

Iterative Detection and Decoding Using Approximate Bayesian Theorem Based PDA Method Over MIMO Nakagami- m Fading Channels

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Abstract—In this paper, the design of iterative detection and decoding (IDD) schemes relying on a low-complexity probabilistic data association (PDA) aided method is conceived for turbo-coded multiple-input multiple-output (MIMO) systems communicating over Nakagami- m fading channels. The known PDA based MIMO detectors typically operate purely in the probability-domain. We show that the classic relationship where the extrinsic LLRs are given by subtracting the *a priori* LLRs from the *a posteriori* LLRs does not hold for the existing PDA based MIMO detectors. Therefore, the PDA method is not readily applicable to the IDD receiver. To overcome this predicament, we propose an approximate Bayesian theorem based log-domain PDA (AB-Log-PDA) detector, as well as a novel simple approach of calculating the bit-wise extrinsic LLRs for the AB-Log-PDA, which makes the AB-Log-PDA well-suited for employment in IDD receivers. It is shown that the proposed AB-Log-PDA based IDD scheme is capable of achieving a comparable performance to that of the optimal maximum *a posteriori* (MAP) detector based IDD receiver, while imposing a much lower computational complexity in the scenarios considered.

Index Terms—Extrinsic information transfer (EXIT) chart, iterative detection and decoding, multiple-input multiple-output (MIMO), probabilistic data association (PDA), Nakagami- m fading channel

I. INTRODUCTION

Iterative detection and decoding (IDD) is capable of achieving a near-optimum performance at a significantly lower complexity than the optimal joint detector/decoder [1]. Even so, the computational complexity imposed by the IDD might remain the limiting factor in practical applications.

The Probabilistic Data Association (PDA) method is a reduced-complexity design alternative of maximum *a posteriori* (MAP) decoders/detectors/equalizers [2]–[7]. As an efficient interference-modelling process, the key feature of PDA is the repeated conversion of a multimodal Gaussian mixture probability to a single multivariate Gaussian distribution. Therefore, the accuracy of the Gaussian approximation dominates the attainable performance. In uncoded MIMO systems using quadrature amplitude modulation (QAM), the quality of the Gaussian approximation in PDA may be improved by transforming the symbol-based model into a bit-based model,

which *in effect* increases the length of the effective transmitted signal vector by the number of bits per symbol, and reduces the effective constellation to a binary constellation [6]. Regarding improving the quality of the Gaussian approximation in FEC coded MIMO systems, we benefit from having an increased degree of freedom to exploit. For example, the soft information gleaned from the output of the FEC decoder via feedback is deemed to be a more reliable information source than the raw received signal at the fading channel's output, and therefore it may reduce the bias in modelling the inter-antenna interference (IAI) components.

Against this background, in this paper we aim for designing a low-complexity IDD scheme relying on the PDA method for FEC coded MIMO systems using M -QAM for transmission over Nakagami- m fading channels. However, there are further particular challenges that render the IDD design using PDA less straightforward than it seems to be. Firstly, to the best of our knowledge, all the existing PDA detectors conceived for uncoded systems [2]–[7] operate purely in the probability-domain, which results in a poor numerical stability and low accuracy in IDD scenarios and hence leads to a degraded performance. By contrast, in this paper we propose a new approximate Bayesian theorem based log-domain PDA (AB-Log-PDA) MIMO detector, which is better suited for the IDD scheme. Furthermore, for the existing family of PDA methods [2]–[7] as well as the proposed AB-Log-PDA, it is unclear how to produce the “correct” extrinsic log-likelihood ratios (LLRs), required by the concatenated outer FEC decoder. A natural way of generating the bit-wise extrinsic LLRs is to subtract the bit-wise *a priori* LLRs from the bit-wise *a posteriori* LLRs generated from the estimated symbol-wise APPs of the PDA. However, our study demonstrate that the symbol-wise APPs produced by the approximate Bayesian theorem based PDA detector are *not* the “true” APPs, but rather some sort of “nominal” APPs, for which the classic relationship where the *a posteriori* LLRs are given by the sum of the *a priori* LLRs and the extrinsic LLRs does not hold. Nonetheless, the symbol-wise probabilities output by the PDA based methods have been referred to as APPs [2]–[7], and the distinctions between these “nominal” APPs and the true APPs have never been reported before. Therefore, a novel approach of producing the bit-wise extrinsic LLRs for the proposed AB-Log-PDA is presented, which makes it possible to design a

