

# Correspondence

## OFDM-Aided Differential Space-Time Shift Keying Using Iterative Soft Multiple-Symbol Differential Sphere Decoding

Mohammad Ismat Kadir, Sheng Chen, K. V. S. Hari,  
K. Giridhar, and Lajos Hanzo

**Abstract**—Soft-decision multiple-symbol differential sphere decoding (MSDSD) is proposed for orthogonal frequency-division multiplexing (OFDM)-aided differential space-time shift keying (DSTSK)-aided transmission over frequency-selective channels. Specifically, the DSTSK signaling blocks are generated by the channel-encoded source information and the space-time (ST) blocks are appropriately mapped to a number of OFDM subcarriers. After OFDM demodulation, the DSTSK signal is noncoherently detected by our soft-decision MSDSD detector. A novel soft-decision MSDSD detector is designed, and the associated decision rule is derived for the DSTSK scheme. Our simulation results demonstrate that an SNR reduction of 2 dB is achieved by the proposed scheme using an MSDSD window size of  $N_w = 4$  over the conventional soft-decision-aided differential detection benchmark, while communicating over dispersive channels and dispensing with channel estimation (CE).

**Index Terms**—Extrinsic information transfer (EXIT) chart, iterative decoding, multiple-symbol differential sphere decoding (MSDSD), orthogonal frequency-division multiplexing (OFDM), space-time shift keying (STSK).

### I. INTRODUCTION

Space-time shift keying (STSK) [1]–[3] has emerged as a beneficial multiple-input-multiple-output (MIMO) concept. STSK bridges the gap between the flexible diversity-multiplexing tradeoff provided by linear dispersion codes (LDCs) [4], [5] and the low-complexity design of spatial modulation (SM) [6]. Similar to the LDCs, STSK spreads the user information to both the spatial and time dimensions, but instead of simultaneously activating all the dispersion matrices (DMs), it transmits an additional  $\log_2 Q$  bits by activating one out of  $Q$  DMs. To overcome the performance degradation of the STSK scheme in wideband channels, orthogonal frequency-division multiplexing (OFDM)-aided STSK [7] and orthogonal frequency-division multiple-access/single-carrier frequency-division multiple-access-aided STSK [8] have also been proposed. The previous STSK studies [1], [2] demonstrate that coherent STSK performs well in conjunction with perfect channel state information (CSI) but exhibits a severe error floor in the presence of channel estimation (CE) errors.

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M. I. Kadir, S. Chen, and L. Hanzo are with the School of Electronics and Computer Science, University of Southampton, Southampton SO17 1BJ, U.K. (e-mail: mik1g09@ecs.soton.ac.uk; sqc@ecs.soton.ac.uk; lh@ecs.soton.ac.uk).

K. V. S. Hari is with the Department of Electrical Communication Engineering, Indian Institute of Science, Bangalore 560012, India (e-mail: hari@ece.iisc.ernet.in).

K. Giridhar is with the Department of Electrical Engineering, Indian Institute of Technology Madras, Chennai 600036, India (e-mail: giri@tenet.res.in).

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Differential STSK (DSTSK) employing conventional differential detection (CDD) has also been proposed for the sake of dispensing with the CE [2] and thus to eliminate the potentially high-Doppler-dependent pilot overhead. However, CDD suffers from a typical 3-dB performance penalty in low-Doppler scenarios. Furthermore, an irreducible error floor may be observed in a high-mobility scenario characterized by a high Doppler frequency. To circumvent the performance degradation of CDD, multiple-symbol differential detection (MSDD) was proposed for differential phase-shift keying (DPSK) in [9]. MSDD uses the fading-plus-noise statistics of the channel for jointly detecting  $(N_w - 1)$  information symbols from  $N_w$  number of consecutively received symbols, where  $N_w$  is usually referred to as the *observation window size*. The performance improvement of MSDD is, however, achieved at the cost of increased complexity, which increases exponentially with  $N_w$ . For mitigating this potentially excessive complexity, sphere decoding (SD) was invoked for MSDD in the context of multiple-symbol differential sphere decoding (MSDSD) in [10] and [11]. Hard-decision MSDSD was conceived in [12] for a DSTSK scheme operating in nondispersive channels. As a further advance, inspired by the near-capacity performance of turbo detection [13], [14], a soft-decision MSDSD scheme was also designed for DPSK in [15]. Furthermore, the concept of differential space-frequency modulation employing MSDSD in conjunction with a specific subcarrier allocation was proposed in [16] for exploiting both the achievable spatial- and frequency-domain diversity. However, the conception of the soft-decision-MSDSD-aided DSTSK designed for realistic dispersive scenarios constitutes an unexplored open problem.

Against this background, we conceive a novel soft-decision MSDSD for OFDM-based DSTSK operating in frequency-selective channels. The main contributions of this paper are as follows.

- 1) A novel soft-decision-aided MSDSD is proposed for OFDM-aided DSTSK operating in dispersive channels. The decision rule of the soft-decision MSDSD is deduced by considering the construction of DSTSK codewords based on the DMs, the Doppler frequency, the OFDM system parameters, and the generation of soft information.
- 2) A lower bound of the detection complexity is deduced, which is verified by simulations.

The remainder of this paper is organized as follows. In Section II, an overview of the proposed channel-coded OFDM-aided DSTSK scheme is provided. The soft-decision MSDSD is modeled in Section III. In Section IV, both the complexity imposed by the system is quantified. The performance of the soft-decision MSDSD-aided DSTSK scheme is investigated in Section V. Finally, we conclude in Section VI.

**Notations:** We use capital boldface letters to denote matrices, whereas  $\{\cdot\}^T$ ,  $\{\cdot\}^H$ ,  $\text{tr}(\cdot)$ ,  $\det[\cdot]$ , and  $\|\cdot\|$  are used to represent the transpose, the Hermitian transpose, the trace, the determinant, and the Euclidean norm of the matrix “ $\cdot$ ,” respectively. The notations  $\mathcal{E}\{\cdot\}$ ,  $\cdot^*$ , and  $P(\cdot)$  are used to denote the expected value, the complex conjugate, and the probability of “ $\cdot$ ” respectively, whereas  $\otimes$  and  $\mathbf{I}_T$  represent the Kronecker product and the  $(T \times T)$  identity matrix, respectively. A symmetric  $(N_w \times N_w)$  Toeplitz matrix is denoted  $\text{toeplitz}\{x_1, \dots, x_{N_w}\}$ , whereas  $\text{diag}\{\mathbf{X}_1, \dots, \mathbf{X}_{N_w}\}$  indicates a block-diagonal matrix with the matrices  $\mathbf{X}_1, \dots, \mathbf{X}_{N_w}$  on its main diagonal. Furthermore,  $\mathcal{CN}(\mu, \sigma^2)$  refers to the circularly symmetric complex Gaussian distribution with mean  $\mu$  and variance  $\sigma^2$ .

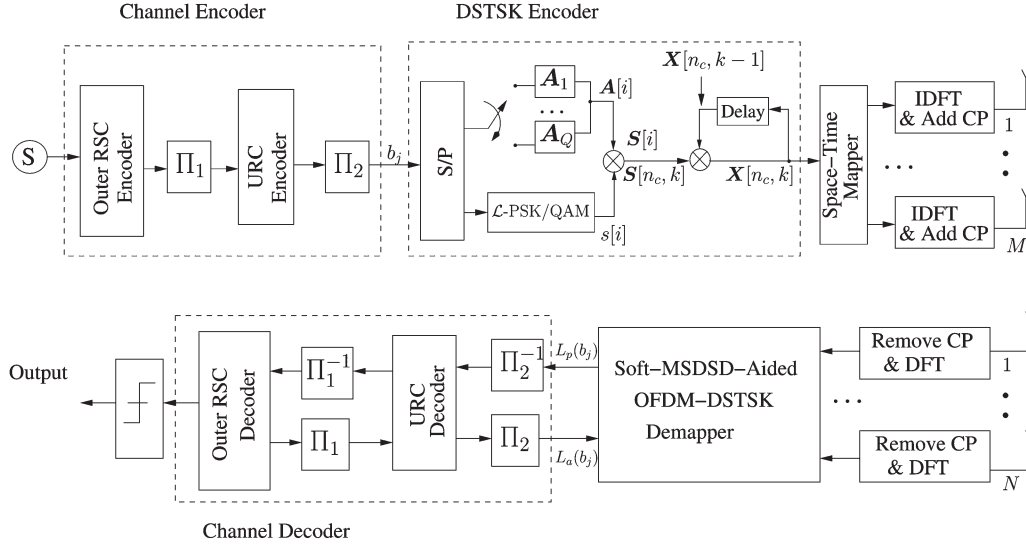


Fig. 1. Transceiver architecture of the proposed concatenated channel-coding-aided DSTSK scheme relying on the soft-decision MSDSD as the inner decoder.

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## II. SYSTEM MODEL

We consider a channel-coded OFDM-aided DSTSK transceiver employing  $M$  transmit and  $N$  receive antenna elements (AEs), as shown in Fig. 1. The channel encoder/decoder blocks of Fig. 1 may incorporate a general channel coding scheme that supports soft-decision decoding at affordable complexity. A pragmatic coding architecture might be an appropriately interleaved serially concatenated recursive convolutional code (RSC) and unity rate code (URC)-aided scheme [17]–[19], as shown in Fig. 1.

The source bits are first channel encoded by the RSC code, and the encoded bits are then interleaved by a random bit interleaver  $\Pi_1$ . Following URC precoding, the interleaved bits are further interleaved by a second interleaver  $\Pi_2$ . The resultant bits are then mapped to STSK codewords, which are further mapped to  $N_c$  parallel subcarriers and then differentially encoded in the time domain (TD), i.e., across the consecutive OFDM symbols of the same subcarrier. The DSTSK codewords are then OFDM modulated, while incorporating appropriate cyclic prefixes (CPs).

The signal received is first OFDM demodulated and then input to the DSTSK soft-decision MSDSD demapper. The extrinsic soft information is then iteratively exchanged between the three soft-in-soft-out components, namely, the DSTSK demapper, the URC decoder, and the RSC decoder, before finally outputting the estimated bits [8], [19].

### A. DSTSK Architecture and OFDM Layout

The STSK encoder generates space-time (ST) codewords from the source information by activating a single DM in any symbol duration in conjunction with the classic modulated symbols for transmission over  $T$  time slots using  $M$  transmit AEs [1], [2]. More specifically, each STSK signaling block  $\mathbf{S}[i] \in \mathbb{C}^{T \times M}$  is created from  $\log_2(\mathcal{L} \cdot Q)$  source bits according to [1], [2]

$$\mathbf{S}[i] = s[i] \mathbf{A}[i] \quad (1)$$

where  $s[i]$  is an  $\mathcal{L}$ -ary constellation symbol represented by  $\log_2 \mathcal{L}$  bits, and  $\mathbf{A}[i] \in \mathbb{C}^{T \times M}$  is the specific DM activated from the set of  $Q$  DMs  $\mathbf{A}_q (q = 1, \dots, Q)$ , as determined by the remaining  $\log_2 Q$  bits. The DMs  $\mathbf{A}_q (q = 1, \dots, Q)$  are unitary matrices generated by employing an exhaustive search for minimizing the objective function constituted by the pairwise error probability of the codewords [5], [12], [20], [21] under the power constraint in [2] expressed by  $\text{tr}(\mathbf{A}_q^H \mathbf{A}_q) =$

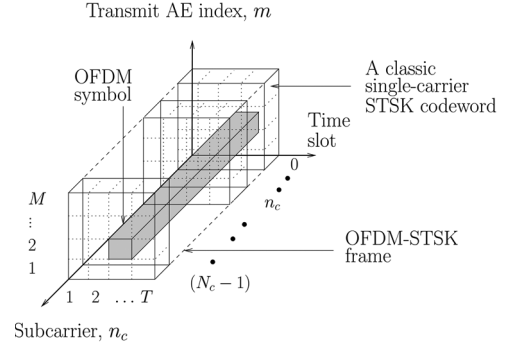


Fig. 2. Mapping of the STSK codewords to  $N_c$  parallel OFDM subcarriers showing the construction of an OFDM-STSK frame and OFDM symbols. After being appropriately mapped to the subcarriers, the codewords are differentially encoded in the TD and transmitted over dispersive channels by  $M$  transmit AEs over  $T$  time-slots.

