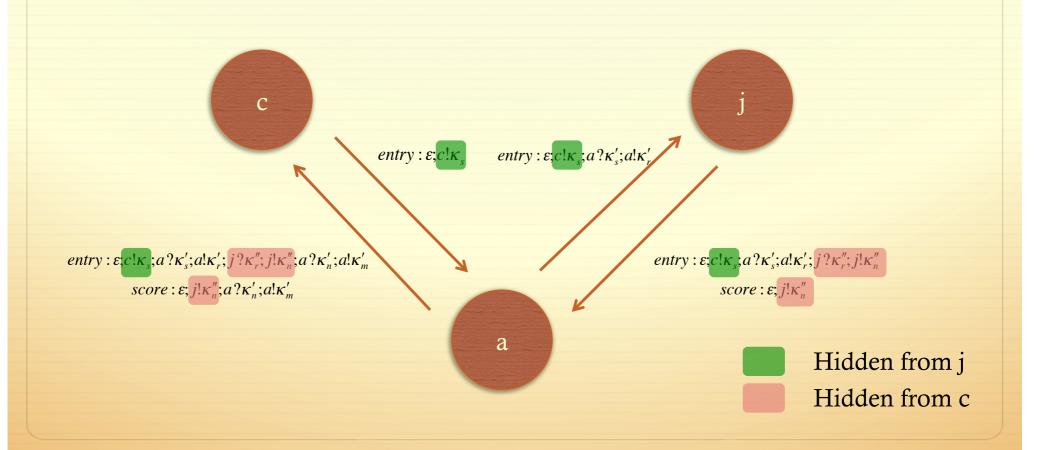
Hiding provenance trees

Example: photography competition



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Confidentiality in provenance systems

a promising approach

- ♦ One value, multiple views
 - → Different principals have different views of the same provenance list based on their privileges

entry: ε ; $c!\kappa_s$; $a?\kappa'_s$; $a!\kappa'_r$; $j?\kappa''_r$; $j!\kappa''_n$; $a?\kappa'_n$; $a!\kappa'_m$

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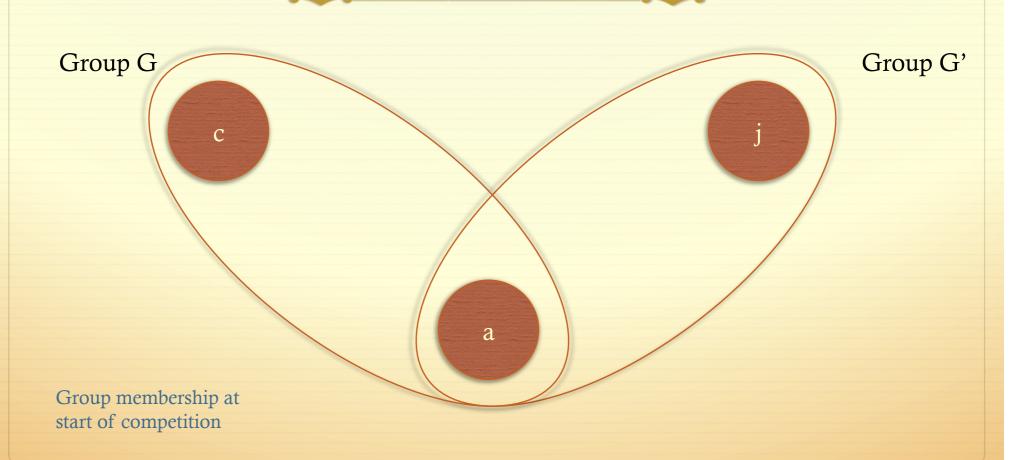
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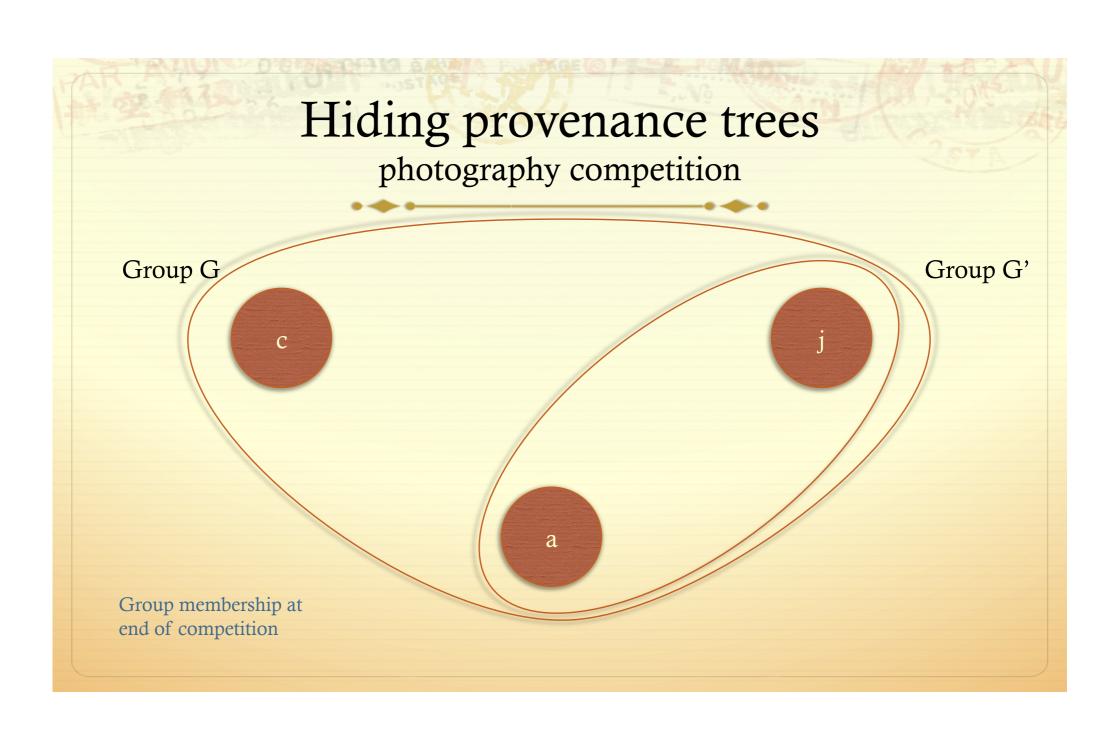
Confidentiality in provenance systems a promising approach

- ♦ To achieve this, we use groups
- ♦ Different principals belong to different groups
- ♦ Group membership determines what parts of a provenance list a principal has access to
- ♦ Principals
 - \leftarrow Can create new groups: new G
 - \leftarrow Can add other principals to their groups: add(a,G)
 - Can **restrict** access to particular parts of a provenance tree to a particular group: $hide(v : \kappa, G)$

Hiding provenance trees

photography competition





Current work

- ♦ Correctness of provenance tracking: the provenance information determines the history of each piece of data accurately "enough"
 - * Express this as a form of testing (on traces):

$$\forall S(\forall t \in [S](t \to^* v : \kappa \implies \forall s \in [v : \kappa](s|t \to^* \checkmark)))$$

- ♦ Using provenance:
 - Provenance queries vs pattern restricted input
 - Trust in quality of data based on trust in principals and provenance of data
- ♦ Policies and types