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TABLE I
PROBABILITIES COMPUTED IN ONE ITERATION BETWEEN THE
AB-LOG-PDA DETECTOR AND THE FEC DECODER

	1	2	...	m	...	M
$\mathbf{P}^{(z)}(1)$	$P_1^{(z)}(s_1 \mathbf{y})$	$P_2^{(z)}(s_1 \mathbf{y})$...	$P_m^{(z)}(s_1 \mathbf{y})$...	$P_M^{(z)}(s_1 \mathbf{y})$
$\mathbf{P}^{(z)}(2)$	$P_1^{(z)}(s_2 \mathbf{y})$	$P_2^{(z)}(s_2 \mathbf{y})$...	$P_m^{(z)}(s_2 \mathbf{y})$...	$P_M^{(z)}(s_2 \mathbf{y})$
\vdots	\vdots	\vdots	...	\vdots	...	\vdots
$\mathbf{P}^{(z)}(i)$	$P_1^{(z)}(s_i \mathbf{y})$	$P_2^{(z)}(s_i \mathbf{y})$...	$P_m^{(z)}(s_i \mathbf{y})$...	$P_M^{(z)}(s_i \mathbf{y})$
\vdots	\vdots	\vdots	...	\vdots	...	\vdots
$\mathbf{P}^{(z)}(N_t)$	$P_1^{(z)}(s_{N_t} \mathbf{y})$	$P_2^{(z)}(s_{N_t} \mathbf{y})$...	$P_m^{(z)}(s_{N_t} \mathbf{y})$...	$P_M^{(z)}(s_{N_t} \mathbf{y})$

specify the mean as

$$\boldsymbol{\mu}_i \triangleq \mathbb{E}(\mathbf{v}_i) = \sum_{k \neq i} \mathbb{E}(s_k) \mathbf{h}_k, \quad (4)$$

the covariance as

$$\mathbf{\Upsilon}_i \triangleq \mathbf{C}(\mathbf{v}_i) = \sum_{k \neq i} \mathbf{C}(s_k) \mathbf{h}_k \mathbf{h}_k^H + 2\sigma^2 \mathbf{I}_{N_r}, \quad (5)$$

and the pseudo-covariance as

$$\underline{\mathbf{\Upsilon}}_i \triangleq \mathbf{C}_p(\mathbf{v}_i) = \sum_{k \neq i} \mathbf{C}_p(s_k) \mathbf{h}_k \mathbf{h}_k^T. \quad (6)$$

For each symbol s_i , $i = 1, \dots, N_t$, we define an M -element symbol probability vector $\mathbf{P}^{(z)}(i)$, whose m -th element $P_m^{(z)}(s_i|\mathbf{y}) \triangleq P^{(z)}(s_i = a_m|\mathbf{y})$ is the estimate of the APP that we have $s_i = a_m$ at the z -th iteration between the ‘‘inner’’ AB-Log-PDA and the ‘‘outer’’ soft FEC decoder, where z is a nonnegative integer, and $m = 1, \dots, M$. For clarity, these probabilities are explicitly presented in form of the probability-matrix of Table I. Then we have

$$\mathbb{E}(s_k) = \sum_{m=1}^M a_m P^{(z)}(s_k = a_m|\mathbf{y}), \quad (7)$$

$$\mathbf{C}(s_k) = \sum_{m=1}^M (a_m - \mathbb{E}(s_k))(a_m - \mathbb{E}(s_k))^* P^{(z)}(s_k = a_m|\mathbf{y}), \quad (8)$$