$T \forall q$ . The resultant STSK system is then uniquely and unambiguously described by the parameters  $(M, N, T, Q, \mathcal{L})$ .

We observe that the STSK codeword  $\mathbf{S}[i]$  belongs to a set  $\mathcal{S}$  of  $(\mathcal{L} \cdot Q)$  codeword matrices defined by

$$\mathcal{S} \triangleq \{s_l \mathbf{A}_q | (q \in \{1, \dots, Q\}, l \in \{1, \dots, \mathcal{L}\})\}. \quad (2)$$

The STSK codewords are mapped to  $N_c$  parallel subcarriers, as shown in Fig. 2, before being differentially encoded. As shown in Fig. 2,  $N_c$  consecutive STSK codewords are arranged in parallel to form an OFDM-STSK frame, and OFDM modulation is carried out over each shaded symbol pipe, which constitutes an OFDM symbol. We may represent the codeword  $\mathbf{S}[i]$  by  $\mathbf{S}[n_c, k]$ , so that the overall codeword index  $i$  is related to the OFDM frame index  $k$  and the subcarrier index  $n_c$  by  $i = kN_c + n_c, n_c = 0, 1, \dots, (N_c - 1)$ . Additionally, we invoke differential encoding in the TD, i.e., differential encoding is performed across the consecutive OFDM symbols of the same subcarrier. We have chosen TD differential encoding/decoding because we have conceived our scheme for continuous transmissions as opposed to the FD differential encoding/decoding across adjacent FD subcarriers, which is more suitable for burst transmissions. To facilitate convenient differential encoding, we assume  $M = T$ . Furthermore, directly generated unitary DMs are used in the proposed scheme for avoiding the nonlinear Cayley transform [2], [12]. The codewords

156  $\mathbf{S}[n_c, k]$  are thus differentially encoded to form the transmit blocks  
 157  $\mathbf{X}[n_c, k]$  ( $k = 0, 1, 2, \dots$ ) according to [16]

$$\mathbf{X}[n_c, k] = \begin{cases} \mathbf{X}[n_c, k-1]\mathbf{S}[n_c, k], & k = 1, 2, \dots \\ \mathbf{I}_T & k = 0. \end{cases} \quad (3)$$

158 The DSTSK codewords are then transmitted after the  $N_c$ -point inverse  
 159 discrete Fourier transform (DFT) operation and appropriate CP incor-  
 160 poration.

#### 161 B. Channel Model

162 Each link between the  $m$ th transmit and  $n$ th receive AE is assumed  
 163 a frequency-selective channel, but as a benefit of OFDM-based trans-  
 164 mission, each dispersive channel is then partitioned into  $N_c$  low-rate  
 165 parallel frequency-flat subchannels [22]. The complex-valued fading  
 166 gain  $h_{m,n}[n_c, k]$  ( $m = 1, 2, \dots, M$ ;  $n = 1, 2, \dots, N$ ) obeys the dis-  
 167 tribution  $\mathcal{CN}(0, 1)$  associated with an autocorrelation function based  
 168 on Clarke's model [23]:  $\varphi_{hh}[n_c, \kappa] \triangleq \mathcal{E}\{h_{m,n}[n_c, k]h_{m,n}^*[n_c, k +$   
 169  $\kappa]\} = J_0(2\pi\kappa f_d)$ , where  $J_0$  denotes the zeroth-order Bessel function  
 170 of the first kind, and  $f_d = f_m T$  is the normalized maximum Doppler  
 171 frequency, whereas  $f_m$  and  $1/T$  represent the maximum Doppler  
 172 frequency and the symbol rate, respectively. The fading is assumed  
 173 quasi-static, i.e., the channel's complex-valued envelope remains ap-  
 174 proximately constant during the transmission of an OFDM STSK  
 175 frame.

176 Given the aforementioned assumptions, the received signal  
 177  $\mathbf{Y}[n_c, k] \in \mathbb{C}^{T \times N}$  obtained after CP removal and DFT may be ex-  
 178 pressed by [16], [22]

$$\mathbf{Y}[n_c, k] = \mathbf{X}[n_c, k]\mathbf{H}[n_c, k] + \mathbf{V}[n_c, k] \quad (4)$$

179 where  $\mathbf{X}[n_c, k] \in \mathbb{C}^{T \times M}$  represents the codeword transmitted and  
 180  $\mathbf{H}[n_c, k] \in \mathbb{C}^{M \times N}$  denotes the FD channel transfer matrix, with  
 181 its  $(m, n)$ th entry given by  $h_{m,n}[n_c, k]$ . Furthermore,  $\mathbf{V}[n_c, k] \in$   
 182  $\mathbb{C}^{T \times N}$  is the additive white Gaussian noise (AWGN) with entries of  
 183  $v_{T_i, n}[n_c, k] \sim \mathcal{CN}(0, \sigma_n^2)$ .

### 184 III. MSDSD RECEIVER

185 This section introduces the maximum-likelihood MSDSD (ML-  
 186 MSDSD), the maximum *a posteriori* MSDSD (MAP-MSDSD) algo-  
 187 rithm, and the generation of the log-likelihood ratios (LLRs) for the  
 188 soft-decision-MSDSD-aided OFDM DSTSK.

#### 189 A. ML-MSDSD for OFDM-Aided DSTSK

190 The ML-MSDD processes  $N_w$  consecutively received space-time  
 191 blocks corresponding to the  $n_c$ th subcarrier given by  $\bar{\mathbf{Y}}[n_c, k] \triangleq$   
 192  $[\mathbf{Y}^T[n_c, k - N_w + 1], \dots, \mathbf{Y}^T[n_c, k]]^T$  and finds the ML estimates  
 193  $\hat{\bar{\mathbf{X}}}[n_c, k]$  of the corresponding  $N_w$  transmitted blocks  $\bar{\mathbf{X}}[n_c, k] \triangleq$   
 194  $[\mathbf{X}^T[n_c, k - N_w + 1], \dots, \mathbf{X}^T[n_c, k]]^T$  [24]. Since the differen-  
 195 tially encoded blocks  $\mathbf{X}[n_c, k]$  are related to the STSK codewords  
 196 by the one-to-one relationship expressed by (3), the ML-MSDD  
 197 in turn estimates  $(N_w - 1)$  STSK codewords given by  $\bar{\mathbf{S}}[n_c, k] \triangleq$   
 198  $[\mathbf{S}^T[n_c, k - N_w + 2], \dots, \mathbf{S}^T[n_c, k]]^T$ , which further estimates the  
 199 source bits mapped to the STSK codewords.

200 Defining a block-diagonal matrix  $\bar{\mathbf{X}}_D[n_c, k]$  by  $\bar{\mathbf{X}}_D[n_c, k] \triangleq$   
 201  $\text{diag}\{\mathbf{X}[n_c, k - N_w + 1], \dots, \mathbf{X}[n_c, k]\}$ ,  $\bar{\mathbf{H}}[n_c, k] \triangleq [\mathbf{H}^T[n_c, k -$   
 202  $N_w + 1], \dots, \mathbf{H}^T[n_c, k]]^T$ , and  $\bar{\mathbf{V}}[n_c, k] \triangleq [\mathbf{V}^T[n_c, k - N_w + 1],$   
 203  $\dots, \mathbf{V}^T[n_c, k]]^T$ , the  $N_w$ -block received sequence can be expressed  
 204 by [11], [24]

$$\bar{\mathbf{Y}}[n_c, k] = \bar{\mathbf{X}}_D[n_c, k]\bar{\mathbf{H}}[n_c, k] + \bar{\mathbf{V}}[n_c, k] \quad (5)$$

where

$$\begin{aligned} \bar{\mathbf{Y}}[n_c, k] &\in \mathbb{C}^{N_w M \times N} & \bar{\mathbf{V}}[n_c, k] &\in \mathbb{C}^{N_w M \times N} \\ \bar{\mathbf{H}}[n_c, k] &\in \mathbb{C}^{N_w M \times N} & \bar{\mathbf{X}}_D[n_c, k] &\in \mathbb{C}^{N_w M \times N_w M}. \end{aligned}$$

For the sake of notational simplicity, we omit the subcarrier index and  
 time index  $[n_c, k]$  in the following and refer to the  $\mu$ th submatrix of a  
 block matrix, e.g.,  $\mathbf{B}$  by the subscripted matrix  $\mathbf{B}_\mu$ . Under the assump-  
 tion that  $h_{m,n}$  and  $v_{T_i, n}$  ( $m = 1, 2, \dots, M$ ,  $n = 1, 2, \dots, N$ ) are  
 zero-mean Gaussian random processes, the probability density  
 function of  $\bar{\mathbf{Y}}$  conditioned on  $\bar{\mathbf{X}}_D$  is given by [25], [26]

$$P(\bar{\mathbf{Y}}|\bar{\mathbf{X}}_D) = \frac{1}{(\pi^{N_w M} \det[\mathbf{\Lambda}_Y])^N} \exp \left\{ -\text{tr} \left( \bar{\mathbf{Y}}^H \mathbf{\Lambda}_Y^{-1} \bar{\mathbf{Y}} \right) \right\} \quad (6)$$

where  $\mathbf{\Lambda}_Y$  is defined by  $\mathbf{\Lambda}_Y \triangleq \mathcal{E}\{\bar{\mathbf{Y}}\bar{\mathbf{Y}}^H|\bar{\mathbf{X}}_D\}$ . The ML estimate  $\hat{\bar{\mathbf{X}}}$   
 under the assumption of quasi-static fading and unitary  $\bar{\mathbf{X}}_D$  reduces  
 to [26], [27]

$$\hat{\bar{\mathbf{X}}} = \arg \max_{\bar{\mathbf{X}}} P(\bar{\mathbf{Y}}|\bar{\mathbf{X}}_D) = \arg \min_{\bar{\mathbf{X}}} \left\{ \text{tr} \left( \bar{\mathbf{Y}}^H \mathbf{\Lambda}_Y^{-1} \bar{\mathbf{Y}} \right) \right\}. \quad (7)$$

Here, the conditional covariance matrix  $\mathbf{\Lambda}_Y$  is related to the channel  
 parameters [25], [26] by

$$\mathbf{\Lambda}_Y^{-1} = \frac{1}{N} \bar{\mathbf{X}}_D (\mathbf{\Lambda}^{-1} \otimes \mathbf{I}_M) \bar{\mathbf{X}}_D^H \quad (8)$$

where we have  $\mathbf{\Lambda} \triangleq (\psi_{hh} + \sigma_n^2 \mathbf{I}_{N_w})$  and  $\psi_{hh} \triangleq \text{toeplitz}\{\varphi_{hh}[n_c, 217$   
 $0], \dots, \varphi_{hh}[n_c, (N_w - 1)]\}$ , with the component autocorrelation  
 functions  $\varphi_{hh}[n_c, \kappa]$  being identical for all spatial channels. Applying  
 the Cholesky factorization of  $\mathbf{\Lambda}^{-1} = \mathbf{U}^H \mathbf{U}$  with the upper triangular  
 matrix  $\mathbf{U}$  and considering the identity  $\text{tr}(\mathbf{X}\mathbf{X}^H) = \|\mathbf{X}\|^2$  for any  
 matrix  $\mathbf{X}$ , the ML-MSDD decision rule can be deduced from (6),  
 yielding [11]

$$\hat{\bar{\mathbf{X}}} = \arg \min_{\bar{\mathbf{X}}} \left\{ \sum_{\mu=1}^{N_w} \left\| \mathbf{Y}_{\mu, \mu}^H \bar{\mathbf{X}}_\mu + \sum_{\nu=\mu+1}^{N_w} (\mathbf{Y}_{\mu, \nu}^H \bar{\mathbf{X}}_\nu) \right\|^2 \right\} \quad (9)$$

where  $\mathbf{Y}_{\mu, \nu}^H$  is defined by  $\mathbf{Y}_{\mu, \nu}^H \triangleq \mathbf{Y}_{\nu}^H u_{\mu, \nu}$ , and  $u_{\mu, \nu}$  represents the  
 $(\mu, \nu)$ th element of  $\mathbf{U}$ . Still referring to (9),  $\hat{\bar{\mathbf{X}}}$  denotes the estimate  
 of  $\bar{\mathbf{X}}_D$ , whereas  $\bar{\mathbf{X}}_\mu$  refers to the  $\mu$ th candidate submatrix of  $\bar{\mathbf{X}}_D$ .

