Error Probability and Capacity Analysis of Generalised Pre-Coding Aided Spatial Modulation

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Abstract—The recently proposed multiple input multiple output (MIMO) transmission scheme termed as generalized pre-coding aided spatial modulation (GPSM) is analyzed, where the key idea is that a particular subset of receive antennas is activated and the specific activation pattern itself conveys useful implicit information. We provide the upper bound of both the symbol error ratio (SER) and bit error ratio (BER) expression of the GPSM scheme of a low-complexity decoupled detector. Furthermore, the corresponding discrete-input continuous-output memoryless channel (DCMC) capacity as well as the achievable rate is quantified. Our analytical SER and BER upper bound expressions are confirmed to be tight by our numerical results. We also show that our GPSM scheme constitutes a flexible MIMO arrangement and there is always a beneficial configuration for our GPSM scheme that offers the same bandwidth efficiency as that of its conventional MIMO counterpart at a lower signal to noise ratio (SNR) per bit.

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I. INTRODUCTION

MULTIPLE INPUT MULTIPLE OUTPUT (MIMO) systems constitute one of the most promising recent technical advances in wireless communications, since they facilitate high-throughput transmissions in the context of various standards [1]. Hence, they attracted substantial research interests, leading to the Vertical-Bell Laboratories Layered Space-Time (V-BLAST) scheme [2] and to the classic Space Time Block Coding (STBC) arrangement [3]. The point-to-point single-user MIMO systems are capable of offering diverse transmission functionalities in terms of multiplexing-diversity-and beam-forming gains. Similarly, Spatial Division Multiple Access (SDMA) employed in the uplink and multi-user MIMO techniques invoked in the downlink also constitute beneficial building blocks [4], [5]. The basic benefits of MIMOs have also been recently exploited in the context of the network MIMO concept [6], [7], for constructing large-scale MIMOs [8], [9] and for conceiving beneficial arrangements for interference-limited MIMO scenarios [10].

Despite having a plethora of studies on classic MIMO systems, their practical constraints, such as their I/Q imbalance, their transmitter and receiver complexity as well as the cost of their multiple Radio Frequency (RF) Power Amplifiers (PA) chains as well as their Digital-Analogue/Analogue-Digital (DA/AD) converters have received limited attention. To circumvent these problems, low complexity alternatives to conventional MIMO transmission schemes have also been proposed, such as the Antenna Selection (AS) [11], [12] and the Spatial Modulation (SM) [13], [14] philosophies. More specifically, SM and generalised SM [15] constitute novel MIMO techniques, which were conceived for providing a higher throughput than a single-antenna aided system, while maintaining both a lower complexity and a lower cost than the conventional 56 MIMOs, since they may rely on a reduced number of RF conversion chains. To elaborate a little further, SM conveys extra information by mapping log2(Nt) bits to the Transmit Antenna (TA) indices of the Nt TAs, in addition to the classic 60 modulation schemes, as detailed in [13].

By contrast, the family of Pre-coding aided Spatial Modulation (PSM) schemes is capable of conveying extra information by appropriately selecting the Receive Antenna (RA) indices, 64 as detailed in [16]. More explicitly, in PSM the indices of the 65 RA represent additional information in the spatial domain. As a specific counterpart of the original SM, PSM benefits from both a low cost and a low complexity at the receiver side, therefore it may be considered to be eminently suitable for downlink transmissions [16]. The further improved concept of 70 Generalised PSM (GPSM) was proposed in [17], where comprehensive performance comparisons were carried out between the GPSM scheme as well as the conventional MIMO scheme and the associated detection complexity issues were discussed. Furthermore, a range of practical issues were investigated, namely the detrimental effects of realistic imperfect Channel 76 State Information at the Transmitter (CSIT), followed by a 77 low-rank approximation invoked for large-dimensional MI-78 MOs. Finally, the main difference between our GPSM scheme and the classic SM is that the former requires downlink pre-processing and CSIT, although they may be considered as a dual counterpart of each other and may hence be used in a hybrid manner. Other efforts on robust PSM was reported 83 in [18].

As a further development, in this paper, we provide the theoretical analysis of the recently proposed GPSM scheme [17], which is not available in the literature. More explicitly, both the discrete-input continuous-output memoryless channel (DCMC) capacity as well as the achievable rate are characterized. Importantly, tight upper bounds of the symbol error ratio (SER) and bit error ratio (BER) expressions are derived, when a decoupled low-complexity detector is employed.

The rest of our paper is organised as follows. In Section II, we introduce the underlying concept as well as the detection methods of the GPSM scheme. This is followed by our analytical study in Section III, where both the DCMC capacity and the achievable rate are characterized. In Section IV, we provide the theoretical analysis of the recently proposed GPSM scheme. Finally, it is plausible that in Section V, we conclude in Section V.

II. SYSTEM MODEL

A. Conceptual Description

Consider a MIMO system equipped with $N_t$ TAs and $N_r$ RAs, where we assume $N_t \geq N_r$. In this MIMO set-up, a maximum of $N_r$ parallel data streams may be supported, conveying a total of $k_{RF} = N_r k_{mod}$ bits altogether, where $k_{mod} = \log_2(M)$ denotes the number of bits per symbol of a conventional $M$-ary PSK/QAM scheme and its alphabet is denoted by $A$. Transmitter Pre-Coding (TPC) relying on the TPC matrix of $P \in \mathbb{C}^{N_r \times N_r}$ may be used for pre-processing the source signal before its transmission upon exploiting the knowledge of the CSI.

In contrast to the above-mentioned classic multiplexing of $N_r$ data streams, in our GPSM scheme a total of $N_s < N_r$ RAs are activated so as to facilitate the simultaneous transmission of $N_a$ data streams, where the particular pattern of the $N_a$ RAs activated conveys extra information in form of so-called spatial symbols in addition to the information carried by the conventional modulated symbols. Hence, the number of bits in GPSM conveyed by a spatial symbol becomes $k_{ant} = \lfloor \log_2(\{|\mathcal{C}_i|\}) \rfloor$, where the set $\mathcal{C}_i$ contains all the combinations associated with choosing $N_a$ activated RAs out of $N_r$ RAs.

As a result, the total number of bits transmitted by the GPSM scheme is $k_{RF} = k_{ant} + N_a k_{mod}$. Finally, it is plausible that the conventional MIMO scheme obeys $N_s = N_r$. For assisting further discussions, we also let $C(k)$ and $C(k,i)$ denote the 6th RA activation pattern and the $i$th activated RA in the $k$th activation pattern, respectively.

B. GPSM Transmitter

More specifically, let $s_{m}^{k}$ be an explicit representation of a so-called super-symbol $s \in \mathbb{C}^{N_t \times 1}$, indicating that the RA pattern $k$ is activated and $N_a$ conventional modulated symbols $b_m \in \mathbb{A}$ are transmitted, where we have $b_m \in \mathbb{A}$ and $\mathbb{E}[|b_m|^2] = 1$, $\forall i \in [1, N_a]$. In other words, we have the relationship

$$s_{m}^{k} = \Omega_k b_m,$$

where $\Omega_k = I[;C(k)]$ is constituted by the specifically selected columns determined by $C(k)$ of an identity matrix of $I_{N_r}$. Following TPC, the resultant transmit signal $x \in \mathbb{C}^{N_r \times 1}$ may be written as

$$x = \sqrt{\beta/N_a} P s_{m}^{k},$$

To avoid dramatic power fluctuation during the pre-processing, we introduce the scaling factor of $\beta$ designed for maintaining either the loose power-constraint of $\mathbb{E}[\|x\|^2] = 1$ or the strict power-constraint of $\|x\|^2 = 1$, which are thus denoted by $\beta_l$ and $\beta_s$, respectively.

As a natural design, the TPC matrix has to ensure that no energy leaks into the unintended RA patterns. Hence, the classic linear Channel Inversion (CI)-based TPC [19], [20] may be used, which is formulated as

$$P = H^H (H H^H)^{-1},$$

where the power-normalisation factor of the output power after pre-processing is given by

$$\beta_l = \frac{N_r}{\text{Tr} \left( H H^H \right)^{-1}},$$

$$\beta_s = \frac{N_a}{\text{Tr} \left( H H^H \right)^{-1}} N_a,$$

The stringent power-constraint of (5) is less common than the 150 loose power-constraint of (4). The former prevents any of the 151 power fluctuations at the transmitter, which was also considered 152 in [19]. For completeness, we include both power-constraints in this paper.