and

$$\mathbf{C}_p(s_k) = \sum_{m=1}^M (a_m - \mathbb{E}(s_k))^2 P^{(z)}(s_k = a_m|\mathbf{y}), \quad (9)$$

where the pseudo-covariance of a complex random vector \mathbf{x} is defined as [9]

$$\mathbf{C}_p(\mathbf{x}) \triangleq \mathbb{E} \left[(\mathbf{x} - \mathbb{E}(\mathbf{x})) (\mathbf{x} - \mathbb{E}(\mathbf{x}))^T \right]. \quad (10)$$

Note that Eq. (4) - Eq. (9) effectively use all $\{\mathbf{P}^{(z)}(k)\}_{k \neq i}$ associated with the interfering signal $\{s_k\}_{k \neq i}$ to model \mathbf{v}_i . Since we do not have any outer *a priori* knowledge about the distribution of $s_i|\mathbf{y}$ at the beginning, an all-zero LLR vector will be provided as the input to the AB-Log-PDA. This all-zero LLR vector is equivalent to initializing $P^{(z)}(s_i = a_m|\mathbf{y})$ with a uniform distribution, i.e.

$$P^{(z)}(s_i = a_m|\mathbf{y}) = P^{(z)}(s_i = a_m) = \frac{1}{M}, \quad (11)$$

where $z = 0, \forall i = 1, \dots, N_t$ and $\forall m = 1, \dots, M$.

Based on the assumption that \mathbf{v}_i obeys the Gaussian distribution, $\mathbf{y}|s_i$ is also Gaussian distributed. Let us now define

$$\mathbf{w} \triangleq \mathbf{y} - s_i \mathbf{h}_i - \sum_{k \neq i} \mathbb{E}(s_k) \mathbf{h}_k \quad (12)$$

and

$$\beta_{m,i}^{(z+1)} \triangleq - \begin{bmatrix} \Re(\mathbf{w}) \\ \Im(\mathbf{w}) \end{bmatrix}^T \boldsymbol{\Lambda}_i^{-1} \begin{bmatrix} \Re(\mathbf{w}) \\ \Im(\mathbf{w}) \end{bmatrix}, \quad (13)$$

in which the *composite* covariance matrix $\boldsymbol{\Lambda}_i$ is defined as [5]

$$\boldsymbol{\Lambda}_i \triangleq \begin{bmatrix} \Re(\mathbf{\Upsilon}_i + \underline{\mathbf{\Upsilon}}_i) & -\Im(\mathbf{\Upsilon}_i - \underline{\mathbf{\Upsilon}}_i) \\ \Im(\mathbf{\Upsilon}_i + \underline{\mathbf{\Upsilon}}_i) & \Re(\mathbf{\Upsilon}_i - \underline{\mathbf{\Upsilon}}_i) \end{bmatrix}, \quad (14)$$

where $\Re(\cdot)$ and $\Im(\cdot)$ represent the real and imaginary part of a complex variable, respectively. Then the likelihood function of $\mathbf{y}|s_i = a_m$ at the $(z+1)$ -st iteration satisfies

$$p^{(z+1)}(\mathbf{y}|s_i = a_m) \propto \exp\left(\beta_{m,i}^{(z+1)}\right). \quad (15)$$