Since the ML metric of (9) is invariant to a phase shift common to all  
 elements of  $\bar{\mathbf{X}}_\mu \forall \mu$  corresponding to the same  $\bar{\mathbf{S}}$  (where  $\mathbf{S}_\lambda^T \in \mathcal{S} \forall \lambda$ ),  
 the accumulated differential matrices may be expressed as [15]

$$\mathcal{A}_\nu = \begin{cases} \prod_{\lambda=\nu}^{N_w-1} \mathbf{S}_\lambda^H, & 1 \leq \nu \leq (N_w - 1) \\ \mathbf{I}_T, & \nu = N_w. \end{cases} \quad (10)$$

For the sake of reducing the complexity associated with an exhaustive  
 search, we employ MSDSD similar to [10] and [11] to search through  
 the candidate set lying within a sphere of radius  $\rho_s$  as follows:

$$\sum_{\mu=1}^{N_w} \left\| \mathbf{Y}_{\mu, \mu}^H \mathcal{A}_\mu + \sum_{\nu=\mu+1}^{N_w} (\mathbf{Y}_{\mu, \nu}^H \mathcal{A}_\nu) \right\|^2 \leq \rho_s^2. \quad (11)$$

#### 233 B. MAP-MSDSD

Assuming the STSK codewords to be mutually independent, (6) and  
 (9) yield [15]:

$$\begin{aligned} & -\ln(P(\bar{\mathbf{S}}|\bar{\mathbf{Y}})) \\ & \propto -\ln(P(\bar{\mathbf{Y}}|\bar{\mathbf{S}})) - \ln(P(\bar{\mathbf{S}})) \\ & \propto \sum_{\mu=1}^{N_w} \left\{ \left\| \mathbf{Y}_{\mu, \mu}^H \mathcal{A}_\mu + \sum_{\nu=\mu+1}^{N_w} \mathbf{Y}_{\mu, \nu}^H \mathcal{A}_\nu \right\|^2 - \ln(P(\mathbf{S}_\mu)) \right\}. \end{aligned} \quad (12)$$



236 The MAP-MSDSD may be thus expressed as

$$\sum_{\mu=1}^{(N_w-1)} \left( \left\| \sum_{\nu=\mu}^{N_w} (\mathbf{Y}_{\mu,\nu}^H \mathbf{A}_{\nu}) \right\|^2 - \ln(P(\mathbf{S}_{\mu})) \right) \leq \rho_s^2 - \|u_{N_w, N_w} \mathbf{Y}_{N_w}\|^2 \triangleq \rho^2. \quad (13)$$

237 Clearly, the codeword  $\mathbf{S}_{\mu}$  obeys the specific distance criterion [10],  
238 [12], [28] that the current partial Euclidean distance (PED)  $d_{\mu}^2$  is the  
239 sum of the previous PED  $d_{\mu+1}^2$  and the distance increment  $\Delta_{\mu}^2$ , i.e.,

$$\begin{aligned} d_{\mu}^2 &\triangleq \Delta_{\mu}^2 + d_{\mu+1}^2 \\ &= \left\| u_{\mu,\mu} \mathbf{Y}_{\mu} \mathbf{A}_{\mu+1} \mathbf{S}_{\mu}^H + \sum_{\nu=\mu+1}^{N_w} u_{\mu,\nu} \mathbf{Y}_{\nu} \mathbf{A}_{\nu} \right\|^2 - \ln(P(\mathbf{S}_{\mu})) \\ &\quad + \sum_{\iota=\mu+1}^{(N_w-1)} \left( \left\| \sum_{\nu=\iota}^{N_w} u_{\iota,\nu} \mathbf{Y}_{\nu} \mathbf{A}_{\nu} \right\|^2 - \ln(P(\mathbf{S}_{\iota})) \right) \leq \rho^2. \quad (14) \end{aligned}$$

240 Similar to the MSDSD principle described in [10] and [15], the  
241 MAP-MSDSD is initialized with  $\mu = (N_w - 1)$  and then proceeds by  
242 applying the search criterion of (14) until  $\mu = 1$ , where the search ra-  
243 dius is updated to  $\rho^2 = d_1^2$ , and the search is repeated by commencing  
244 from  $\mu = 2$  until  $\mu = (N_w - 1)$  is reached. If the new search does not  
245 provide a better estimate, the previous estimate is retained.

#### 246 C. Log-Likelihood Ratio and Soft-Decision-MSDSD-Aided 247 OFDM DSTSK

248 The soft demapper relies on the *a priori* information gleaned from  
249 the URC decoder and the MAP-MSDSD. A high interleaver depth is  
250 assumed so that the permuted bits may be treated as being independent.  
251 The LLR corresponding to the bit  $b_j$  interleaved by the interleaver  
252  $\Pi_2$  of Fig. 1 is defined by [29]  $L_a(b_j) \triangleq \ln(P(b_j = b)/P(b_j = \bar{b}))$ ,  
253 where  $b \in \{0, 1\}$ , and the  $j$ th bit  $b_j = b$  corresponds to the MAP-  
254 MSDSD estimate  $\hat{\mathbf{S}}$ , whereas  $\bar{b}$  indicates its complement. The  
255 *a posteriori* LLR  $L_p(\cdot)$  of  $b_j$  may be then approximated by the  
256 maximum-logarithmic-MAP (max-log-MAP) algorithm [14], [15], i.e.,

$$\begin{aligned} L_p(b_j) &= \ln \frac{P(b_j = b|\mathbf{Y})}{P(b_j = \bar{b}|\mathbf{Y})} \\ &\approx \ln \frac{\max_{\hat{\mathbf{S}}: b_j = b} \left[ -\sum_{\mu=1}^{N_w} \left\| \sum_{\nu=\mu}^{N_w} (\mathbf{Y}_{\mu,\nu}^H \hat{\mathbf{A}}_{\nu}) \right\|^2 + \ln(P(\hat{\mathbf{S}})) \right]}{\max_{\hat{\mathbf{S}}: b_j = \bar{b}} \left[ -\sum_{\mu=1}^{N_w} \left\| \sum_{\nu=\mu}^{N_w} (\mathbf{Y}_{\mu,\nu}^H \hat{\mathbf{A}}_{\nu}) \right\|^2 + \ln(P(\hat{\mathbf{S}})) \right]} \\ &= -\sum_{\mu=1}^{N_w} \left\| \sum_{\nu=\mu}^{N_w} (\mathbf{Y}_{\mu,\nu}^H \hat{\mathbf{A}}_{\nu}) \right\|^2 + \ln(P(\hat{\mathbf{S}}^b)) \\ &\quad + \sum_{\mu=1}^{N_w} \left\| \sum_{\nu=\mu}^{N_w} (\mathbf{Y}_{\mu,\nu}^H \hat{\mathbf{A}}_{\nu}) \right\|^2 - \ln(P(\hat{\mathbf{S}}^{\bar{b}})) \quad (15) \end{aligned}$$

257 where  $\hat{\mathbf{S}}^b$  and  $\hat{\mathbf{S}}^{\bar{b}}$  represent the MAP-MSDSD estimate and the  
258 constrained estimate associated with  $b_j = b$ , respectively.

The extrinsic LLR  $L_e(\cdot)$  for  $b_j$  is now evaluated by combining the  
*a posteriori* and *a priori* LLR:  $L_e(b_j) = L_p(b_j) - L_a(b_j)$ . The ex-  
trinsic information extracted from the soft-decision MSDSD demapper  
is iteratively exchanged with the URC decoder of Fig. 1, which forms  
the *inner* iteration, whereas the exchange of extrinsic information  
between the URC decoder and the RSC decoder of Fig. 1 may be  
termed as the *outer* iteration. Note that, for each outer iteration be-  
tween the RSC decoder and the URC decoder, several inner iterations  
may be invoked between the URC and the soft-decision-MSDSD-  
aided DSTSK demapper [2], [19]. Finally, the RSC decoder generates  
*a posteriori* LLRs, from which the source bits are estimated.

#### IV. COMPLEXITY

Here, the complexity of the proposed scheme is detailed, and the  
complexity imposed by the MAP-MSDSD is quantified.

Equation (2) shows that there exists  $(\mathcal{L} \cdot Q)$  legitimate code-  
word matrices for each  $\log_2(\mathcal{L} \cdot Q)$  bits of source information. The  
exhaustive-search-based solution to (9) involves a search in a  $(\mathcal{L} \cdot Q)^{(N_w-1)}$   
element space of candidate matrices  $\hat{\mathbf{X}}$  corresponding to all possible  
choices of  $\hat{\mathbf{S}}$ . The ML-MSDSD associated with chosen sphere  
radius  $\rho$  imposes average complexity, which is lower bounded by [11]

$$C \geq \frac{(\mathcal{L} \cdot Q)^{N_w \zeta - 1} - (\mathcal{L} \cdot Q)}{(\mathcal{L} \cdot Q) - 1} \quad (16)$$

where

$$\zeta \triangleq \frac{\sigma_n^2(1 + \epsilon)}{2(1 + \sigma_n^2)} \quad \rho^2 = (1 + \epsilon)NMN_w, \quad \epsilon > 0. \quad (17)$$

To quantify the complexity of the MAP-MSDSD scheme, we con-  
sider the number of real-valued multiplication operations (RMOs)  
required for obtaining a single soft-output value, which is used as  
our complexity metric. The lower bound of the complexity may be  
obtained if the number of RMOs required for computing the soft  
outputs corresponding to the first codeword estimate  $\hat{\mathbf{S}}^b$  is counted  
and if a single constrained estimate  $\hat{\mathbf{S}}^{\bar{b}}$  is taken into account [15].  
Considering the upper diagonal nature of the matrix  $\mathbf{U}$ , we observe  
that  $\mathbf{Y}_{\mu,\nu}^H$  is defined only for  $\nu \geq \mu$  in the context of (9), although each  
 $\mathbf{Y}_{\mu,\nu}^H$  is an  $(N \times T)$ -element matrix, where  $T = M$ . The computation  
of the  $\mathbf{Y}_{\mu,\nu}^H$  terms in (15) thus involves a total of  $2MN[1 + 2 + \dots + N_w] = MN N_w(N_w + 1)$  RMOs, assuming real-valued autocorrela-  
tion functions of  $\varphi_{hh}[n_c, \kappa]$ . To compute the *a posteriori* LLRs given  
by (15), the number of RMOs associated with the computation of each  
 $\left\| \sum_{\nu=\mu}^{N_w} (\mathbf{Y}_{\mu,\nu}^H \hat{\mathbf{A}}_{\nu}) \right\|^2$  is  $4M^2N(N_w - \mu + 1) + 2$ . The number of  
RMOs required for generating  $\log_2(\mathcal{L} \cdot Q)$  soft outputs corresponding  
to a single codeword estimate  $\hat{\mathbf{S}}^b$  is thus given by

$$\begin{aligned} \text{RMO}[\hat{\mathbf{S}}^b] &= MN N_w(N_w + 1) + \sum_{\mu=1}^{N_w} [4M^2N(N_w - \mu + 1) + 2] \\ &= MN(2M + 1)N_w(N_w + 1) + 2N_w. \quad (18) \end{aligned}$$

On the other hand, the number of RMOs related to each bit  $\bar{b}$  of  
the constrained estimate  $\hat{\mathbf{S}}^{\bar{b}}$  is found to be  $2N_w[M^2N(N_w + 1) + 1]$ .  
The lower bound for the number of RMOs associated with the gen-  
eration of a single soft output is thus  $[MN(2M + 1)/\log_2(\mathcal{L} \cdot Q) + 300]$   
 $2M^2N]N_w^2$  for large  $N_w$ . The complexity of the scheme, however,  
depends on a number of parameters, such as on the channel SNR, on  
the autocorrelation function of the channel's fading plus noise, and  
most importantly, on the *a priori* mutual information  $I_A$  of the inner  
decoder [15]. In Section V, the complexity of the MAP-MSDSD will  
be investigated as a function of the observation window width  $N_w$   
parameterized by the available *a priori* information  $I_A$ .