C. GPSM Receiver

The signal observed at the $N_r$ RAs may be written as

$$y = \sqrt{\beta/N_a} H P s_{m}^{k} + w,$$

where $w \in \mathbb{C}^{N_r \times 1}$ is the circularly symmetric complex Gaussian noise vector with each entry having a zero mean and a variance of $\sigma^2$, i.e. we have $\mathbb{E}[|w|^2] = \sigma^2 I_{N_r}$, while $H \in \mathbb{C}^{N_r \times N_t}$ represents the MIMO channel involved. We assume furthermore that each entry of $H$ undergoes frequency-flat Rayleigh fading and it is uncorrelated between different super-symbols, while remains constant within the du-163 ration of a super-symbol’s transmission. The super-symbols transmitted are statistically independent from the noise.

At the receiver, the joint detection of both the conventional 166 modulated symbols $b_{m}$ and of the spatial symbol $k$ obeys the 167 Maximum Likelihood (ML) criterion, which is formulated as

$$\hat{m}_1, \ldots, \hat{m}_{N_a}, \hat{k} = \arg \min_{s_{m}^{k} \in \mathcal{B}} \left\{ \| y - \sqrt{\beta/N_a} H P s_{m}^{k} \|^2 \right\},$$

where $\mathcal{B} = C \times A^{N_a}$ is the joint search space of the super-symbol $s_{m}^{k}$. Alternatively, decoupled or separate detection may also be employed, which treats the detection of the conventional
modulated symbols $b_m$ and the spatial symbol $k$ separately. In this reduced-complexity variant, \(^1\) we have

$$\hat{k} = \arg \max_{k \in [1, |\mathcal{C}|]} \left\{ \sum_{i=1}^{N_c} \left| C(t,i) \right|^2 \right\},$$  

(8)

$$\hat{m}_i = \arg \min_{n_i \in [1, |\mathcal{M}|]} \left\{ \left| y_{\hat{v}_i} - \sqrt{\frac{\beta}{N_c}} h_{\hat{v}_i} p_{\hat{v}_i} b_{\hat{m}_i} \right|^2 \right\},$$  

(9)

where $h_{\hat{v}_i}$ is the $\hat{v}_i$th row of $H$ representing the channel between the $\hat{v}_i$th RA and the transmitter, while $p_{\hat{v}_i}$ is the $\hat{v}_i$th column of $P$ representing the $\hat{v}_i$th TPC vector. Thus, correct detection is declared, when we have $\hat{k} = k$ and $\hat{m}_i = m_i, \forall i$. \(^2\)

Remarks: Note that the complexity of the ML detection of (7) is quite high, which is on the order determined by the super-alphabet $\mathcal{B}$, hence obeying $O(|\mathcal{C}|M^{N_c})$. By contrast, the decoupled detection of (8) and (9) facilitates a substantially reduced complexity compared to that of (7). More explicitly, the complexity imposed by detecting $N_c$ conventional modulated symbols, plus the complexity ($\kappa$) imposed by the comparisons invoked for non-coherently detecting the spatial symbol of (8), which may be written as $O(N_cM + \kappa)$. Further discussions about the detection complexity of the decoupled detection of the GPSM scheme may be found in [17], where the main conclusion is that the complexity of the decoupled detection of the GPSM scheme is no higher than that of the conventional MIMO scheme corresponding to $N_a = N_r$. \(^3\)

III. PERFORMANCE ANALYSIS

We continue by investigating the DCMC capacity of our GPSM scheme, when the joint detection scheme of (7) is used and then quantify its achievable rate, when the realistic decoupled detection of (8) and (9) is employed. The achievable rate expression requires the theoretical BER/SER analysis of the GPSM scheme, which provides more insights into the inner nature of our GPSM scheme.\(^4\)

A. DCMC Capacity and Achievable Rate

Both Shannon’s channel capacity and its MIMO generalisation are maximized, when the input signal obeys a Gaussian distribution [22]. Our GPSM scheme is special in the sense that the spatial symbol conveys integer values constituted by the RA pattern index, which does not obey the shaping requirements of Gaussian signalling. This implies that the channel capacity of the GPSM scheme depends on a mixture of a continuous and a discrete input. Hence, for simplicity’s sake, we discuss the DCMC capacity and the achievable rate of our GPSM scheme in the context of discrete-input signalling for both the spatial and for the conventional modulated symbols mapped to it.

1) DCMC Capacity: Upon recalling the received signal observed at the $N_r$ RAs expressed in (6), the conditional probability of receiving $y$ given that a $M = |\mathcal{C}|M^{N_c}$-ary super-symbol $s_r \in \mathcal{B}$ was transmitted over Rayleigh channel and subjected to the TPC of (3) is formulated as

$$p(y|s_r) = \frac{1}{\pi \sigma^2} \exp \left\{ -\frac{\|y - Gs_r\|^2}{\sigma^2} \right\},$$  

(10)

where $G = \sqrt{\beta/N_c}HP$. The DCMC capacity of the ML-based joint detection of our GPSM scheme is given by [23]

$$C = \max_{p(s_1,\ldots,s_M)} \sum_{\tau=1}^M \int p(y, s_{\tau}) \log_2 \left( \frac{p(y|s_{\tau})}{\sum_{\epsilon=1}^M p(y, s_\epsilon)} \right) dy,$$  

(11)

which is maximized, when we have $p(s_{\tau}) = 1/M, \forall \tau$ [23]. Furthermore, we have

$$\log_2 \left( \frac{p(y|s_{\tau})}{\sum_{\epsilon=1}^M p(y, s_\epsilon)} \right) = \log_2 \left( \frac{p(y|s_{\tau})}{\sum_{\epsilon=1}^M p(y, s_\epsilon) p(s_{\epsilon})} \right) = -\log_2 \left( \frac{1}{M} \sum_{\epsilon=1}^M p(y|s_\epsilon) \right)$$

$$= \log_2(M) - \log_2 \sum_{\epsilon=1}^M \exp(\Psi),$$  

(12)

where substituting (10) into (12), the term $\Psi$ is expressed as

$$\Psi = -\frac{\|G(s_r - s_{\tau}) + w\|^2 + \|w\|^2}{\sigma^2}. $$  

(13)

Finally, by substituting (12) into (11) and exploiting that $p(s_{\tau}) = 233/1/M, \forall \tau$, we have

$$C = \log_2(M) - \frac{1}{M} \sum_{\tau=1}^M \mathbb{E}_{G,w} \left[ \log_2 \sum_{\epsilon=1}^M \exp(\Psi) \right].$$  

(14)

2) Achievable Rate: The above DCMC capacity expression implicitly relies on the ML-based joint detection of (7), which has a complexity on the order of $O(\mathcal{M})$. When the reduced-complexity decoupled detection of (8) and (9) is employed, we estimate the achievable rate based on the mutual information $I(z; \hat{z})$ per bit measured for our GPSM scheme between the input bits $z \in [0, 1]$ and the corresponding demodulated output bits $\hat{z} \in [0, 1]$.\(^5\)

The mutual information per bit $I(z; \hat{z})$ is given for the Binary Symmetric Channel (BSC) by [22]:

$$I(z; \hat{z}) = H(z) - H(z|\hat{z}),$$  

(15)
where \( H(z) = -\sum z P_z \log_2 P_z \) represents the entropy of the input bits \( z \) and \( P_z \) is the Probability Mass Function (PMF) of \( z \). It is noted furthermore that we have \( H(z) = 1 \), when we adopt the common assumption of equal-probability bits, i.e. \( P_z = 1/2 \). On the other hand, the conditional entropy \( H(z|\tilde{z}) \) represents the average uncertainty about \( z \) after observing \( \tilde{z} \).

Importantly, we have Lemma III.1 for the expression of \( \delta_{k_{\text{ant}}} \) acting as a correction factor in (22).

**Lemma III.1.** (Proof in Appendix A): The generic expression of the correction factor \( \delta_{k_{\text{ant}}} \) for \( k_{\text{ant}} \) bits of information is given by:

\[
\delta_{k_{\text{ant}}} = \delta_{k_{\text{ant}}-1} + \frac{2k_{\text{ant}}-1 - 2k_{\text{ant}}-1}{2k_{\text{ant}}-1} 
\]

where given \( \delta_0 = 0 \), we can recursively determine \( \delta_{k_{\text{ant}}} \).