Upon invoking an approximate form of the Bayesian theorem, the *estimated APP* of symbol s_i at the $(z+1)$ -st iteration may be calculated in the probability-domain as

$$\begin{aligned} P^{(z+1)}(s_i = a_m|\mathbf{y}) &\approx \frac{p^{(z+1)}(\mathbf{y}|s_i = a_m)}{\sum_{m=1}^M p^{(z+1)}(\mathbf{y}|s_i = a_m)} \\ &= \frac{\exp\left(\beta_{m,i}^{(z+1)} - \gamma\right)}{\sum_{m=1}^M \exp\left(\beta_{m,i}^{(z+1)} - \gamma\right)}, \end{aligned} \quad (16)$$

where $\gamma \triangleq \max_{m=1, \dots, M} \beta_{m,i}^{(z+1)}$. In order to further improve the achievable numerical stability and accuracy, the log-domain form of (16) is formulated as

$$\begin{aligned} \psi_{m,i}^{(z+1)} &\triangleq \ln\left(P^{(z+1)}(s_i = a_m|\mathbf{y})\right) \\ &= \left(\tilde{\beta}_{m,i}^{(z+1)}\right) - \ln\left(\sum_{m=1}^M \exp\left(\tilde{\beta}_{m,i}^{(z+1)}\right)\right), \end{aligned} \quad (17)$$

in which we have $\tilde{\beta}_{m,i}^{(z+1)} \triangleq \beta_{m,i}^{(z+1)} - \gamma$, and the second term of the right-hand-side expression may be computed by invoking the ‘‘Jacobian logarithm’’ of [11]. Upon invoking the Max-log approximation, (17) may be further simplified as

$$\psi_{m,i}^{(z+1)} = \tilde{\beta}_{m,i}^{(z+1)} - \max_{m=1, \dots, M} \tilde{\beta}_{m,i}^{(z+1)}. \quad (18)$$

As a result, the estimated symbol-wise APP of s_i is given by

$$P^{(z+1)}(s_i = a_m|\mathbf{y}) \approx e^{\psi_{m,i}^{(z+1)}}, \quad (19)$$

which will replace the value of $P^{(z)}(s_i = a_m|\mathbf{y})$ in Table I. These updated symbol-wise APPs have to be converted to equivalent bit-wise LLRs, of which the *extrinsic* parts will be delivered to the outer FEC decoder of Fig. 1. Then, the extrinsic LLRs output by the FEC decoder will be converted to symbol-wise probabilities in the next iteration for the sake of generating new estimates of the symbol-wise APPs using the AB-Log-PDA. For reasons of explicit clarity, the proposed

TABLE II
SUMMARY OF THE AB-LOG-PDA ALGORITHM

Given the received signal \mathbf{y} , the channel matrix \mathbf{H} and the bit-wise *a priori* LLRs .

Step 1. Set the initial value of the iteration index to $z = 0$.

Step 2. Convert the *a priori* LLRs feedback by the FEC decoder to symbol-wise probabilities shown in Table I. When $z = 0$, the bit-wise *a priori* LLRs are all zeros, and the values of the symbol-wise probabilities are actually initialized as $P^{(z)}(s_i = a_m | \mathbf{y}) = 1/M$, for $\forall i = 1, 2, \dots, N_t$ and $\forall m = 1, 2, \dots, M$.

Step 3. Based on the values of $\{P^{(z)}(k)\}_{k \neq i}$, calculate $P^{(z+1)}(s_i = a_m | \mathbf{y})$ by

for $i = 1 : N_t$

calculate the statistics of the interference-plus-noise term \mathbf{v}_i using (4) - (9), as well as the inverse of Λ_i in (14),

for $m = 1 : M$

calculate $P^{(z+1)}(s_i = a_m | \mathbf{y})$ using (12), (13), (17) and (19).

end

end

Step 4. Convert the symbol-wise probabilities $P^{(z+1)}(s_i = a_m | \mathbf{y})$ to bit-wise LLRs, of which the extrinsic parts are delivered to the outer FEC decoder.

Step 5. If the iteration index z has reached a given number of iterations, terminate the iteration. Otherwise, let $z = z + 1$ and return to Step 2.

AB-Log-PDA algorithm is summarized in Table II.