TABLE I  
ADOPTED VALUES OF MAIN SIMULATION PARAMETERS

Parameter	Value
Dispersive channel model	COST207-TU12
Fast fading envelope	Correlated Rayleigh fading
Normalized Doppler spread, $f_d$	0.01
Number of subcarriers, $N_c$	128
Length of cyclic prefix	32
Overall symbol duration	300 ns
STSK ( $M, N, T, Q, \mathcal{L}$ )	(2, 2, 2, 4, 4)
RSC encoder and decoder	Half rate
	Constraint length=2
Generator polynomial	(011, 010) <sub>2</sub>
Length of interleavers	200,000 bits
Outer iterations	11
Inner iterations	2

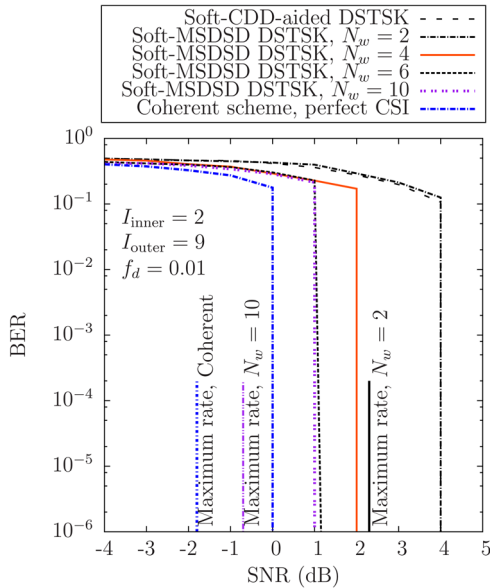


Fig. 3. Simulated BER performance of the soft-decision-MSDSD-aided OFDM DSTSK (2, 2, 2, 4, 4) scheme for transmission over dispersive COST207-TU12 channel with normalized Doppler frequency  $f_d = 0.01$  and different observation window size  $N_w = 2, 4, 6, 10$ . The BER falls sharply after  $I_{\text{outer}} = 9$  outer iterations as a benefit of employing the URC and the performance approaches that of the coherent scheme with perfect CSI with an increasing value of  $N_w$ .

## V. PERFORMANCE RESULTS

Here, the performance of the proposed scheme is investigated using the parameters listed in Table I. We have employed the COST207-TU12 channel model for the links between each transmit-receive antenna pair. The power delay profile characterizing the 12 taps of the COST207-TU12 channel is detailed in [30] and [31]. As mentioned in Table I, we employ an RSC (2, 1, 2) outer code having octally represented generator polynomials of  $(g_r, g) = (3, 2)_8$  as well as two random interleavers with a length of 200 000 bits.

Fig. 3 characterizes the achievable bit error rate (BER) of the proposed soft-decision-MSDSD-aided OFDM DSTSK scheme associated with observation window sizes of  $N_w = 2, 4, 6, 10$  and compares it to that achievable by the corresponding coherent scheme relying on perfect CSI. We observe that the proposed scheme has the benefit of dispensing with CE due to differential encoding, while mitigating the performance erosion of classic STSK by employing OFDM. Further-

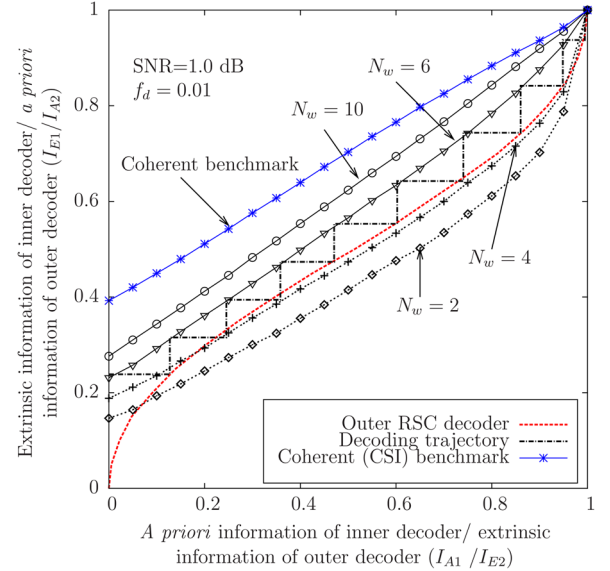


Fig. 4. EXIT charts of the inner decoders of the soft-decision-MSDSD-aided OFDM DSTSK (2, 2, 2, 4, 4) scheme at SNR = 1 dB with normalized Doppler frequency  $f_d = 0.01$  and different observation window sizes  $N_w = 2, 4, 6, 10$  and of the corresponding coherent system having perfect CSI at the receiver. At this SNR, we observe the inner EXIT charts with  $N_w = 6, 10$  have an open EXIT tunnel and converge to the (1.0, 1.0) point of perfect convergence, indicating a sharp fall in the BER curve after  $I_{\text{outer}} = 9$  outer iterations, which is confirmed by the decoding trajectory for  $N_w = 6$ . The EXIT charts with  $N_w = 2, 4$  are, however, “pinched off”; thus, the BER at this SNR do not converge.

more, the multiple-symbol detection partially mitigates the inherent performance penalty imposed by noncoherent detection. We observe in Fig. 3 that, as  $N_w$  increases, the BER performance gradually approaches that of the perfect CSI-oriented coherent scheme. Note that all the performance characteristics exhibit a vanishingly low BER after  $I_{\text{outer}} = 9$  outer iterations, which is the explicit benefit of using the URC in the system. The URC is a low-complexity code, which has an infinite impulse response and hence assists the inner decoder in efficiently spreading the soft information [2], [19]. As a result, the extrinsic information transfer (EXIT) charts of Figs. 4 and 5 converge to the (1.0, 1.0) point of perfect decoding convergence, leading to a vanishingly low BER, thus eliminating the system’s error floor. The maximum achievable rates for the corresponding scheme, where the scheme still exhibits an infinitesimally low BER were computed by exploiting the area property of EXIT charts [18], [32], [33] and are shown in Fig. 3 as the ultimate benchmark of the scheme.

To elaborate further, Figs. 4 and 5 portray the EXIT charts of our proposed scheme at SNR = 1 and 4 dB, respectively. We observe in Fig. 4 that the inner decoder’s EXIT charts recorded at SNR = 1 dB for  $N_w = 2, 4$  are “pinched off”, i.e., there remains no “open” EXIT tunnel, indicating a high residual BER. By contrast, the BER associated with  $N_w = 6, 10$  may be expected to decrease sharply at this SNR after  $I_{\text{outer}} = 9$  outer iterations, which is confirmed by the staircase-shaped Monte Carlo-simulation-based decoding trajectory [8], [13]. Fig. 5, on the other hand, shows the EXIT charts at SNR = 4 dB, where all the curves associated with  $N_w = 2, 4, 6, 10$  exhibit an open EXIT tunnel, implying an infinitesimally low BER after  $I_{\text{outer}} = 9$  iterations. The EXIT charts of the soft-decision-MSDSD-aided OFDM DSTSK recorded both for SNR = 1 dB and SNR = 4 dB are further compared in Figs. 4 and 5 to the ultimate benchmark of the coherent detector assuming perfect CSI at the receiver.

Fig. 6 characterizes the complexity associated with the MAP-MSDSD of the OFDM-aided DSTSK (2, 2, 2, 4, 4) scheme at SNR = 4 dB as a function of the window size  $N_w$ , parameterized by the

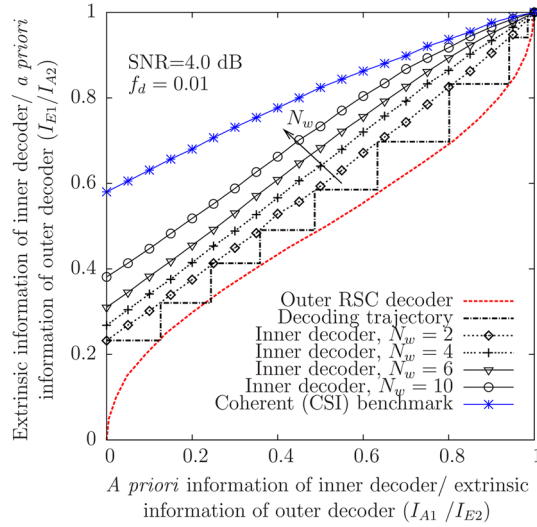


Fig. 5. EXIT charts of the inner decoders of the soft-decision-MSDSD-aided OFDM DSTSK (2, 2, 2, 4, 4) scheme at SNR = 4 dB with normalized Doppler frequency  $f_d = 0.01$  and different observation window sizes  $N_w = 2, 4, 6, 10$  and that of the corresponding coherent inner decoder as a benchmark. All the EXIT charts have a quite open EXIT tunnel at this SNR and converge to the (1.0, 1.0) point as a benefit of employing the URC, indicating a sharp fall in the BER curve after  $I_{\text{outer}} = 9$  outer iterations, which is confirmed by the decoding trajectory.

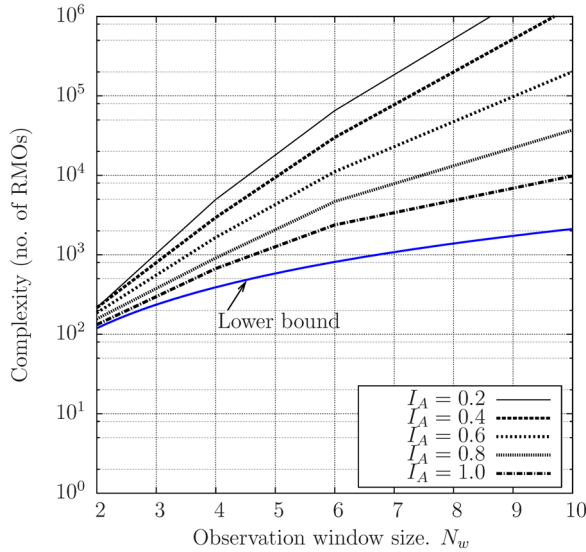


Fig. 6. Complexity in terms of the numbers of RMOs for the proposed DSTSK (2, 2, 2, 4, 4) scheme at SNR = 4 dB using the parameters of Table I as a function of observation window size  $N_w$  parameterized against the  $a$  priori mutual information of the inner decoder  $I_A$ . The complexity shoots up with  $N_w > 6$ , although the rate of increase in complexity slows down with increased  $a$  priori information.

358  $a$  priori information  $I_A$  provided by the outer decoder for the demap-  
359 per of Fig. 1. The  $a$  priori information  $I_A$  is measured by the average  
360 mutual information [13] between the  $a$  priori LLR  $L_a(b_j)$  and the  
361  $a$  posteriori LLR  $L_p(b_j)$  of Fig. 1. The influence of the  $a$  priori  
362 information  $I_A$  on the complexity may be beneficially exploited in the  
363 context of adaptive system design [15], where  $N_w$  may be adaptively  
364 selected depending on the quality of the soft input. To be specific,  
365 the value of  $I_A$  increases during the consecutive decoding iterations,  
366 and we can flexibly increase  $N_w$  when the value of  $I_A$  is higher. The  
367 theoretical lower bound of the complexity quantified by the number  
368 of RMOs in Section IV is also shown as a benchmarker in Fig. 6.  
369 As expected, the complexity rapidly escalates upon increasing  $N_w$ ,

albeit it does not become excessively high, provided that the  $a$  priori 370  
information gleaned from the outer decoder is in the range of  $I_A \geq 0.8$ . 371

## VI. CONCLUSION

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We have proposed a soft-decision-MSDSD-aided multicarrier 373  
DSTSK scheme for communications over wideband channels. The 374  
OFDM-aided DSTSK provides a flexible diversity versus multiplexing 375  
gain tradeoff by spreading the source information across both the spa- 376  
tial and time dimensions, while mitigating the potential performance 377  
degradation imposed by the frequency selectivity of the channel. The 378  
turbo-principle-based soft-decision MSDSD facilitates joint decisions 379  
over a number of DSTSK codewords, while exploiting the fading- 380  
plus-noise statistics of the channel. We have demonstrated that the 381  
proposed soft-decision-MSDSD-aided DSTSK scheme provides sub- 382  
stantial flexibility at moderate complexity owing to dispensing with 383  
CE. Furthermore, the MSDSD mitigates the performance degradation 384  
inflicted by the CDD scheme without an undue increase in computa- 385  
tional complexity. 386

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## AUTHOR QUERY

NO QUERY.