Furthermore, by considering (21) and (22), the achievable rate expressed in (18) may be written as

\[
R \approx k_{\text{ant}} I \left( \frac{e^{\text{ant}}}{k_{\text{ant}}} \right) + N_{a k_{\text{mod}}} I \left( \frac{e^{\text{mod}}}{k_{\text{mod}}} \right). 
\]

Hence, as suggested by (19), (20) and (24), we find that both the 273 average error probability as well as the achievable rate of our 274 GPSM scheme requires the entries of \( e^{\text{ant}} \) and \( e^{\text{mod}} \), which 275 will be discussed as follows.

2) **Upper Bound of** \( e^{\text{ant}} \). We commence our discussion by 277 directly formulating the following lemma:

**Lemma III.2.** (Proof in Appendix B): The upper bound of 279 the analytical SER of the spatial symbol of our GPSM scheme 280 relying on CI TPC may be formulated as:

\[
e^{\text{ant}} \leq e^{\text{ant}} = 1 - N_a \int_0^{\infty} \int \left( \frac{N_e}{N_a} f_{X Z}(g) \right)^{N_a} f_{X Z}(g; \lambda) d\lambda \int \frac{f_{X}(\lambda)}{d\lambda} d\lambda,
\]

where \( f_{X Z}(g) \) represents the Cumulative Distribution Function (CDF) of a chi-square distribution having two degrees of freedom 283 and, while \( f_{X Z}(g; \lambda) \) represents the Probability Distribution 284 Function (PDF) of a non-central chi-square distribution having two 285 degrees of freedom and non-centrality given by

\[
\lambda = \beta \frac{N_a}{N_e},
\]

with its PDF of \( f_{X}(\lambda) \) and \( \sigma^2 = \sigma^2/2 \). Finally, equality of (25) 287 holds when \( N_e = 1 \).

Moreover, the PDF of \( f_{X}(\lambda) \) is formulated in Lemma III.3 and Lemma III.4, respectively, when either the loose or strict 290 power-normalisation factor of (4) and (5) is employed.

**Lemma III.3 (Proof in Appendix C):** When CI TPC is employed and the loose power-normalisation factor of (4) is used, the 293 distribution \( f_{X}(\lambda) \) of the non-centrality \( \lambda \) is given by:

\[
f_{X}(\lambda) = \frac{2N_e}{\lambda^2 N_e \sigma^2} f_U \left( \frac{2N_e}{\lambda N_e \sigma^2} \right),
\]

where by letting \( U = \text{Tr}[(H H^H)^{-1}] \), we have \( f_U(\cdot) \), which constitutes the derivative of \( F_U(\cdot) \) and it is given in (50) of Appendix C.

**Lemma III.4 (Proof in Appendix D):** When CI TPC is 299 employed and the stringent power-normalisation factor of (5) is used, the distribution \( f_{X}(\lambda) \) of the non-centrality \( \lambda \) is given by:

\[
f_{X}(\lambda) = \frac{N_e N_a - N_e + 1}{(N_e - N_e)!} \frac{\sigma^2}{2} e^{-\lambda N_e \sigma^2/2} \left( \frac{\lambda \sigma^2}{2} \right)^{N_e - N_e}.
\]
301 3) Upper Bound of $\overline{e}^s_{\text{mod}}$: Considering a general case of $N_r$ as well as $N_a$ and assuming that the RA pattern $C(k)$ was activated, after substituting (3) into (6), we have:

$$y_{v_i} = \sqrt{\beta/N_a}b_{m_i} + w_{v_i}, \quad \forall v_i \in C(k),$$

$$y_{u_i} = w_{u_i}, \quad \forall u_i \in \tilde{C}(k),$$

304 where $\tilde{C}(k)$ denotes the complementary set of the activated RA pattern $C(k)$ in $C$. Hence, we have the signal to noise ratio (SNR) given as

$$\gamma = \gamma_{v_i} = \frac{\beta N_a}{\sigma^2} = \frac{\lambda}{2}, \quad \forall v_i$$

307 and for the remaining deactivated RAs in $\tilde{C}(k)$, we have only 308 random noises of zero mean and variance of $\sigma^2$.

309 The SER $e^s_{\text{mod}}$ of the conventional modulated symbol $b_{m_i} \in A$ in the absence of spatial symbol errors may be upper 311 bounded by [24]:

$$e^s_{\text{mod}} < N_{\text{min}} \int_0^\infty Q(\sqrt{\gamma/2}) f_\gamma(\gamma) d\gamma = e^{s,\text{ub}}_{\text{mod}},$$

312 where in general $f_\gamma(\gamma)$ has to be acquired by the empirical 313 histogram based method. When Lemma III.3 or Lemma III.4 314 is exploited, $f_\gamma(\gamma)$ is a scaled version of $f_\lambda(\lambda)$, i.e. we have 315 $f_\gamma(\gamma) = 2f_\lambda(2\gamma)$. Moreover, $d_{\text{min}}$ is the minimum Euclidean 316 distance in the conventional modulated symbol constellation, 317 $N_{\text{min}}$ is the average number of the nearest neighbours separated 318 by $d_{\text{min}}$ in the constellation and $Q(\cdot)$ denotes the Gaussian 319 Q-function.

320 When taking into account of the spatial symbol errors, we 321 have Lemma III.5 for the upper bound of $e^s_{\text{mod}}$.

322 Lemma III.5. (Proof in Appendix E): Given the $k$th activated 323 RA pattern, the SER of the conventional modulated symbols in 324 the presence of spatial symbol errors can be upper bounded by:

$$e^s_{\text{mod}} < 1 - e^{s,\text{ub}}_{\text{mod}}$$

$$+ e^{s,\text{ub}}_{\text{eff}} \sum_{\ell \neq k} N_c e^{s,\text{mod}}_{\text{eff}} + N_a e^s_{\text{mod}} = e^{s,\text{ub}}_{\text{mod}},$$

325 where $N_c$ and $N_d = (N_a - N_c)$ represent the number of config- 326 mon and different RA between $C(\ell)$ and $C(k)$, respectively.

327 Mathematically we have $N_c = \sum_{i=1}^{N_a} [C(\ell, i) \in C(k)]$. More- 328 over, $e_c^s = (M - 1)/M$ is SER as a result of random guess.

329 4) Upper Bound of $e_{\text{ant}}^s$ and $e_{\text{eff}}^s$: By substituting (25) and 330 (33) into (19) and (20), we arrive at the upper bound of the 331 average symbol and bit error probability as

$$e_{s,\text{mod}}^s = \frac{(e^{s,\text{ub}}_{\text{mod}} + N_a e^{s,\text{ub}}_{\text{mod}})}{(1 + N_a)}$$

$$e_{s,\text{ant}}^s = \frac{(e^{s,\text{ub}}_{\text{ant}} + N_a e^{s,\text{ub}}_{\text{mod}})}{k_{\text{ant}}},$$

Similarly, by substituting (25) and (33) into (24), we obtain the 332 lower bound of the achievable rate as

$$R_{\text{th}} = k_{\text{ant}} I (\delta_{\text{ant}} e_{s,\text{ant}}^s + N_a e_{s,\text{mod}}^s).$$

IV. NUMERICAL RESULTS

We now provide numerical results for characterizing both the 335 DCMC capacity of our GPSM scheme and for demonstrating the 336 accuracy of our analytical error probability results.

A. DCMC Capacity

1) Effect of the Number of Activated RAs: Fig. 1 character- 339 ises the DCMC capacity versus the SNR of the CI TPC 340 aided GPSM scheme based on the loose power-normalisation 341 factor of (4) under $\{N_t, N_r\} = \{8, 4\}$ and employing QPSK, while having $N_a = \{1, 2, 3, 4\}$ activated RAs.