IV. EXTRINSIC LLR CALCULATION FOR AB-LOG-PDA IN FEC-CODED MIMO SYSTEMS

In a FEC-coded MIMO system employing the IDD scheme of Fig. 1, typically the extrinsic LLRs of the coded bits are exchanged between the soft-input soft-output (SISO) MIMO detector and the SISO FEC decoder. Therefore, in order to employ the AB-Log-PDA in the IDD scheme, the AB-Log-PDA has to output the *correct* extrinsic LLRs of the FEC coded bits, which is however, not quite as straightforward as it seems at first sight.

Denote the l -th bit of the i -th symbol s_i as b_{il} , then the *a posteriori* LLR of b_{il} based on the *true* symbol-wise APPs of $P(s_i = a_m | \mathbf{y})$ may be written as

$$\begin{aligned}
 L_D(b_{il} | \mathbf{y}) &= \ln \frac{P(b_{il} = +1 | \mathbf{y})}{P(b_{il} = -1 | \mathbf{y})} \\
 &= \ln \frac{\sum_{\forall a_m \in \mathcal{A}_i^+} P(s_i = a_m | \mathbf{y})}{\sum_{\forall a_m \in \mathcal{A}_i^-} P(s_i = a_m | \mathbf{y})} \\
 &= L_E(b_{il} | \mathbf{y}) + \underbrace{\ln \frac{P(b_{il} = +1)}{P(b_{il} = -1)}}_{L_A(b_{il})}. \quad (20)
 \end{aligned}$$

In other words, we have

$$\begin{aligned}
 L_E(b_{il} | \mathbf{y}) &= \ln \frac{\sum_{\forall a_m \in \mathcal{A}_i^+} P(s_i = a_m | \mathbf{y})}{\sum_{\forall a_m \in \mathcal{A}_i^-} P(s_i = a_m | \mathbf{y})} \\
 &\quad - \underbrace{\ln \frac{P(b_{il} = +1)}{P(b_{il} = -1)}}_{L_A(b_{il})}. \quad (21)
 \end{aligned}$$

It is noteworthy that (21) represents a simple approach of generating the bit-wise extrinsic LLR of $L_E(b_{il} | \mathbf{y})$, as long as the *true* symbol APPs of $P(s_i = a_m | \mathbf{y})$ can be obtained.

Surprisingly, although we can directly obtain the estimated symbol APPs of $P(s_i = a_m | \mathbf{y})$ from the output of the AB-

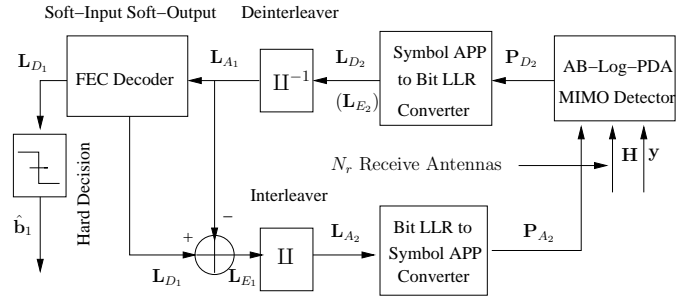


Fig. 2. Simplified structure of the AB-Log-PDA based IDD receiver, where we have $L_{E_2} = L_{D_2}$ rather than the classical $L_{E_2} = L_{D_2} - L_{A_2}$.

Log-PDA, as shown in (16), our study shows that this sort of estimated symbol APPs are not capable of yielding correct extrinsic bit-wise LLRs when invoking (21)¹. Therefore, the result of (16) should be regarded as “nominal” symbol APPs, rather than the “true” symbol APPs satisfying (21). Below we will show that it is still possible to obtain the extrinsic LLR directly based on the approximate Bayesian Theorem aided symbol APP, which is calculated using (19).

Conjecture 1. *The bit-wise extrinsic LLR of the AB-Log-PDA algorithm relying on (16) is given by*

$$L_E(b_{il} | \mathbf{y}) \approx \ln \frac{\sum_{\forall a_m \in \mathcal{A}_i^+} P(s_i = a_m | \mathbf{y})}{\sum_{\forall a_m \in \mathcal{A}_i^-} P(s_i = a_m | \mathbf{y})}, \quad (22)$$

where $P(s_i = a_m | \mathbf{y})$ is calculated by invoking (19).