# Correspondence

## OFDM-Aided Differential Space–Time Shift Keying Using Iterative Soft Multiple-Symbol Differential Sphere Decoding

Mohammad Ismat Kadir, Sheng Chen, K. V. S. Hari,  
K. Giridhar, and Lajos Hanzo

**Abstract**—Soft-decision multiple-symbol differential sphere decoding (MSDSD) is proposed for orthogonal frequency-division multiplexing (OFDM)-aided differential space–time shift keying (DSTSK)-aided transmission over frequency-selective channels. Specifically, the DSTSK signaling blocks are generated by the channel-encoded source information and the space–time (ST) blocks are appropriately mapped to a number of OFDM subcarriers. After OFDM demodulation, the DSTSK signal is noncoherently detected by our soft-decision MSDSD detector. A novel soft-decision MSDSD detector is designed, and the associated decision rule is derived for the DSTSK scheme. Our simulation results demonstrate that an SNR reduction of 2 dB is achieved by the proposed scheme using an MSDSD window size of  $N_w = 4$  over the conventional soft-decision-aided differential detection benchmark, while communicating over dispersive channels and dispensing with channel estimation (CE).

**Index Terms**—Extrinsic information transfer (EXIT) chart, iterative decoding, multiple-symbol differential sphere decoding (MSDSD), orthogonal frequency-division multiplexing (OFDM), space–time shift keying (STSK).

### I. INTRODUCTION

Space–time shift keying (STSK) [1]–[3] has emerged as a beneficial multiple-input–multiple-output (MIMO) concept. STSK bridges the gap between the flexible diversity–multiplexing tradeoff provided by linear dispersion codes (LDCs) [4], [5] and the low-complexity design of spatial modulation (SM) [6]. Similar to the LDCs, STSK spreads the user information to both the spatial and time dimensions, but instead of simultaneously activating all the dispersion matrices (DMs), it transmits an additional  $\log_2 Q$  bits by activating one out of  $Q$  DMs. To overcome the performance degradation of the STSK scheme in wideband channels, orthogonal frequency-division multiplexing (OFDM)-aided STSK [7] and orthogonal frequency-division multiple-access/single-carrier frequency-division multiple-access-aided STSK [8] have also been proposed. The previous STSK studies [1], [2] demonstrate that coherent STSK performs well in conjunction with perfect channel state information (CSI) but exhibits a severe error floor in the presence of channel estimation (CE) errors.

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M. I. Kadir, S. Chen, and L. Hanzo are with the School of Electronics and Computer Science, University of Southampton, Southampton SO17 1BJ, U.K. (e-mail: mik1g09@ecs.soton.ac.uk; sqc@ecs.soton.ac.uk; lh@ecs.soton.ac.uk).

K. V. S. Hari is with the Department of Electrical Communication Engineering, Indian Institute of Science, Bangalore 560012, India (e-mail: hari@ece.iisc.ernet.in).

K. Giridhar is with the Department of Electrical Engineering, Indian Institute of Technology Madras, Chennai 600036, India (e-mail: giri@tenet.res.in).

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Differential STSK (DSTSK) employing conventional differential detection (CDD) has also been proposed for the sake of dispensing with the CE [2] and thus to eliminate the potentially high-Doppler-dependent pilot overhead. However, CDD suffers from a typical 3-dB performance penalty in low-Doppler scenarios. Furthermore, an irreducible error floor may be observed in a high-mobility scenario characterized by a high Doppler frequency. To circumvent the performance degradation of CDD, multiple-symbol differential detection (MSDD) was proposed for differential phase-shift keying (DPSK) in [9]. MSDD uses the fading-plus-noise statistics of the channel for jointly detecting  $(N_w - 1)$  information symbols from  $N_w$  number of consecutively received symbols, where  $N_w$  is usually referred to as the *observation window size*. The performance improvement of MSDD is, however, achieved at the cost of increased complexity, which increases exponentially with  $N_w$ . For mitigating this potentially excessive complexity, sphere decoding (SD) was invoked for MSDD in the context of multiple-symbol differential sphere decoding (MSDSD) in [10] and [11]. Hard-decision MSDSD was conceived in [12] for a DSTSK scheme operating in nondispersive channels. As a further advance, inspired by the near-capacity performance of turbo detection [13], [14], a soft-decision MSDSD scheme was also designed for DPSK in [15]. Furthermore, the concept of differential space–frequency modulation employing MSDSD in conjunction with a specific subcarrier allocation was proposed in [16] for exploiting both the achievable spatial- and frequency-domain diversity. However, the conception of the soft-decision-MSDSD-aided DSTSK designed for realistic dispersive scenarios constitutes an unexplored open problem.

Against this background, we conceive a novel soft-decision MSDSD for OFDM-based DSTSK operating in frequency-selective channels. The main contributions of this paper are as follows.

- 1) A novel soft-decision-aided MSDSD is proposed for OFDM-aided DSTSK operating in dispersive channels. The decision rule of the soft-decision MSDSD is deduced by considering the construction of DSTSK codewords based on the DMs, the Doppler frequency, the OFDM system parameters, and the generation of soft information.
- 2) A lower bound of the detection complexity is deduced, which is verified by simulations.

The remainder of this paper is organized as follows. In Section II, an overview of the proposed channel-coded OFDM-aided DSTSK scheme is provided. The soft-decision MSDSD is modeled in Section III. In Section IV, both the complexity imposed by the system is quantified. The performance of the soft-decision MSDSD-aided DSTSK scheme is investigated in Section V. Finally, we conclude in Section VI.

**Notations:** We use capital boldface letters to denote matrices, whereas  $\{\cdot\}^T$ ,  $\{\cdot\}^H$ ,  $\text{tr}(\cdot)$ ,  $\det[\cdot]$ , and  $\|\cdot\|$  are used to represent the transpose, the Hermitian transpose, the trace, the determinant, and the Euclidean norm of the matrix “ $\cdot$ ,” respectively. The notations  $\mathcal{E}\{\cdot\}$ ,  $\cdot^*$ , and  $P(\cdot)$  are used to denote the expected value, the complex conjugate, and the probability of “ $\cdot$ ” respectively, whereas  $\otimes$  and  $I_T$  represent the Kronecker product and the  $(T \times T)$  identity matrix, respectively. A symmetric  $(N_w \times N_w)$  Toeplitz matrix is denoted  $\text{toeplitz}\{x_1, \dots, x_{N_w}\}$ , whereas  $\text{diag}\{\mathbf{X}_1, \dots, \mathbf{X}_{N_w}\}$  indicates a block-diagonal matrix with the matrices  $\mathbf{X}_1, \dots, \mathbf{X}_{N_w}$  on its main diagonal. Furthermore,  $\mathcal{CN}(\mu, \sigma^2)$  refers to the circularly symmetric complex Gaussian distribution with mean  $\mu$  and variance  $\sigma^2$ .

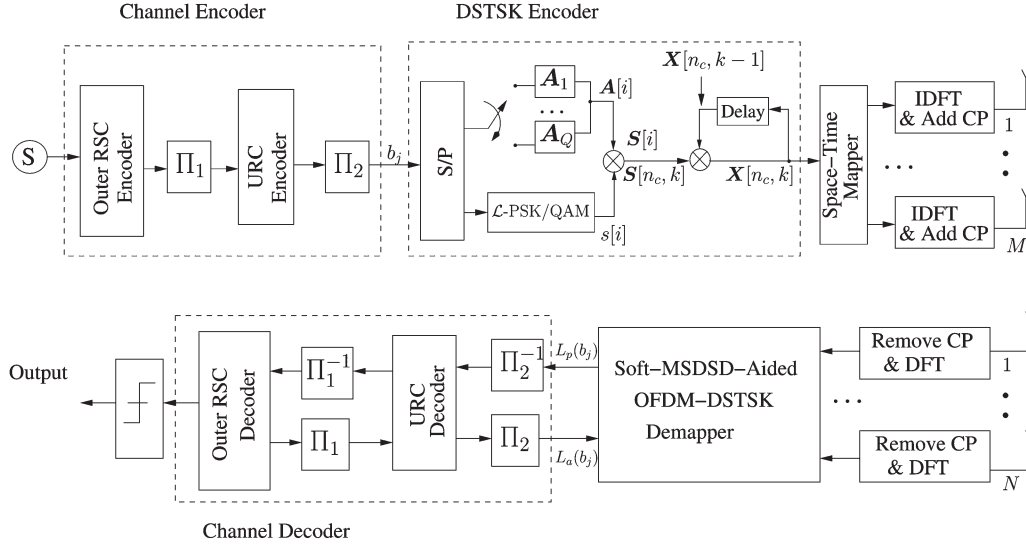


Fig. 1. Transceiver architecture of the proposed concatenated channel-coding-aided DSTSK scheme relying on the soft-decision MSDSD as the inner decoder.

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## II. SYSTEM MODEL

We consider a channel-coded OFDM-aided DSTSK transceiver employing  $M$  transmit and  $N$  receive antenna elements (AEs), as shown in Fig. 1. The channel encoder/decoder blocks of Fig. 1 may incorporate a general channel coding scheme that supports soft-decision decoding at affordable complexity. A pragmatic coding architecture might be an appropriately interleaved serially concatenated recursive convolutional code (RSC) and unity rate code (URC)-aided scheme [17]–[19], as shown in Fig. 1.

The source bits are first channel encoded by the RSC code, and the encoded bits are then interleaved by a random bit interleaver  $\Pi_1$ . Following URC precoding, the interleaved bits are further interleaved by a second interleaver  $\Pi_2$ . The resultant bits are then mapped to STSK codewords, which are further mapped to  $N_c$  parallel subcarriers and then differentially encoded in the time domain (TD), i.e., across the consecutive OFDM symbols of the same subcarrier. The DSTSK codewords are then OFDM modulated, while incorporating appropriate cyclic prefixes (CPs).

The signal received is first OFDM demodulated and then input to the DSTSK soft-decision MSDSD demapper. The extrinsic soft information is then iteratively exchanged between the three soft-in-soft-out components, namely, the DSTSK demapper, the URC decoder, and the RSC decoder, before finally outputting the estimated bits [8], [19].

### A. DSTSK Architecture and OFDM Layout

The STSK encoder generates space-time (ST) codewords from the source information by activating a single DM in any symbol duration in conjunction with the classic modulated symbols for transmission over  $T$  time slots using  $M$  transmit AEs [1], [2]. More specifically, each STSK signaling block  $\mathbf{S}[i] \in \mathbb{C}^{T \times M}$  is created from  $\log_2(\mathcal{L} \cdot Q)$  source bits according to [1], [2]

$$\mathbf{S}[i] = s[i] \mathbf{A}[i] \quad (1)$$

where  $s[i]$  is an  $\mathcal{L}$ -ary constellation symbol represented by  $\log_2 \mathcal{L}$  bits, and  $\mathbf{A}[i] \in \mathbb{C}^{T \times M}$  is the specific DM activated from the set of  $Q$  DMs  $\mathbf{A}_q (q = 1, \dots, Q)$ , as determined by the remaining  $\log_2 Q$  bits. The DMs  $\mathbf{A}_q (q = 1, \dots, Q)$  are unitary matrices generated by employing an exhaustive search for minimizing the objective function constituted by the pairwise error probability of the codewords [5], [12], [20], [21] under the power constraint in [2] expressed by  $\text{tr}(\mathbf{A}_q^H \mathbf{A}_q) =$

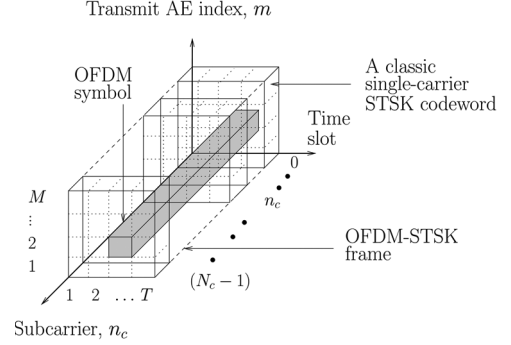


Fig. 2. Mapping of the STSK codewords to  $N_c$  parallel OFDM subcarriers showing the construction of an OFDM-STSK frame and OFDM symbols. After being appropriately mapped to the subcarriers, the codewords are differentially encoded in the TD and transmitted over dispersive channels by  $M$  transmit AEs over  $T$  time-slots.