2) Effect of the Number of Antennas: We further investigate the attainable bandwidth efficiency by replacing the SNR used in Fig. 1 by the SNR per bit in Fig. 2, 357 where we have $\text{SNR}_{\text{db}} = \frac{\text{SNR}_{\text{db}} - 10\log_{10}(C/N_a)}{N_a}$. It is clear that the lower $N_a$, the higher the bandwidth efficiency attained in the low range of $\text{SNR}_{\text{db}}$. Importantly, the achievable bandwidth efficiency of $N_a = 3$ is 361
consistent and significantly higher than that achieved by
$N_a = 4$, before they both converge to 8 bits/symbol/Hz at their
maximum. Overall, there is always a beneficial configuration
for our GPSM scheme that offers the same bandwidth efficiency
as that of its conventional MIMO counterpart, which is achieved
at a lower SNR per bit.

2) Robustness to Impairments: Like in all TPC schemes,
an important aspect related to GPSM is its resilience to CSIT
inaccuracies. In this paper, we let $H = H_a + H_i$, where $H_a$
represents the matrix hosting the average CSI, with each entry
obeying the complex Gaussian distribution of $h_a \sim \mathcal{CN}(0, \sigma_a^2)$
and $H_i$ is the instantaneous CSI error matrix obeying the
complex Gaussian distribution of $h_i \sim \mathcal{CN}(0, \sigma_i^2)$, where we
have $\sigma_a^2 + \sigma_i^2 = 1$. As a result, only $H_a$ is available at the
transmitter for pre-processing.

Another typical impairment is antenna correlation. The
correlated MIMO channel is modelled by the widely-used
Kronecker model, which is written as
$$H = (R_t^{1/2}) G (R_r^{1/2})^T,$$
with $G$ representing the original MIMO channel imposing no
correlation, while $R_t$ and $R_r$ represents the correlations at the
transmitter and receiver side, respectively, with the correlation
entries given by $R_t(i,j) = \rho_t|^{i-j}$ and $R_r(i,j) = \rho_r|^{i-j}$.

Figs. 3 and 4 characterise the effect of imperfect CSIT
associated with $\sigma_i = 0.4$ and of antenna correlation of $\rho_t =
\rho_r = 0.3$ on the attainable DCMC capacity versus the SNR
for our CI TPC aided GPSM scheme with the loose power-
normalisation factor of (4), respectively, under $\{N_t, N_r\} =
\{8, 4\}$ and employing QPSK having $N_a = \{1, 2, 3, 4\}$ activated
RAS. It can be seen that both impairments result
into a degraded DCMC capacity. Observe in Fig. 3 for imperfect
CSIT that the degradation of the conventional MIMO
associated with $N_a = 4$ and marked by the triangle is larger
than that of our GPSM scheme corresponding $N_a = \{1, 2, 3\}$.
On the other hand, as seen in Fig. 4, roughly the same level of
degradation is observed owing to antenna correlation.

3) Effect of Modulation Order and MIMO Configuration:
Fig. 5 characterises the DCMC capacity versus the SNR
of our CI TPC aided GPSM scheme relying on the loose 399
power-normalisation factor of (4) under $\{N_t, N_r\} = \{8, 4\}$ and employing various conventional modulation schemes having 400
$N_a = \{1, 2\}$ activated RAS. It can be seen that the higher the 402
modulation order $M$, the higher the achievable DCMC capac-
ity. Furthermore, for a fixed modulation order $M$, the higher 404
the value of $N_a$, the higher the achievable DCMC capacity 405
becomes as a result of the information embedded in the spatial 406
symbol.

Fig. 6 characterises the DCMC capacity versus the SNR 408
for our CI TPC aided GPSM scheme for the loose power- 409
normalisation factor of (4) under different settings of $\{N_t, N_r\} =\{4\}
with $N_t/N_r = 2$ and employing QPSK, while having $N_a = 411
\{1, 2\}$ activated RAs. It can be seen in Fig. 6 that for a fixed 412
MIMO setting, the higher the value of $N_a$, the higher the 413
Fig. 5. DCMC capacity versus the SNR of our CI TPC aided GPSM scheme relying on the loose power-normalisation factor of (4) under $\{N_t,N_r\} = (8,4)$ and employing various conventional modulation schemes having $N_a = \{1,2\}$ activated RAs.

Fig. 6. DCMC capacity versus the SNR for our CI TPC aided GPSM scheme for the loose power-normalisation factor of (4) under different settings of $\{N_t,N_r\}$ with $N_t/N_r = 2$ and employing QPSK, while having $N_a = \{1,2\}$ activated RAs.

414 DCMC capacity becomes. Importantly, for a fixed $N_a$, the larger the size of the MIMO antenna configuration, the higher 415 the DCMC capacity.

417 B. Achievable Rate

418 1) Error Probability: Figs. 7–10 characterize the GPSM scheme’s SER as well as the BER under both the loose 420 power-normalisation factor of (4) and the stringent power-normalisation factor of (5) for $\{N_t,N_r\} = \{16,8\}$ and employing QPSK, respectively. From Figs. 7–10, we recorded the 423 curves from left to right corresponding to $N_a = \{1,2,4,6\}$. For 424 reasons of space-economy and to avoid crowded figures, our 425 results for $N_a = \{3,5,7\}$ were not shown here, but they obey 426 the same trends.

It can be seen from Figs. 7 and 9 that our analytical SER 427 results of (34) form tight upper bounds for the empirical sim- 428 ulation results. Hence they are explicitly referred to as ‘tight 429 upper bound’ in both figures. Additionally, a loose upper bound 430 of the GPSM scheme’s SER is also included, which may be 431 written as

$$e_{\text{SER}}^{\text{ub},\text{loose}} = 1 - e_{\text{ant}}^{\text{ub}} \left(1 - e_{\text{mod}}^{\text{ub}}\right).$$ (37)

Note that in this loose upper bound expression, $e_{\text{mod}}^{\text{ub}}$ of (32) is 433 required rather than $e_{\text{ant}}^{\text{ub}}$ of (33). This expression implicitly 434 assumes that the detection of (8) and (9) are independent. 435 However, the first-step detection of (8) significantly affects the 436 second-step detection of (9). Hence, the loose upper bound 437 shown by the dash-dot line is only tight for $N_a = 1$ and 438 becomes much looser upon increasing $N_a$, when compared to 439 the tight upper bound of (34).
Similarly, when the GPSM scheme’s BER is considered in Figs. 8 and 10, our analytical results of (35) again form tight upper bounds for the empirical results.

2) Separability: To access the inner nature of first-step detection of (8), Fig. 11 reveals the separability between the activated RAs and deactivated RAs in our GPSM scheme, where the PDF of (44) and (45) were recorded both for $\text{SNR} = -5$ dB (left subplot) and for $\text{SNR} = 0$ dB (right subplot) respectively for the same snapshot of MIMO channel realisation with the aid of CI TPC and the loose power-normalisation factor of (4) under \( \{N_t, N_r\} = \{16, 8\} \) and employing QPSK. By comparing the left subplot to the right subplot, it becomes clear that the higher the SNR, the better the separability between the activated and the deactivated RAs, since the mean of the solid curves representing (44) move further apart from that of the dashed curve representing (45). Furthermore, as expected, the lower $N_a$, the better the separability becomes, as demonstrated in both subplots of Fig. 11.

3) Comparison: Finally, Fig. 12 characterizes the comparison between the DCMC capacity (14) of our GPSM scheme relying implicitly on the ML-based joint detection and its lower bound of the achievable rate in (36) relying on the low-complexity decoupled detection, where we use CI TPC with the loose power-normalisation factor of (4) under \( \{N_t, N_r\} = \{16, 8\} \) and employing QPSK having $N_a = \{1, 2, 3\}$. It is clear that the DCMC capacity is higher than the achievable rate for each $N_a$ considered, although both of them converge to the same value, when the SNR is sufficiently high. Noticeably, the discrepancy between the two quantities before their convergence is wider, when $N_a$ is higher. This is because the higher $N_a$, the lower the achievable rate at low SNRs.
which is shown by comparing the solid curves. This echoes our observations of Fig. 11, namely that a higher \( N_a \) leads to a reduced separability and consequently both to a higher overall error probability and to a lower achievable rate. In fact, the achievable rate becomes especially insightful after being compared to the DCMC capacity, where we may tell how a realistic decoupled detection performs and how far its performance is from the DCMC capacity.