The $L_E(b_{il} | \mathbf{y})$ values calculated from (22) using the “nominal” symbol APPs are typically not equivalent to $L_E(b_{il} | \mathbf{y})$ calculated from (21) using the “true” symbol APPs, but nonetheless, they constitute a good approximation of the latter without inducing any significant performance loss, as it will be demonstrated by our simulations in Section V. As a result, the classic IDD receiver structure of Fig. 1 is simplified to the structure of Fig. 2, where we have $L_{E_2} = L_{D_2}$, rather than $L_{E_2} = L_{D_2} - L_{A_2}$.

V. SIMULATION RESULTS AND COMPLEXITY ANALYSIS

In this section, both the performance and the computational complexity of the proposed AB-Log-PDA based IDD scheme are characterized, which further confirms the attractive performance versus complexity tradeoff achieved by the proposed scheme. The FEC employed is the parallel concatenated recursive systematic convolutional (RSC) code based turbo code having a coding rate of $R = \frac{k}{n} = 1/2$, constraint length of $L = 3$ and generator polynomials of (7, 5) in octal form. The turbo code is decoded by the Approximate-Log-MAP algorithm using $it_{tc} = 4$ inner iterations. The interleaver employed is the 2400-bit random sequence interleaver. The

¹In fact, if $L_E(b_{il} | \mathbf{y})$ is calculated by substituting the estimated symbol APPs of $P(s_i = a_m | \mathbf{y})$, i.e. the output of the AB-Log-PDA, into Eq. (21), the resultant BER curve of the IDD scheme of Fig. 1 exhibits an anti-bell shape when increasing SNR values. Due to the limitations of space, this BER curve is not presented in this paper.

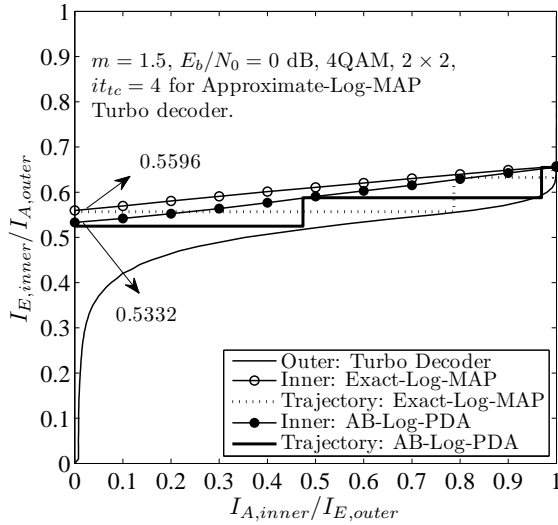


Fig. 3. EXIT chart analysis of the AB-Log-PDA and the Exact-Log-MAP based IDD schemes.

MIMO arrangement is represented in form of $(N_t \times N_r)$, and the Nakagami- m fading parameter is denoted by m .

A. Performance of the AB-Log-PDA based IDD

Fig. 3 compares the convergence behavior of the proposed AB-Log-PDA based IDD and that of the optimal Exact-Log-MAP based IDD scheme using EXIT chart [12] analysis, where the EXIT curve of the AB-Log-PDA is close to that of the Exact-Log-MAP. For example, when the *a priori* mutual information is $I_{A,inner} = 0$, the extrinsic mutual information of the AB-Log-PDA and of the Exact-Log-MAP is $I_{E,outer} = 0.5332$ and $I_{E,outer} = 0.5596$, respectively. This indicates that the performance of the AB-Log-PDA is close to that of the Exact-Log-MAP in the scenario considered. Additionally, the detection/decoding trajectories indicate that both the AB-Log-PDA and the Exact-Log-MAP based IDD schemes converge after three iterations, although the respective performance improvements achieved at each iteration are different.