$T \forall q$ . The resultant STSK system is then uniquely and unambiguously described by the parameters  $(M, N, T, Q, \mathcal{L})$ .

We observe that the STSK codeword  $\mathbf{S}[i]$  belongs to a set  $\mathcal{S}$  of  $(\mathcal{L} \cdot Q)$  codeword matrices defined by

$$\mathcal{S} \triangleq \{s_l \mathbf{A}_q | (q \in \{1, \dots, Q\}, l \in \{1, \dots, \mathcal{L}\})\}. \quad (2)$$

The STSK codewords are mapped to  $N_c$  parallel subcarriers, as shown in Fig. 2, before being differentially encoded. As shown in Fig. 2,  $N_c$  consecutive STSK codewords are arranged in parallel to form an OFDM-STSK frame, and OFDM modulation is carried out over each shaded symbol pipe, which constitutes an OFDM symbol. We may represent the codeword  $\mathbf{S}[i]$  by  $\mathbf{S}[n_c, k]$ , so that the overall codeword index  $i$  is related to the OFDM frame index  $k$  and the subcarrier index  $n_c$  by  $i = kN_c + n_c, n_c = 0, 1, \dots, (N_c - 1)$ . Additionally, we invoke differential encoding in the TD, i.e., differential encoding is performed across the consecutive OFDM symbols of the same subcarrier. We have chosen TD differential encoding/decoding because we have conceived our scheme for continuous transmissions as opposed to the FD differential encoding/decoding across adjacent FD subcarriers, which is more suitable for burst transmissions. To facilitate convenient differential encoding, we assume  $M = T$ . Furthermore, directly generated unitary DMs are used in the proposed scheme for avoiding the nonlinear Cayley transform [2], [12]. The codewords

156  $\mathbf{S}[n_c, k]$  are thus differentially encoded to form the transmit blocks  
 157  $\mathbf{X}[n_c, k]$  ( $k = 0, 1, 2, \dots$ ) according to [16]

$$\mathbf{X}[n_c, k] = \begin{cases} \mathbf{X}[n_c, k-1]\mathbf{S}[n_c, k], & k = 1, 2, \dots \\ \mathbf{I}_T & k = 0. \end{cases} \quad (3)$$

158 The DSTSK codewords are then transmitted after the  $N_c$ -point inverse  
 159 discrete Fourier transform (DFT) operation and appropriate CP incor-  
 160 poration.

#### 161 B. Channel Model

162 Each link between the  $m$ th transmit and  $n$ th receive AE is assumed  
 163 a frequency-selective channel, but as a benefit of OFDM-based trans-  
 164 mission, each dispersive channel is then partitioned into  $N_c$  low-rate  
 165 parallel frequency-flat subchannels [22]. The complex-valued fading  
 166 gain  $h_{m,n}[n_c, k]$  ( $m = 1, 2, \dots, M$ ;  $n = 1, 2, \dots, N$ ) obeys the dis-  
 167 tribution  $\mathcal{CN}(0, 1)$  associated with an autocorrelation function based  
 168 on Clarke's model [23]:  $\varphi_{hh}[n_c, \kappa] \triangleq \mathcal{E}\{h_{m,n}[n_c, k]h_{m,n}^*[n_c, k +$   
 169  $\kappa]\} = J_0(2\pi\kappa f_d)$ , where  $J_0$  denotes the zeroth-order Bessel function  
 170 of the first kind, and  $f_d = f_m T$  is the normalized maximum Doppler  
 171 frequency, whereas  $f_m$  and  $1/T$  represent the maximum Doppler  
 172 frequency and the symbol rate, respectively. The fading is assumed  
 173 quasi-static, i.e., the channel's complex-valued envelope remains ap-  
 174 proximately constant during the transmission of an OFDM STSK  
 175 frame.

176 Given the aforementioned assumptions, the received signal  
 177  $\mathbf{Y}[n_c, k] \in \mathbb{C}^{T \times N}$  obtained after CP removal and DFT may be ex-  
 178 pressed by [16], [22]

$$\mathbf{Y}[n_c, k] = \mathbf{X}[n_c, k]\mathbf{H}[n_c, k] + \mathbf{V}[n_c, k] \quad (4)$$

179 where  $\mathbf{X}[n_c, k] \in \mathbb{C}^{T \times M}$  represents the codeword transmitted and  
 180  $\mathbf{H}[n_c, k] \in \mathbb{C}^{M \times N}$  denotes the FD channel transfer matrix, with  
 181 its  $(m, n)$ th entry given by  $h_{m,n}[n_c, k]$ . Furthermore,  $\mathbf{V}[n_c, k] \in$   
 182  $\mathbb{C}^{T \times N}$  is the additive white Gaussian noise (AWGN) with entries of  
 183  $v_{T_i, n}[n_c, k] \sim \mathcal{CN}(0, \sigma_n^2)$ .

### 184 III. MSDSD RECEIVER

185 This section introduces the maximum-likelihood MSDSD (ML-  
 186 MSDSD), the maximum *a posteriori* MSDSD (MAP-MSDSD) algo-  
 187 rithm, and the generation of the log-likelihood ratios (LLRs) for the  
 188 soft-decision-MSDSD-aided OFDM DSTSK.

#### 189 A. ML-MSDSD for OFDM-Aided DSTSK

190 The ML-MSDD processes  $N_w$  consecutively received space-time  
 191 blocks corresponding to the  $n_c$ th subcarrier given by  $\bar{\mathbf{Y}}[n_c, k] \triangleq$   
 192  $[\mathbf{Y}^T[n_c, k - N_w + 1], \dots, \mathbf{Y}^T[n_c, k]]^T$  and finds the ML estimates  
 193  $\hat{\bar{\mathbf{X}}}[n_c, k]$  of the corresponding  $N_w$  transmitted blocks  $\bar{\mathbf{X}}[n_c, k] \triangleq$   
 194  $[\mathbf{X}^T[n_c, k - N_w + 1], \dots, \mathbf{X}^T[n_c, k]]^T$  [24]. Since the differen-  
 195 tially encoded blocks  $\mathbf{X}[n_c, k]$  are related to the STSK codewords  
 196 by the one-to-one relationship expressed by (3), the ML-MSDD  
 197 in turn estimates  $(N_w - 1)$  STSK codewords given by  $\bar{\mathbf{S}}[n_c, k] \triangleq$   
 198  $[\mathbf{S}^T[n_c, k - N_w + 2], \dots, \mathbf{S}^T[n_c, k]]^T$ , which further estimates the  
 199 source bits mapped to the STSK codewords.

200 Defining a block-diagonal matrix  $\bar{\mathbf{X}}_D[n_c, k]$  by  $\bar{\mathbf{X}}_D[n_c, k] \triangleq$   
 201  $\text{diag}\{\mathbf{X}[n_c, k - N_w + 1], \dots, \mathbf{X}[n_c, k]\}$ ,  $\bar{\mathbf{H}}[n_c, k] \triangleq [\mathbf{H}^T[n_c, k -$   
 202  $N_w + 1], \dots, \mathbf{H}^T[n_c, k]]^T$ , and  $\bar{\mathbf{V}}[n_c, k] \triangleq [\mathbf{V}^T[n_c, k - N_w + 1],$   
 203  $\dots, \mathbf{V}^T[n_c, k]]^T$ , the  $N_w$ -block received sequence can be expressed  
 204 by [11], [24]

$$\bar{\mathbf{Y}}[n_c, k] = \bar{\mathbf{X}}_D[n_c, k]\bar{\mathbf{H}}[n_c, k] + \bar{\mathbf{V}}[n_c, k] \quad (5)$$

where

$$\begin{aligned} \bar{\mathbf{Y}}[n_c, k] &\in \mathbb{C}^{N_w M \times N} & \bar{\mathbf{V}}[n_c, k] &\in \mathbb{C}^{N_w M \times N} \\ \bar{\mathbf{H}}[n_c, k] &\in \mathbb{C}^{N_w M \times N} & \bar{\mathbf{X}}_D[n_c, k] &\in \mathbb{C}^{N_w M \times N_w M}. \end{aligned}$$

For the sake of notational simplicity, we omit the subcarrier index and  
 time index  $[n_c, k]$  in the following and refer to the  $\mu$ th submatrix of a  
 block matrix, e.g.,  $\mathbf{B}$  by the subscripted matrix  $\mathbf{B}_\mu$ . Under the assump-  
 tion that  $h_{m,n}$  and  $v_{T_i, n}$  ( $m = T_i = 1, 2, \dots, M$ ,  $n = 1, 2, \dots, N$ )  
 are zero-mean Gaussian random processes, the probability density  
 function of  $\bar{\mathbf{Y}}$  conditioned on  $\bar{\mathbf{X}}_D$  is given by [25], [26]

$$P(\bar{\mathbf{Y}}|\bar{\mathbf{X}}_D) = \frac{1}{(\pi^{N_w M} \det[\mathbf{\Lambda}_Y])^N} \exp \left\{ -\text{tr} \left( \bar{\mathbf{Y}}^H \mathbf{\Lambda}_Y^{-1} \bar{\mathbf{Y}} \right) \right\} \quad (6)$$

where  $\mathbf{\Lambda}_Y$  is defined by  $\mathbf{\Lambda}_Y \triangleq \mathcal{E}\{\bar{\mathbf{Y}}\bar{\mathbf{Y}}^H|\bar{\mathbf{X}}_D\}$ . The ML estimate  $\hat{\bar{\mathbf{X}}}$   
 under the assumption of quasi-static fading and unitary  $\bar{\mathbf{X}}_D$  reduces  
 to [26], [27]

$$\hat{\bar{\mathbf{X}}} = \arg \max_{\bar{\mathbf{X}}} P(\bar{\mathbf{Y}}|\bar{\mathbf{X}}_D) = \arg \min_{\bar{\mathbf{X}}} \left\{ \text{tr} \left( \bar{\mathbf{Y}}^H \mathbf{\Lambda}_Y^{-1} \bar{\mathbf{Y}} \right) \right\}. \quad (7)$$

Here, the conditional covariance matrix  $\mathbf{\Lambda}_Y$  is related to the channel  
 parameters [25], [26] by

$$\mathbf{\Lambda}_Y^{-1} = \frac{1}{N} \bar{\mathbf{X}}_D (\mathbf{\Lambda}^{-1} \otimes \mathbf{I}_M) \bar{\mathbf{X}}_D^H \quad (8)$$

where we have  $\mathbf{\Lambda} \triangleq (\psi_{hh} + \sigma_n^2 \mathbf{I}_{N_w})$  and  $\psi_{hh} \triangleq \text{toeplitz}\{\varphi_{hh}[n_c, 217$   
 $0], \dots, \varphi_{hh}[n_c, (N_w - 1)]\}$ , with the component autocorrelation  
 functions  $\varphi_{hh}[n_c, \kappa]$  being identical for all spatial channels. Applying  
 the Cholesky factorization of  $\mathbf{\Lambda}^{-1} = \mathbf{U}^H \mathbf{U}$  with the upper triangular  
 matrix  $\mathbf{U}$  and considering the identity  $\text{tr}(\mathbf{X}\mathbf{X}^H) = \|\mathbf{X}\|^2$  for any  
 matrix  $\mathbf{X}$ , the ML-MSDD decision rule can be deduced from (6),  
 yielding [11]

$$\hat{\bar{\mathbf{X}}} = \arg \min_{\bar{\mathbf{X}}} \left\{ \sum_{\mu=1}^{N_w} \left\| \mathbf{Y}_{\mu, \mu}^H \bar{\mathbf{X}}_\mu + \sum_{\nu=\mu+1}^{N_w} (\mathbf{Y}_{\mu, \nu}^H \bar{\mathbf{X}}_\nu) \right\|^2 \right\} \quad (9)$$

where  $\mathbf{Y}_{\mu, \nu}^H$  is defined by  $\mathbf{Y}_{\mu, \nu}^H \triangleq \mathbf{Y}_{\nu}^H u_{\mu, \nu}$ , and  $u_{\mu, \nu}$  represents the  
 $(\mu, \nu)$ th element of  $\mathbf{U}$ . Still referring to (9),  $\hat{\bar{\mathbf{X}}}$  denotes the estimate  
 of  $\bar{\mathbf{X}}_D$ , whereas  $\bar{\mathbf{X}}_\mu$  refers to the  $\mu$ th candidate submatrix of  $\bar{\mathbf{X}}_D$ .