V. CONCLUSION

In this paper, we introduced the concept of our GPSM scheme and carried out its theoretical analysis in terms of both its DCMC capacity as well as its achievable rate relying on our analytical upper bound of the SER and the BER expressions, when a low-complexity decoupled detector is employed. Our numerical results demonstrate that the upper bound introduced is tight and the DCMC capacity analysis indicates that our GPSM scheme constitutes a flexible MIMO arrangement. Our future work will consider a range of other low-complexity MIMO schemes, such as the receive antenna selection and the classic SM, in the context of large-scale MIMOS.

Furthermore, the insights of our error probability and capacity analysis are multi-folds:

- It can be seen that there is a gap between the DCMC capacity relying on ML detection and the achievable rate of decoupled detection. Thus, a novel detection method is desired for closing this gap and for striking a better trade-off between the performance attained and the complexity imposed.
- The error probability derived serves as a tight upper bound of our GPSM performance. This facilitates the convenient study of finding beneficial bit-to-symbol mapping and error-probability balancing between the spatial symbols and conventional modulated symbols [25]. Otherwise, excessive-complexity bit-by-bit Monte-Carlo simulations would be required.
- Furthermore, both the capacity and error probability analysis provide a bench-mark for conducting further research on antenna selection techniques for our GPSM scheme, where different criteria may be adopted either for maximizing the capacity or for minimizing the error probability, again without excessive-complexity bit-by-bit Monte-Carlo simulations.

APPENDIX A

PROOF OF LEMMA III.1

Let \( A_{k^*\text{nt}} \) denote the alphabet of the spatial symbol having \( k^*\text{nt} \) bits of information. Then the cardinality of the alphabet \( A_{k^*\text{nt}} \) is twice higher compared to that of \( A_{k^*\text{nt}-1} \). Thus, \( A_{k^*\text{nt}} \) may be constructed by two sub-alphabets of \( A_{k^*\text{nt}-1} \), represented by 0 and 1, respectively. We may thereafter refer to the alphabet of \( A_{k^*\text{nt}-1} \) preceded by the above-mentioned with 0 (1) as zero-alphabet (one-alphabet).

Assuming that the spatial symbol representing \( k^*\text{nt} \) zeros was transmitted, we may then calculate the total number of \( k^*\text{nt} \) pair-wise bit errors \( \epsilon_0 \) in the above zero-alphabet. Hence, the number of pair-wise bit errors \( \epsilon_1 \) in the one-alphabet is simply \( \epsilon_1 = \epsilon_0 + A \), where \( A = 2^{k^*\text{nt}} \) accounts for the difference in the first preceding bit. Hence the total number of pair-wise bit errors \( \epsilon \) is \( \epsilon = 2\epsilon_0 + 2^{k^*\text{nt}} \). Taking into account an equal probability of \( 1/(2^{k^*\text{nt}} - 1) \) for each possible spatial symbol \( 531 \), we arrive at the correction factor given by \( \delta_{k^*\text{nt}} = (2\epsilon_0 + 2^{k^*\text{nt}})/(2^{k^*\text{nt}} - 1) \).

Since \( \epsilon_0 \) represents the total number of pair-wise bit errors \( \epsilon \) corresponding to case of \( (k^*\text{nt}) \) bits of information, we have \( \epsilon_0 = (2^{k^*\text{nt}} - 1)\delta_{k^*\text{nt}} \). Hence the resultant expression of the correction factor may be calculated recursively according to (23) after some further manipulations.

APPENDIX B

PROOF OF LEMMA III.2

Considering a general case of \( N_a \) as well as \( N_a \) and assuming that the RA pattern \( C(k) \) was activated, after substituting (3) into (6), we have:

\[
y_{i_1} = \sqrt{\beta/N_a} b_{m_1} + u_{i_1}, \quad \forall i_1 \in C(k),
\]

\[
y_{i_2} = u_{i_2}, \quad \forall i_2 \in \bar{C}(k),
\]

where \( C(k) \) denotes the complementary set of the activated RA pattern \( C(k) \) in \( C \). Furthermore, upon introducing \( \sigma_0^2 = \sigma^2/2 \), we have:

\[
|y_{i_1}|^2 = \mathcal{R}(y_{i_1})^2 + \mathcal{I}(y_{i_1})^2 \sim \mathcal{N}(\sqrt{\beta/N_a} \mathcal{R}(b_{m_1}), \sigma_0^2) + \mathcal{N}(\sqrt{\beta/N_a} \mathcal{I}(b_{m_1}), \sigma_0^2),
\]

\[
|y_{i_2}|^2 = \mathcal{R}(u_{i_2})^2 + \mathcal{I}(u_{i_2})^2 \sim \mathcal{N}(0, \sigma_0^2) + \mathcal{N}(0, \sigma_0^2),
\]

where \( \mathcal{R}(\cdot) \) and \( \mathcal{I}(\cdot) \) represent the real and imaginary operators, respectively. As a result, by normalisation with respect to \( \sigma_0^2 \), we have the following observations:

\[
|y_{i_1}|^2 \sim \chi_2^2(g; \lambda_{i_1}), \quad \forall i_1 \in C(k),
\]

\[
|y_{i_2}|^2 \sim \chi_2^2(g), \quad \forall i_2 \in \bar{C}(k),
\]

where the non-centrality is given by \( \lambda_{i_1} = \beta |b_{m_1}|^2/N_a \sigma_0^2 \). Exploiting the fact that \( \mathbb{E}[|b_{m_1}|^2] = 1, \forall i \) (or \( |b_{m_1}|^2 = 1, \forall i \) for 551 PSK modulation), we have \( \lambda = \lambda_{i_1}, \forall i_1 \). Note that \( \lambda \) is also a random variable obeying the distribution of \( f_\lambda(\lambda) \).

Recall from (8) that the correct decision concerning the spatial symbols occurs, when \( \sum_{i=1}^{N_a} |y_{i_1}|^2 \) is the maximum. 555

By exploiting the fact that \( \mathbb{E}[|C(k)|] = \Delta \), the correct detection probability \( \Delta \) of the spatial symbols given the non-centrality \( \lambda \), 557

\footnote{By assuming equal-probability erroneously detected patterns, a spatial symbol may be mistakenly detected as any of the other spatial symbols with equal probability. Let us now give an example for highlighting the rationale of introducing the correction factor. For example, spatial symbol ‘0’ carrying bits [0,0] was transmitted, it would result into a one-bit difference when the spatial symbol ‘1’ carrying [0,1] or ‘2’ carrying [1,0] was erroneously detected. However, it would result into a two-bits difference when spatial symbol ‘3’ carrying [1,1] was erroneously detected. This corresponds to four bit errors in total for these three cases, thus a correction factor of \( 4/3 \) is needed when converting the symbol error ratio to bit error ratio.}
when the RA pattern $C(k)$ was activated may be lower bounded as in (46). (See equation at bottom of page) More explicitly,

- equation (a) serves as the lower bound, since it sets the most strict condition for the correct detection, when each metric $y_{u_j}$ of the inactivated RA indices in $C(k)$ is lower than each metric $g_{v_i}$ of the activated RA indices in $C(k)$.

Note that, equality holds when $N_a = 1$;

- equation (b) follows from the fact that the $N_a$ random variables $|y_{u_j}|$ are independent of each other;

- equation (c) follows from the fact that the $(N_r - N_a)$ random variables $|y_{u_j}|$ are independent and equation (d) follows from the fact that the $N_a$ independent variables of $|y_{u_j}|$ and the $(N_r - N_a)$ independent variables of $|y_{v_i}|$ are both identically distributed.

As a result, after averaging over the distribution of $f_\lambda(\lambda)$, the analytical SER $c_{\text{ant}}^*$ of the spatial symbol in our GPSM scheme may be upper bounded as in (25). In general, the expression of $f_\lambda(\lambda)$ can be acquired with the aid of the empirical histograms based method, while in case the loose/stringent power normalisation factor of (4)/(5) is used, the analytical expression for $f_\lambda(\lambda)$ is given in Lemma III.3/Lemma III.4.