The above EXIT chart based performance prediction and the convergence behavior of the IDD schemes considered are also characterized by the BER performance results of Fig. 4, where the Nakagami- m fading parameter is set to $m = 1.0$, which corresponds to the Rayleigh fading channel. Observe from Fig. 4 that the performance of the AB-Log-PDA based IDD scheme is improved upon increasing the number of outer iterations it_o , where $it_o = 0$ represents the conventional receiver structure in which the signal detector and the FEC decoder are serially concatenated, but operate without exchanging soft information. However, the attainable improvement gradually becomes smaller and the performance achieved after three outer iterations becomes almost the same as that of four outer iterations. This implies that the AB-Log-PDA based IDD scheme essentially converges after three outer iterations. A similar convergence profile is also observed for the optimal Exact-Log-MAP based IDD, although its performance is always marginally better than that of the

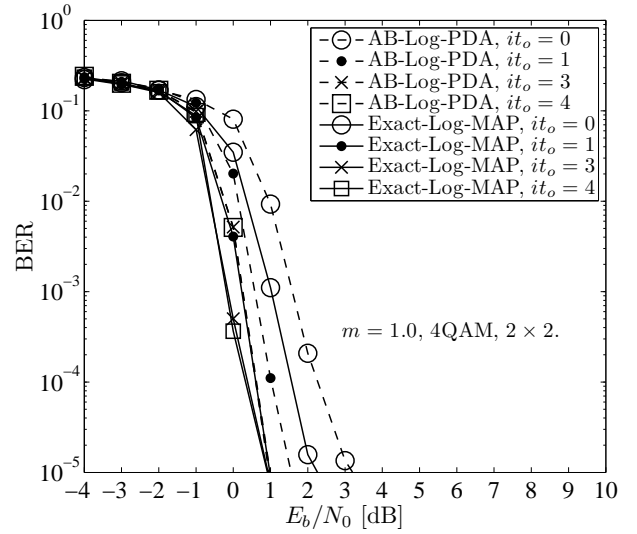


Fig. 4. Impact of the number of outer iterations on the achievable BER of the AB-Log-PDA and the Exact-Log-MAP based IDD schemes.

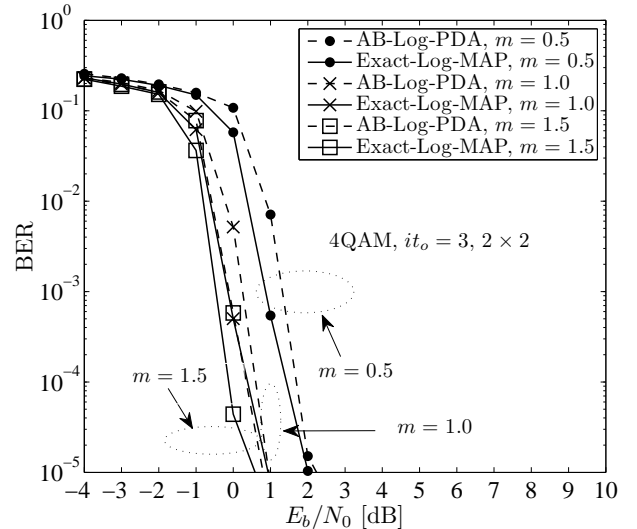


Fig. 5. Impact of Nakagami- m fading parameter m on the achievable BER of the AB-Log-PDA and the Exact-Log-MAP based IDD schemes.

corresponding AB-Log-PDA based IDD. Notably, both IDD schemes considered achieve $\text{BER} = 10^{-5}$ at about $E_b/N_0 = 1$ dB after three iterations.

Fig. 5 shows the impact of different m values on the achievable BER performance of the IDD schemes considered. As m decreases, the achievable performance of both the IDD schemes considered is degraded, since the fading becomes more severe. However, the performance gap between the AB-Log-PDA and the Exact-Log-MAP based IDD schemes is marginal for all values of m considered.