Since the ML metric of (9) is invariant to a phase shift common to all  
 elements of  $\bar{\mathbf{X}}_\mu \forall \mu$  corresponding to the same  $\bar{\mathbf{S}}$  (where  $\mathbf{S}_\lambda^T \in \mathcal{S} \forall \lambda$ ),  
 the accumulated differential matrices may be expressed as [15]

$$\mathcal{A}_\nu = \begin{cases} \prod_{\lambda=\nu}^{N_w-1} \mathbf{S}_\lambda^H, & 1 \leq \nu \leq (N_w - 1) \\ \mathbf{I}_T, & \nu = N_w. \end{cases} \quad (10)$$

For the sake of reducing the complexity associated with an exhaustive  
 search, we employ MSDSD similar to [10] and [11] to search through  
 the candidate set lying within a sphere of radius  $\rho_s$  as follows:

$$\sum_{\mu=1}^{N_w} \left\| \mathbf{Y}_{\mu, \mu}^H \mathcal{A}_\mu + \sum_{\nu=\mu+1}^{N_w} (\mathbf{Y}_{\mu, \nu}^H \mathcal{A}_\nu) \right\|^2 \leq \rho_s^2. \quad (11)$$

#### 233 B. MAP-MSDSD

Assuming the STSK codewords to be mutually independent, (6) and  
 (9) yield [15]:

$$\begin{aligned} & -\ln(P(\bar{\mathbf{S}}|\bar{\mathbf{Y}})) \\ & \propto -\ln(P(\bar{\mathbf{Y}}|\bar{\mathbf{S}})) - \ln(P(\bar{\mathbf{S}})) \\ & \propto \sum_{\mu=1}^{N_w} \left\{ \left\| \mathbf{Y}_{\mu, \mu}^H \mathcal{A}_\mu + \sum_{\nu=\mu+1}^{N_w} \mathbf{Y}_{\mu, \nu}^H \mathcal{A}_\nu \right\|^2 - \ln(P(\mathbf{S}_\mu)) \right\}. \end{aligned} \quad (12)$$



236 The MAP-MSDSD may be thus expressed as

$$\sum_{\mu=1}^{(N_w-1)} \left( \left\| \sum_{\nu=\mu}^{N_w} (\mathbf{Y}_{\mu,\nu}^H \mathbf{A}_{\nu}) \right\|^2 - \ln(P(\mathbf{S}_{\mu})) \right) \leq \rho_s^2 - \|u_{N_w, N_w} \mathbf{Y}_{N_w}\|^2 \triangleq \rho^2. \quad (13)$$

237 Clearly, the codeword  $\mathbf{S}_{\mu}$  obeys the specific distance criterion [10],  
238 [12], [28] that the current partial Euclidean distance (PED)  $d_{\mu}^2$  is the  
239 sum of the previous PED  $d_{\mu+1}^2$  and the distance increment  $\Delta_{\mu}^2$ , i.e.,

$$\begin{aligned} d_{\mu}^2 &\triangleq \Delta_{\mu}^2 + d_{\mu+1}^2 \\ &= \left\| u_{\mu,\mu} \mathbf{Y}_{\mu} \mathbf{A}_{\mu+1} \mathbf{S}_{\mu}^H + \sum_{\nu=\mu+1}^{N_w} u_{\mu,\nu} \mathbf{Y}_{\nu} \mathbf{A}_{\nu} \right\|^2 - \ln(P(\mathbf{S}_{\mu})) \\ &\quad + \sum_{\iota=\mu+1}^{(N_w-1)} \left( \left\| \sum_{\nu=\iota}^{N_w} u_{\iota,\nu} \mathbf{Y}_{\nu} \mathbf{A}_{\nu} \right\|^2 - \ln(P(\mathbf{S}_{\iota})) \right) \leq \rho^2. \quad (14) \end{aligned}$$

240 Similar to the MSDSD principle described in [10] and [15], the  
241 MAP-MSDSD is initialized with  $\mu = (N_w - 1)$  and then proceeds by  
242 applying the search criterion of (14) until  $\mu = 1$ , where the search ra-  
243 dius is updated to  $\rho^2 = d_1^2$ , and the search is repeated by commencing  
244 from  $\mu = 2$  until  $\mu = (N_w - 1)$  is reached. If the new search does not  
245 provide a better estimate, the previous estimate is retained.

#### 246 C. Log-Likelihood Ratio and Soft-Decision-MSDSD-Aided 247 OFDM DSTSK

248 The soft demapper relies on the *a priori* information gleaned from  
249 the URC decoder and the MAP-MSDSD. A high interleaver depth is  
250 assumed so that the permuted bits may be treated as being independent.  
251 The LLR corresponding to the bit  $b_j$  interleaved by the interleaver  
252  $\Pi_2$  of Fig. 1 is defined by [29]  $L_a(b_j) \triangleq \ln(P(b_j = b)/P(b_j = \bar{b}))$ ,  
253 where  $b \in \{0, 1\}$ , and the  $j$ th bit  $b_j = b$  corresponds to the MAP-  
254 MSDSD estimate  $\hat{\mathbf{S}}$ , whereas  $\bar{b}$  indicates its complement. The  
255 *a posteriori* LLR  $L_p(\cdot)$  of  $b_j$  may be then approximated by the  
256 maximum-logarithmic-MAP (max-log-MAP) algorithm [14], [15], i.e.,

$$\begin{aligned} L_p(b_j) &= \ln \frac{P(b_j = b|\mathbf{Y})}{P(b_j = \bar{b}|\mathbf{Y})} \\ &\approx \ln \frac{\max_{\hat{\mathbf{S}}: b_j \neq \bar{b}} \left[ -\sum_{\mu=1}^{N_w} \left\| \sum_{\nu=\mu}^{N_w} (\mathbf{Y}_{\mu,\nu}^H \hat{\mathbf{A}}_{\nu}) \right\|^2 + \ln(P(\hat{\mathbf{S}})) \right]}{\max_{\hat{\mathbf{S}}: b_j = \bar{b}} \left[ -\sum_{\mu=1}^{N_w} \left\| \sum_{\nu=\mu}^{N_w} (\mathbf{Y}_{\mu,\nu}^H \hat{\mathbf{A}}_{\nu}) \right\|^2 + \ln(P(\hat{\mathbf{S}})) \right]} \\ &= -\sum_{\mu=1}^{N_w} \left\| \sum_{\nu=\mu}^{N_w} (\mathbf{Y}_{\mu,\nu}^H \hat{\mathbf{A}}_{\nu}) \right\|^2 + \ln(P(\hat{\mathbf{S}}^b)) \\ &\quad + \sum_{\mu=1}^{N_w} \left\| \sum_{\nu=\mu}^{N_w} (\mathbf{Y}_{\mu,\nu}^H \hat{\mathbf{A}}_{\nu}^{\bar{b}}) \right\|^2 - \ln(P(\hat{\mathbf{S}}^{\bar{b}})) \quad (15) \end{aligned}$$

257 where  $\hat{\mathbf{S}}^b$  and  $\hat{\mathbf{S}}^{\bar{b}}$  represent the MAP-MSDSD estimate and the  
258 constrained estimate associated with  $b_j = b$ , respectively.

The extrinsic LLR  $L_e(\cdot)$  for  $b_j$  is now evaluated by combining the  
*a posteriori* and *a priori* LLR:  $L_e(b_j) = L_p(b_j) - L_a(b_j)$ . The ex-  
trinsic information extracted from the soft-decision MSDSD demapper  
is iteratively exchanged with the URC decoder of Fig. 1, which forms  
the *inner* iteration, whereas the exchange of extrinsic information  
between the URC decoder and the RSC decoder of Fig. 1 may be  
termed as the *outer* iteration. Note that, for each outer iteration be-  
tween the RSC decoder and the URC decoder, several inner iterations  
may be invoked between the URC and the soft-decision-MSDSD-  
aided DSTSK demapper [2], [19]. Finally, the RSC decoder generates  
*a posteriori* LLRs, from which the source bits are estimated.

#### IV. COMPLEXITY

Here, the complexity of the proposed scheme is detailed, and the  
complexity imposed by the MAP-MSDSD is quantified.

Equation (2) shows that there exists  $(\mathcal{L} \cdot Q)$  legitimate code-  
word matrices for each  $\log_2(\mathcal{L} \cdot Q)$  bits of source information. The  
exhaustive-search-based solution to (9) involves a search in a  $(\mathcal{L} \cdot Q)^{(N_w-1)}$   
element space of candidate matrices  $\hat{\mathbf{X}}$  corresponding to all possible  
choices of  $\hat{\mathbf{S}}$ . The ML-MSDSD associated with chosen sphere  
radius  $\rho$  imposes average complexity, which is lower bounded by [11]

$$C \geq \frac{(\mathcal{L} \cdot Q)^{N_w \zeta - 1} - (\mathcal{L} \cdot Q)}{(\mathcal{L} \cdot Q) - 1} \quad (16)$$

where

$$\zeta \triangleq \frac{\sigma_n^2(1+\epsilon)}{2(1+\sigma_n^2)} \quad \rho^2 = (1+\epsilon)NMN_w, \quad \epsilon > 0. \quad (17)$$

To quantify the complexity of the MAP-MSDSD scheme, we con-  
sider the number of real-valued multiplication operations (RMOs)  
required for obtaining a single soft-output value, which is used as  
our complexity metric. The lower bound of the complexity may be  
obtained if the number of RMOs required for computing the soft  
outputs corresponding to the first codeword estimate  $\hat{\mathbf{S}}^b$  is counted  
and if a single constrained estimate  $\hat{\mathbf{S}}^{\bar{b}}$  is taken into account [15].  
Considering the upper diagonal nature of the matrix  $\mathbf{U}$ , we observe  
that  $\mathbf{Y}_{\mu,\nu}^H$  is defined only for  $\nu \geq \mu$  in the context of (9), although each  
 $\mathbf{Y}_{\mu,\nu}^H$  is an  $(N \times T)$ -element matrix, where  $T = M$ . The computation  
of the  $\mathbf{Y}_{\mu,\nu}^H$  terms in (15) thus involves a total of  $2MN[1 + 2 + \dots + N_w] = MN N_w(N_w + 1)$  RMOs, assuming real-valued autocorrela-  
tion functions of  $\varphi_{hh}[n_c, \kappa]$ . To compute the *a posteriori* LLRs given  
by (15), the number of RMOs associated with the computation of each  
 $\|\sum_{\nu=\mu}^{N_w} (\mathbf{Y}_{\mu,\nu}^H \hat{\mathbf{A}}_{\nu}^b)\|^2$  is  $4M^2N(N_w - \mu + 1) + 2$ . The number of  
RMOs required for generating  $\log_2(\mathcal{L} \cdot Q)$  soft outputs corresponding  
to a single codeword estimate  $\hat{\mathbf{S}}^b$  is thus given by

$$\begin{aligned} \text{RMO}[\hat{\mathbf{S}}^b] &= MN N_w(N_w + 1) + \sum_{\mu=1}^{N_w} [4M^2N(N_w - \mu + 1) + 2] \\ &= MN(2M + 1)N_w(N_w + 1) + 2N_w. \quad (18) \end{aligned}$$

On the other hand, the number of RMOs related to each bit  $\bar{b}$  of  
the constrained estimate  $\hat{\mathbf{S}}^{\bar{b}}$  is found to be  $2N_w[M^2N(N_w + 1) + 1]$ .  
The lower bound for the number of RMOs associated with the gen-  
eration of a single soft output is thus  $[MN(2M + 1)/\log_2(\mathcal{L} \cdot Q) + 300]$   
 $2M^2N]N_w^2$  for large  $N_w$ . The complexity of the scheme, however,  
depends on a number of parameters, such as on the channel SNR, on  
the autocorrelation function of the channel's fading plus noise, and  
most importantly, on the *a priori* mutual information  $I_A$  of the inner  
decoder [15]. In Section V, the complexity of the MAP-MSDSD will  
be investigated as a function of the observation window width  $N_w$   
parameterized by the available *a priori* information  $I_A$ .