**APPENDIX C**

**PROOF OF LEMMA III.3**

Upon expanding the expression of $\lambda$ in (26) by taking into account (4), we have:

$$\lambda = \frac{\beta_1}{N_a\sigma_0^2} = \frac{N_r}{N_a\sigma_0^2 \text{Tr}[(HH^H)^{-1}]}.$$  \hspace{1cm} (47)

Consider first the distribution of $\text{Tr}[(HH^H)^{-1}]$ and let $W = HH^H$. Since the entries of $H$ are i.i.d. zero-mean unit-variance complex Gaussian random variables, $W$ obeys a complex Wishart distribution. Hence the joint PDF of its eigenvalues $\{\lambda_{W_i}\}_{i=1}^{N_r}$ is given by [26], [27]

$$f_W(\lambda_i) = \frac{K^{-1}}{N_r!} \prod_i e^{-\lambda_i} \lambda_i^{N_r - N_r} \prod_{i<j} (\lambda_i - \lambda_j)^2,$$

where $K$ is a normalising factor. Thus for its inverse $U = W^{-1}$, we have

$$f_U(\lambda_i) = \frac{1}{\lambda_i^{N_r}} f_W(\lambda_i).$$  \hspace{1cm} (49)

Furthermore, since $\text{Tr}[U] = \sum_{i=1}^{N_r} \lambda_i$, where $\{\lambda_i\}_{i=1}^{N_r}$ is the eigenvalues of $U$, we have the CDF of $\text{Tr}[U]$ given by (50), where $T_1 = T$ and $T_2 = 1/T$, while $\forall j > 1$.

$$T_j = T - \sum_{i=1}^{j-1} \lambda_i, \quad T_j = 1/T.$$  \hspace{1cm} (51)

**APPENDIX D**

**PROOF OF LEMMA III.4**

Upon expanding the expression of $\lambda$ in (26) by taking into account (5), we have:

$$\lambda = \frac{\beta_2}{N_a\sigma_0^2} = \frac{1}{\sigma_0^2 s^H (HH^H)^{-1} s}.$$  \hspace{1cm} (52)

$$\Delta \geq \prod_{i=1}^{N_r} \prod_{u_j \in C(k)} P\left(\{y_{u_j}\}_i^N < g_{v_i}\right)$$

$$\geq \prod_{i=1}^{N_a} \prod_{u_j \in C(k)} P\left(\{y_{u_j}\}_i^N < g_{v_i}\right)$$

$$\geq \prod_{i=1}^{N_a} \prod_{u_j \in C(k)} P\left(\{y_{u_j}\}_i^N < g_{v_i}\right)$$

$$\geq \prod_{i=1}^{N_a} \prod_{u_j \in C(k)} \left\{ \int_0^{\infty} f_{\lambda_i^2}(\lambda_i) \frac{N_a - N_r}{N_a} f_{\lambda_i^2}(\lambda_i) \right\}$$

$$\geq \prod_{i=1}^{N_a} \prod_{u_j \in C(k)} \left\{ \int_0^{\infty} f_{\lambda_i^2}(\lambda_i) \frac{N_a - N_r}{N_a} f_{\lambda_i^2}(\lambda_i) \right\}$$

$$\frac{\lambda_1^2}{\sigma_0^2 s^H (HH^H)^{-1} s}.$$  \hspace{1cm} (53)

$$F_{\text{Tr}[U]}(T) = \int_0^{T_1} \cdots \int_0^{T_{N_r}} f_U(\{\lambda_{U_i}\}_{i=1}^{N_r}) d\lambda_{U_{N_r}} \cdots d\lambda_{U_{1}} = \int_0^{\infty} \cdots \int_0^{\infty} f_W(\{\lambda_{W_i}\}_{i=1}^{N_r}) d\lambda_{W_{N_r}} \cdots d\lambda_{W_{1}}$$

$$\frac{\lambda_1^2}{\sigma_0^2 s^H (HH^H)^{-1} s}.$$  \hspace{1cm} (54)
Since the entries of $H$ are i.i.d. zero-mean unit-variance 603 complex Gaussian random variables, $HH^H$ obeys a complex 604 Wishart distribution with $N_r$ dimensions and $2N_t$ degrees of 605 freedom, where we have:

$$HH^H \sim CW(\Sigma, N_r, 2N_t),$$

(52)

with $\Sigma = (1/2)I_{N_r}$ being the variance. By exploiting propo- 607 sition 8.9 from [28] and letting $\lambda_0 = \left(s^H (HH^H)^{-1} s\right)^{-1}$, we 608 have:

$$\lambda_0 \sim CW\left(\left(s^H \Sigma^{-1} s\right)^{-1}, 1, 2(N_t - N_r + 1)\right),$$

(53)

where $A \sim B$ stands for $A$ follows the distribution of $B$. 610 According to [28], the above one-dimensional complex-valued 611 Wishart distribution is actually a chi-square distribution with 612 $2(N_t - N_r + 1)$ degrees of freedom and scaling parameter of 613 $(s^H \Sigma^{-1} s)^{-1} = 1/2N_a$. Thus, the PDF of $\lambda_0$ may be explicitly 615 written as:

$$f_{\lambda_0}(\lambda_0) = \int_2 \left(2N_a \lambda_0; 2(N_t - N_r + 1)\right) \frac{e^{-\lambda_0 N_a} (2N_a \lambda_0)^{N_t - N_r}}{2(N_t - N_r + 1) !} \left(1 - \lambda_0 \frac{N_a^{N_t - N_r}}{N_t - N_r} \right)^{N_a - N_r} \left(\frac{N_t - N_r}{(N_t - N_r + 1)}\right)!$$

(54)

Finally, since $\lambda_0 = \sigma_0^2 \lambda$, we have $f_{\lambda}(\lambda) = \sigma_0^2 f_{\lambda_0}(\sigma_0^2 \lambda)$, which 617 is (28).

### Appendix E

## Proof of Lemma III.5

The SER of $e^{s}_{\text{mod}}$ is composed by the SER of $e^{s}_{\text{mod}}$, 620 when the detection of the spatial symbol is correct having a 621 probability of $(1 - e^{s}_{\text{ant}})$, plus the SER, when the detection 622 of the spatial symbol is erroneous having a probability of $e^{s}_{\text{ant}}$, 623 which is expressed as

$$e^{s}_{\text{mod}} = \left(1 - e^{s}_{\text{ant}}\right)e^{s}_{\text{mod}} + e^{s}_{\text{ant}} \sum_{\ell \neq k} P_{k \rightarrow \ell} \frac{N_c e^{s}_{\text{mod}} + N_d e^{s}_{o}}{N_a},$$

(55)

$$b \leq \left(1 - e^{s}_{\text{ant}}\right)e^{s}_{\text{ant}} + e^{s}_{\text{ant}} \sum_{\ell \neq k} P_{k \rightarrow \ell} \frac{N_c e^{s}_{\text{mod}} + N_d e^{s}_{o}}{N_a},$$

(56)

$$c \leq \left(1 - e^{s}_{\text{ant}}\right)e^{s}_{\text{mod}} + \frac{e^{s}_{\text{ant}}}{(2^{e^{s}_{\text{ant}} - 1})} \sum_{\ell \neq k} P_{k \rightarrow \ell} \frac{N_c e^{s}_{\text{mod}} + N_d e^{s}_{o}}{N_a},$$

(57)

$$d \leq \left(1 - e^{s}_{\text{ant}}\right)e^{s}_{\text{ant}} + e^{s}_{\text{ant}} \sum_{\ell \neq k} P_{k \rightarrow \ell} \frac{N_c e^{s}_{\text{mod}} + N_d e^{s}_{o}}{N_a},$$

(58)

625 Regarding the second additive term of (a), the true activated RA 626 pattern $C(k)$ may be erroneously deemed to be any of the other legitimate RA patterns $C(\ell) \in C, \ell \neq k$ with a probability of 627 $P_{k \rightarrow \ell}$, which we have to average over. As for the calculation of 628 the per-case error rates $E$, when $C(k)$ was erroneously detected 629 as a particular $C(\ell)$, we found that it was constituted by the error 630 rates of $e^{s}_{\text{mod}}$ for those $N_c$ RAs in common (which maybe 631 regarded as being partially correctly detected) and the error 632 rates of $e^{s}_{o}$ for those RAs that were exclusively hosted by $C(\ell)$, 633 but were excluded from $C(k)$. Furthermore, since only random 634 noise may be received by those $N_d$ RAs in $C(\ell)$, thus $e^{s}_{o}$ simply 635 represents the SER as a result of a random guess, i.e. we have 636 $e^{s}_{o} = (M - 1)/M$. Let us now provide some further detailed 637 discussions of the relations ranging from (b) to (d):