B. Complexity Analysis

Because the turbo codec module is common to both IDD schemes, and since we have shown that both the AB-Log-PDA and the Exact-Log-MAP based IDD schemes converge after

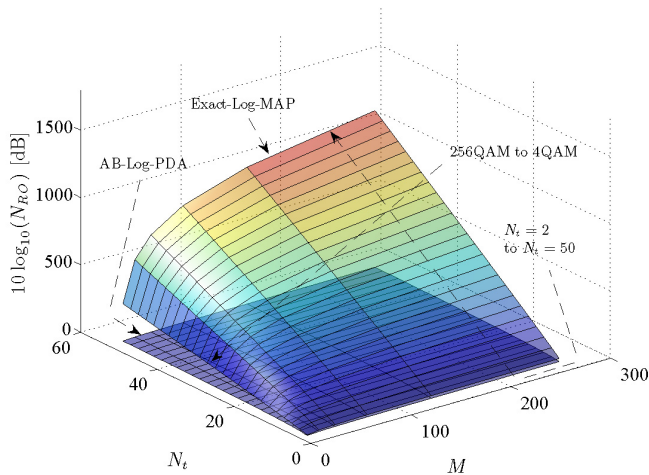


Fig. 6. Computational complexity comparison of the AB-Log-PDA and the Exact-Log-MAP algorithms in terms of the number of real operations N_{RO} .

three iterations in the scenarios considered, the computational complexity of the proposed AB-Log-PDA based IDD scheme can be evaluated by simply comparing its complexity to that of the Exact-Log-MAP in a single iteration. As shown in Table II, the major computational cost of the AB-Log-PDA per transmit symbol is the calculation of Λ_i^{-1} and the matrix multiplication of (13). By using the Sherman-Morrison-Woodbury formula based “speed-up” techniques of [2], the computational cost of calculating Λ_i^{-1} can be reduced to $\mathcal{O}(4N_t N_r^2)$ real operations (additions/multiplications) per iteration. Additionally, the calculation of (13) requires $\mathcal{O}(4MN_t N_r^2 + 2MN_t N_r)$ real operations per iteration. In summary, the computational complexity of the AB-Log-PDA method is $\mathcal{O}(4MN_t N_r^2 + 2MN_t N_r) + \mathcal{O}(4N_t N_r^2)$ per iteration. By comparison, the Exact-Log-MAP algorithm has to calculate the Euclidean distance $\|\mathbf{y} - \mathbf{H}\mathbf{s}\|^2$ M^{N_t} times, hence its complexity order is $\mathcal{O}(M^{N_t})$. More specifically, $\|\mathbf{y} - \mathbf{H}\mathbf{s}\|^2$ requires $\mathcal{O}(4N_r N_t + 6N_r)$ real operations. Therefore, the Exact-Log-MAP algorithm has a computational complexity of $\mathcal{O}[M^{N_t}(4N_r N_t + 6N_r)]$ real operations per iteration, which is significantly higher than that of the AB-Log-PDA, especially when N_t , N_r and M have large values. This observation is further confirmed by the results of Fig. 6, where the number of real operations is denoted by N_{RO} , while considering the scenario of $N_r = N_t$ as an example.

VI. CONCLUSIONS

In this paper, we showed that the classic relationship where the extrinsic LLRs are given by subtracting the *a priori* LLRs from the *a posteriori* LLRs is not valid for calculating the bit-wise extrinsic LLRs of the existing family of PDA based methods when employing M -QAM. Therefore, the PDA methods are not as readily applicable to the IDD receiver as they seem to be. This predicament is overcome by our novel approach of calculating the bit-wise extrinsic LLRs conceived for the proposed AB-Log-PDA MIMO detector. Moreover, we showed that the proposed AB-Log-PDA based IDD scheme is capable of achieving a comparable performance to that of the

optimal Exact-Log-MAP detector based IDD receiver, while imposing a significantly lower computational complexity in the scenarios considered.

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