TABLE I  
ADOPTED VALUES OF MAIN SIMULATION PARAMETERS

Parameter	Value
Dispersive channel model	COST207-TU12
Fast fading envelope	Correlated Rayleigh fading
Normalized Doppler spread, $f_d$	0.01
Number of subcarriers, $N_c$	128
Length of cyclic prefix	32
Overall symbol duration	300 ns
STSK ( $M, N, T, Q, L$ )	(2, 2, 2, 4, 4)
RSC encoder and decoder	Half rate
	Constraint length=2
Generator polynomial	(011, 010) <sub>2</sub>
Length of interleavers	200,000 bits
Outer iterations	11
Inner iterations	2

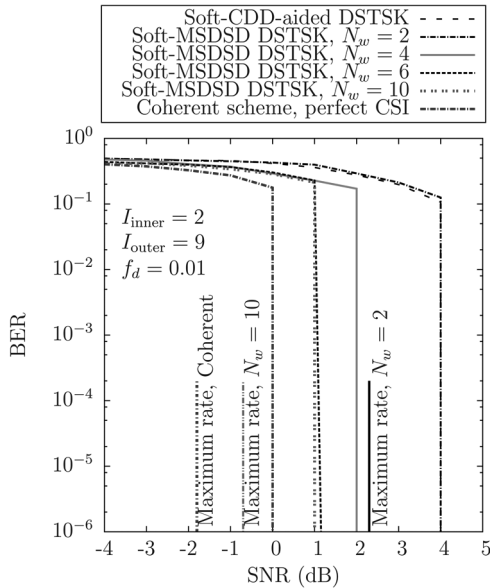


Fig. 3. Simulated BER performance of the soft-decision-MSDSD-aided OFDM DSTSK (2, 2, 2, 4, 4) scheme for transmission over dispersive COST207-TU12 channel with normalized Doppler frequency  $f_d = 0.01$  and different observation window size  $N_w = 2, 4, 6, 10$ . The BER falls sharply after  $I_{\text{outer}} = 9$  outer iterations as a benefit of employing the URC and the performance approaches that of the coherent scheme with perfect CSI with an increasing value of  $N_w$ .

## V. PERFORMANCE RESULTS

Here, the performance of the proposed scheme is investigated using the parameters listed in Table I. We have employed the COST207-TU12 channel model for the links between each transmit-receive antenna pair. The power delay profile characterizing the 12 taps of the COST207-TU12 channel is detailed in [30] and [31]. As mentioned in Table I, we employ an RSC (2, 1, 2) outer code having octally represented generator polynomials of  $(g_r, g) = (3, 2)_8$  as well as two random interleavers with a length of 200 000 bits.

Fig. 3 characterizes the achievable bit error rate (BER) of the proposed soft-decision-MSDSD-aided OFDM DSTSK scheme associated with observation window sizes of  $N_w = 2, 4, 6, 10$  and compares it to that achievable by the corresponding coherent scheme relying on perfect CSI. We observe that the proposed scheme has the benefit of dispensing with CE due to differential encoding, while mitigating the performance erosion of classic STSK by employing OFDM. Further-

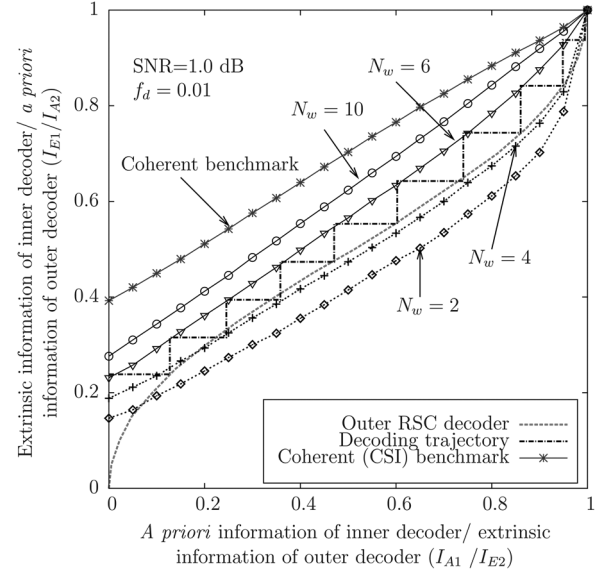


Fig. 4. EXIT charts of the inner decoders of the soft-decision-MSDSD-aided OFDM DSTSK (2, 2, 2, 4, 4) scheme at SNR = 1 dB with normalized Doppler frequency  $f_d = 0.01$  and different observation window sizes  $N_w = 2, 4, 6, 10$  and of the corresponding coherent system having perfect CSI at the receiver. At this SNR, we observe the inner EXIT charts with  $N_w = 6, 10$  have an open EXIT tunnel and converge to the (1.0, 1.0) point of perfect convergence, indicating a sharp fall in the BER curve after  $I_{\text{outer}} = 9$  outer iterations, which is confirmed by the decoding trajectory for  $N_w = 6$ . The EXIT charts with  $N_w = 2, 4$  are, however, “pinched off”; thus, the BER at this SNR do not converge.

more, the multiple-symbol detection partially mitigates the inherent performance penalty imposed by noncoherent detection. We observe in Fig. 3 that, as  $N_w$  increases, the BER performance gradually approaches that of the perfect CSI-oriented coherent scheme. Note that all the performance characteristics exhibit a vanishingly low BER after  $I_{\text{outer}} = 9$  outer iterations, which is the explicit benefit of using the URC in the system. The URC is a low-complexity code, which has an infinite impulse response and hence assists the inner decoder in efficiently spreading the soft information [2], [19]. As a result, the extrinsic information transfer (EXIT) charts of Figs. 4 and 5 converge to the (1.0, 1.0) point of perfect decoding convergence, leading to a vanishingly low BER, thus eliminating the system's error floor. The maximum achievable rates for the corresponding scheme, where the scheme still exhibits an infinitesimally low BER were computed by exploiting the area property of EXIT charts [18], [32], [33] and are shown in Fig. 3 as the ultimate benchmark of the scheme.

To elaborate further, Figs. 4 and 5 portray the EXIT charts of our proposed scheme at SNR = 1 and 4 dB, respectively. We observe in Fig. 4 that the inner decoder's EXIT charts recorded at SNR = 1 dB for  $N_w = 2, 4$  are “pinched off”, i.e., there remains no “open” EXIT tunnel, indicating a high residual BER. By contrast, the BER associated with  $N_w = 6, 10$  may be expected to decrease sharply at this SNR after  $I_{\text{outer}} = 9$  outer iterations, which is confirmed by the staircase-shaped Monte Carlo-simulation-based decoding trajectory [8], [13]. Fig. 5, on the other hand, shows the EXIT charts at SNR = 4 dB, where all the curves associated with  $N_w = 2, 4, 6, 10$  exhibit an open EXIT tunnel, implying an infinitesimally low BER after  $I_{\text{outer}} = 9$  iterations. The EXIT charts of the soft-decision-MSDSD-aided OFDM DSTSK recorded both for SNR = 1 dB and SNR = 4 dB are further compared in Figs. 4 and 5 to the ultimate benchmark of the coherent detector assuming perfect CSI at the receiver.

Fig. 6 characterizes the complexity associated with the MAP-MSDSD of the OFDM-aided DSTSK (2, 2, 2, 4, 4) scheme at SNR = 4 dB as a function of the window size  $N_w$ , parameterized by the

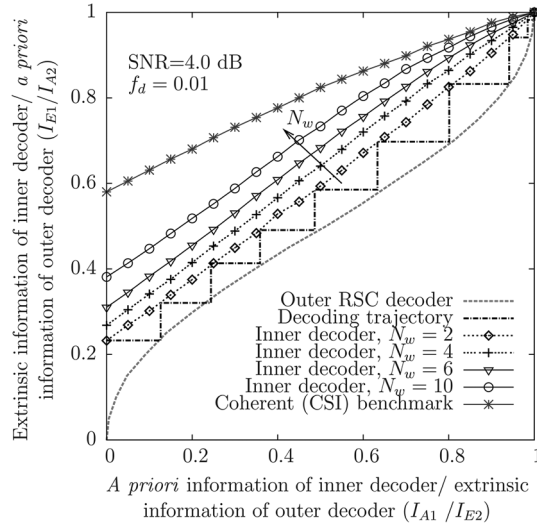


Fig. 5. EXIT charts of the inner decoders of the soft-decision-MSDSD-aided OFDM DSTSK (2, 2, 2, 4, 4) scheme at SNR = 4 dB with normalized Doppler frequency  $f_d = 0.01$  and different observation window sizes  $N_w = 2, 4, 6, 10$  and that of the corresponding coherent inner decoder as a benchmark. All the EXIT charts have a quite open EXIT tunnel at this SNR and converge to the (1.0, 1.0) point as a benefit of employing the URC, indicating a sharp fall in the BER curve after  $I_{\text{outer}} = 9$  outer iterations, which is confirmed by the decoding trajectory.

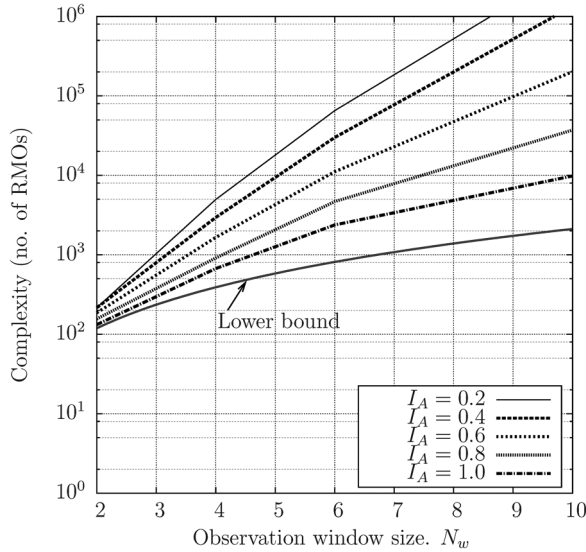


Fig. 6. Complexity in terms of the numbers of RMOs for the proposed DSTSK (2, 2, 2, 4, 4) scheme at SNR = 4 dB using the parameters of Table I as a function of observation window size  $N_w$  parameterized against the *a priori* mutual information of the inner decoder  $I_A$ . The complexity shoots up with  $N_w > 6$ , although the rate of increase in complexity slows down with increased *a priori* information.

358 *a priori* information  $I_A$  provided by the outer decoder for the demap-  
359 per of Fig. 1. The *a priori* information  $I_A$  is measured by the average  
360 mutual information [13] between the *a priori* LLR  $L_a(b_j)$  and the  
361 *a posteriori* LLR  $L_p(b_j)$  of Fig. 1. The influence of the *a priori*  
362 information  $I_A$  on the complexity may be beneficially exploited in the  
363 context of adaptive system design [15], where  $N_w$  may be adaptively  
364 selected depending on the quality of the soft input. To be specific,  
365 the value of  $I_A$  increases during the consecutive decoding iterations,  
366 and we can flexibly increase  $N_w$  when the value of  $I_A$  is higher. The  
367 theoretical lower bound of the complexity quantified by the number  
368 of RMOs in Section IV is also shown as a benchmarker in Fig. 6.  
369 As expected, the complexity rapidly escalates upon increasing  $N_w$ ,

albeit it does not become excessively high, provided that the *a priori* 370  
information gleaned from the outer decoder is in the range of  $I_A \geq 0.8$ . 371

## VI. CONCLUSION

372

We have proposed a soft-decision-MSDSD-aided multicarrier 373  
DSTSK scheme for communications over wideband channels. The 374  
OFDM-aided DSTSK provides a flexible diversity versus multiplexing 375  
gain tradeoff by spreading the source information across both the spa- 376  
tial and time dimensions, while mitigating the potential performance 377  
degradation imposed by the frequency selectivity of the channel. The 378  
turbo-principle-based soft-decision MSDSD facilitates joint decisions 379  
over a number of DSTSK codewords, while exploiting the fading- 380  
plus-noise statistics of the channel. We have demonstrated that the 381  
proposed soft-decision-MSDSD-aided DSTSK scheme provides sub- 382  
stantial flexibility at moderate complexity owing to dispensing with 383  
CE. Furthermore, the MSDSD mitigates the performance degradation 384  
inflicted by the CDD scheme without an undue increase in computa- 385  
tional complexity. 386

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