- relation (b) holds true, since $e^{s}_{\text{mod}}$ is a monotonically decreasing function of $e^{s}_{\text{mod}}$; thus it is upper bounded upon replacing $e^{s}_{\text{mod}}$ by $e^{s}_{\text{ant}}$;
- although it is natural that patterns with a higher $N_c$ would be more likely to cause an erroneous detection, we assume an equal probability of $P_{k \rightarrow \ell} = 1/(2^{e^{s}_{\text{ant}} - 1})$. The equal probability assumption thus puts more weight on the patterns having higher $N_d$, since we have $e^{s}_{o} \approx e^{s}_{\text{ant}}$. This 646 leads to the relation of (c).
- Note that, equality holds when $N_c = 1$, where $N_c = 0$ and $N_d = 1$;
- replacing $e^{s}_{\text{ant}}$ by $e^{s}_{\text{ant}}$ puts more weight on the second 649 additive term of (d), since having $e^{s}_{\text{ant}} > e^{s}_{\text{mod}}$ leads to 650 the relation of $A > e^{s}_{\text{ant}}$. As a result (d) also holds. 651 Again, equality holds when $N_a = 1$, where $e^{s}_{\text{ant}} = e^{s}_{\text{ant}}$ as indicated by Lemma III.2.

### Acknowledgment

The financial support of the EPSRC under the India-UK Advanced Technology Centre (IU-ATC), that of the EU under the Concerto project as well as that of the European Research 657 Council’s (ERC) Advance Fellow Grant is gratefully acknowledged.

### References


AUTHOR QUERIES

AUTHOR PLEASE ANSWER ALL QUERIES

AQ1 = Please be informed that the capital letters were removed from the terms “multiple input multiple output,” “generalised pre-coded aided spatial modulation,” “symbol error ratio,” “bit error ratio,” “discrete-input continuous-output memoryless channel,” and “signal to noise ratio” in the Abstract per IEEE style and also in other occurrences of these terms in lines 88 to 91 and 305 for the sake of consistency. Please check if it is correct.

AQ2 = Please provide keywords.

AQ3 = Please check changes made in first footnote and the addition of an Acknowledgment Section.

AQ4 = Please check if “30 journals” should be “30 papers” instead.

END OF ALL QUERIES
Error Probability and Capacity Analysis of Generalised Pre-Coding Aided Spatial Modulation

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Abstract—The recently proposed multiple input multiple output (MIMO) transmission scheme termed as generalized pre-coding aided spatial modulation (GPSM) is analyzed, where the key idea is that a particular subset of receive antennas is activated and the specific activation pattern itself conveys useful implicit information. We provide the upper bound of both the symbol error ratio (SER) and bit error ratio (BER) expression of the GPSM scheme of a low-complexity decoupled detector. Furthermore, the corresponding discrete-input continuous-output memoryless channel (DCMC) capacity as well as the achievable rate is quantified. Our analytical SER and BER upper bound expressions are confirmed to be tight by our numerical results. We also show that our GPSM scheme constitutes a flexible MIMO arrangement and there is always a beneficial configuration for our GPSM scheme that offers the same bandwidth efficiency as that of its conventional MIMO counterpart at a lower signal to noise ratio (SNR) per bit.

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I. INTRODUCTION

Multiple input multiple output (MIMO) systems constitute one of the most promising recent technical advances in wireless communications, since they facilitate high-throughput transmissions in the context of various standards [1]. Hence, they attracted substantial research interests, leading to the Vertical-Bell Laboratories Layered Space-Time (V-BLAST) scheme [2] and to the classic Space Time Block Coding (STBC) arrangement [3]. The point-to-point single-user MIMO systems are capable of offering diverse transmission functionalities in terms of multiplexing-diversity- and beam-forming gains. Similarly, Spatial Division Multiple Access (SDMA) employed in the uplink and multi-user MIMO techniques invoked in the downlink also constitute beneficial building blocks [4], [5]. The basic benefits of MIMOs have also been recently exploited in the context of the network MIMO concept [6], [7], for constructing large-scale MIMOs [8], [9] and for conceiving beneficial arrangements for interference-limited MIMO scenarios [10].

Despite having a plethora of studies on classic MIMO systems, their practical constraints, such as their I/Q imbalance, their transmitter and receiver complexity as well as the cost of their multiple Radio Frequency (RF) Power Amplifiers (PA) chains as well as their Digital-Analogue/Analogue-Digital (DA/AD) converters have received limited attention. To circumvent these problems, low complexity alternatives to conventional MIMO transmission schemes have also been proposed, such as the Antenna Selection (AS) [11], [12] and the Spatial Modulation (SM) [13], [14] philosophies. More specifically, SM and generalised SM [15] constitute novel MIMO techniques, which were conceived for providing a higher throughput than a single-antenna aided system, while maintaining both a lower complexity and a lower cost than the conventional MIMOs, since they may rely on a reduced number of RF conversion chains. To elaborate a little further, SM conveys extra information by mapping $\log_2(N_t)$ bits to the Transmit Antenna (TA) indices of the $N_t$ TAs, in addition to the classic 60 modulation schemes, as detailed in [13].

By contrast, the family of Pre-coding aided Spatial Modulation (PSM) schemes is capable of conveying extra information by appropriately selecting the Receive Antenna (RA) indices, as detailed in [16]. More explicitly, in PSM the indices of the 65 RA represent additional information in the spatial domain. As 66 a specific counterpart of the original SM, PSM benefits from both a low cost and a low complexity at the receiver side, 68 therefore it may be considered to be eminently suitable for 69 downlink transmissions [16]. The further improved concept of 70 Generalised PSM (GPSM) was proposed in [17], where comprehensive performance comparisons were carried out between 72 the GPSM scheme as well as the conventional MIMO scheme 73 and the associated detection complexity issues were discussed. 74 Furthermore, a range of practical issues were investigated, 75 namely the detrimental effects of realistic imperfect Channel 76 State Information at the Transmitter (CSIT), followed by a 77 low-rank approximation invoked for large-dimensional MI- 78 MOs. Finally, the main difference between our GPSM scheme 79 and the classic SM is that the former requires downlink pre- 80 processing and CSIT, although they may be considered as 81 a dual counterpart of each other and may hence be used in 82 a hybrid manner. Other efforts on robust PSM was reported 83 in [18].

As a further development, in this paper, we provide the theoretical analysis of the recently proposed GPSM scheme [17, 18], which is not available in the literature. More explicitly, both the discrete-input continuous-output memoryless channel (DCMC) capacity as well as the achievable rate are characterized.

Importantly, tight upper bounds of the symbol error ratio (SER) and bit error ratio (BER) expressions are derived, when a decoupled low-complexity detector is employed.

The rest of our paper is organised as follows. In Section II, we introduce the underlying concept as well as the detection methods of the GPSM scheme. This is followed by our analytical study in Section III, where both the DCMC capacity and the achievable rate as well as the SER/BER expressions are derived. Our simulation results are provided in Section IV, while we conclude in Section V.

II. System Model

A. Conceptual Description

Consider a MIMO system equipped with $N_t$ TAs and $N_r$ RAs, where we assume $N_t \geq N_r$. In this MIMO set-up, a maximum of $N_r$ parallel data streams may be supported, conveying a total of $k$ symbols in addition to the information carried by the conventional modulated symbols. Hence, the number of activated RAs out of $N_r$ RAs may be written as

$$\beta = \frac{N_r}{\text{Tr} \left[ (HH^H)^{-1} \right]}.$$  

The stringent power-constraint of (5) is less common than the 150 loose power-constraint of (4). The former prevents any of the 151 power fluctuations at the transmitter, which was also considered 152 in [19]. For completeness, we include both power-constraints in this paper.

C. GPSM Receiver

The signal observed at the $N_r$ RAs may be written as

$$y = \sqrt{\beta} \frac{N_a}{|H|^2} \frac{H^H}{s_m} + w,$$  

where $w \in \mathbb{C}^{N_r \times 1}$ is the circularly symmetric complex Gaussian noise vector with each entry having a zero mean and a variance of $\sigma^2$, i.e. we have $\mathbb{E} \left[ \|w\|^2 \right] = \sigma^2 I_{N_r}$, while $H \in \mathbb{C}^{N_r \times N_t}$ represents the MIMO channel involved. We assume 160 furthermore that each entry of $H$ undergoes frequency-flat 161 Rayleigh fading and it is uncorrelated between different super-162 symbol transmissions, while remains constant within the du-163 ration of a super-symbol’s transmission. The super-symbols 164 transmitted are statistically independent from the noise.

At the receiver, the joint detection of both the conventional 166 modulated symbols $b_m$ and of the spatial symbol $k$ obeys the 167 Maximum Likelihood (ML) criterion, which is formulated as

$$[\hat{m}_1, \ldots, \hat{m}_{N_a}, \hat{k}] = \arg \min_{s_m \in B} \left\{ \|y - \sqrt{\beta / N_a} H P s_m^l\|^2 \right\},$$  

where $B = \mathcal{C} \times \mathbb{C}^{N_a}$ is the joint search space of the super-169 symbol $s_m^l$. Alternatively, decoupled or separate detection may also be employed, which treats the detection of the conventional 171
modulated symbols $b_m$ and the spatial symbol $k$ separately. In this reduced-complexity variant, we have

$$\hat{k} = \arg \max_{l \in [1, |c|]} \left\{ \sum_{l=1}^{N_c} |c_l(\ell, l)|^2 \right\}, \quad (8)$$

$$\hat{m}_i = \arg \min_{n_i \in [1, M]} \left\{ |y_{\ell, i} - \sqrt{\beta/N_c} h_{\ell, i} p_{\ell, i} b_{n_i}|^2 \right\}, \quad (9)$$

where $h_{\ell, i}$ is the $i$th row of $H$ representing the channel between the $i$th RA and the transmitter, while $p_{\ell, i}$ is the $i$th column of $P$ representing the $i$th TPC vector. Thus, correct detection is declared, when we have $\hat{k} = k$ and $\hat{m}_i = m_{n_i}, \forall i$. 

Remarks: Note that the complexity of the ML detection of (7) is quite high, which is on the order determined by the super-alphabet $B$, hence obeying $O(|C|^M N_c)$. By contrast, the decoupled detection of (8) and (9) facilitates a substantially reduced complexity compared to that of (7). More explicitly, the complexity is imposed by detecting $N_c$ conventional modulated symbols, plus the complexity $(k_i)$ imposed by the comparisons invoked for non-coherently detecting the spatial symbol of (8), which may be written as $O(N_c M + k)$. Further discussions about the detection complexity of the decoupled detection of the GPSM scheme may be found in [17], where the main conclusion is that the complexity of the decoupled detection of the GPSM scheme is no higher than that of the conventional MIMO scheme corresponding to $N_a = N_r$.

III. PERFORMANCE ANALYSIS

We continue by investigating the DCMC capacity of our GPSM scheme, when the joint detection scheme of (7) is used and then quantify its achievable rate, when the realistic decoupled detection of (8) and (9) is employed. The achievable rate expression requires the theoretical BER/SER analysis of the GPSM scheme, which provides more insights into the inner nature of our GPSM scheme.\(^1\)

200 A. DCMC Capacity and Achievable Rate

Both Shannon’s channel capacity and its MIMO generalisation are maximized, when the input signal obeys a Gaussian distribution [22]. Our GPSM scheme is special in the sense that the spatial symbol conveys integer values constituted by the RA pattern index, which does not obey the shaping requirements of Gaussian signalling. This implies that the channel capacity of the GPSM scheme depends on a mixture of a continuous and 205 discrete input. Hence, for simplicity’s sake, we discuss the DCMC capacity and the achievable rate of our GPSM scheme in the context of discrete-input signalling for both the spatial 209 symbol and for the conventional modulated symbols mapped 211 to it.

1) DCMC Capacity: Upon recalling the received signal observed at the $N_r$ RAs expressed in (6), the conditional probability of receiving $y$ given that a $M = |C|^M N_c$-ary super-symbol $s_r \in B$ was transmitted over Rayleigh channel and subjected to it. The TPC of (3) is formulated as

$$p(y|s_r) = \frac{1}{\pi \sigma^2} \exp \left\{ -\frac{\|y - G s_r\|^2}{\sigma^2} \right\}, \quad (10)$$

where $G = \sqrt{\beta/N_c} H P$. The DCMC capacity of the ML-based joint detection of our GPSM scheme is given by [23]

$$C = \max_{p(s_1), \ldots, p(s_M)} \sum_{\tau=1}^{\infty} \int p(y, s_r) \log_2 \left( \frac{p(y|s_r)}{\sum_{\ell=1}^{M} p(y|s_\ell)} \right) dy, \quad (11)$$

which is maximized, when we have $p(s_r) = \frac{1}{|\mathcal{M}|}, \forall \tau$. Furthermore, we have

$$\log_2 \left( \frac{p(y|s_r)}{\sum_{\ell=1}^{M} p(y|s_\ell)} \right) = \log_2 \left( \frac{\sum_{\ell=1}^{M} p(y|s_\ell) p(s_\ell)}{\sum_{\ell=1}^{M} p(y|s_\ell) p(s_r)} \right) = -\log_2 \left( \frac{1}{|\mathcal{M}|} \sum_{\ell=1}^{M} p(y|s_\ell) p(s_r) \right) \exp(\Psi), \quad (12)$$

where substituting (10) into (12), the term $\Psi$ is expressed as

$$\Psi = -\frac{\|G(s_r - s) + w\|^2}{\sigma^2} + \|w\|^2. \quad (13)$$

Finally, by substituting (12) into (11) and exploiting that $p(s_r) = \frac{1}{|\mathcal{M}|}, \forall \tau$, we have

$$C = \log_2(|\mathcal{M}|) - \frac{1}{|\mathcal{M}|} \sum_{\tau=1}^{\infty} \mathbb{E}_{G, w} \left[ \log_2 \left( \sum_{\ell=1}^{M} \exp(\Psi) \right) \right]. \quad (14)$$

2) Achievable Rate: The above DCMC capacity expression implicitly relies on the ML-based joint detection of (7), which has a complexity on the order of $O(|\mathcal{M}|)$. When the reduced-complexity decoupled detection of (8) and (9) is employed, we estimate the achievable rate based on the mutual information $I(z; \hat{z})$ per bit measured for our GPSM scheme between the 230 input bits $z \in \{0, 1\}$ and the corresponding demodulated output 231 bits $\hat{z} \in \{0, 1\}$.

The mutual information per bit $I(z; \hat{z})$ is given for the Binary Symmetric Channel (BSC) by [22]:

$$I(z; \hat{z}) = H(z) - H(z|\hat{z}), \quad (15)$$

\(^1\)The reduced complexity receiver operates in a decoupled manner, which is beneficial in the scenario considered, where the spatial symbols and the conventionally modulated symbols are independent. However, this assumption may not be ideal, when correlations exist between the spatial symbols and the conventionally modulated symbols. In this case, an iterative detection exchanging extrinsic soft-information between the spatial symbols and conventionally modulated symbols may be invoked. Importantly, the iterations would exploit the beneficial effects of improving the soft-information by taking channel decoding into account as well for simultaneously exploiting the underlying correlations, which is reminiscent of the detection of correlated source. Further inspiration would be to beneficially map the symbols to both the spatial and to the conventional domain at the transmitter, so that the benefits of unequal protection could be exploited.

\(^2\)The Pair-wise Error Probability (PEP) analysis, relying on error events [21], was conducted in our previous contribution for the specific scenario of ML-based detection [17]. In this paper, our error probability analysis is dedicated to the low-complexity decoupled detection philosophy.
where $H(z) = - \sum_z P_z \log_2 P_z$ represents the entropy of the input bit $z$ and $P_z$ is the Probability Mass Function (PMF) of $z$. It is noted furthermore that we have $H(z) = 1$, when we adopt the common assumption of equal-probability bits, i.e., $P_z = 1/2$. On the other hand, the conditional entropy $H(z|\tilde{z})$ represents the average uncertainty about $z$ after observing $\tilde{z}$, which is given by:

$$H(z|\tilde{z}) = \sum_z P_z \left[ \sum_z P_{z|\tilde{z}} \log_2 P_{z|\tilde{z}} \right] = - e_x \log_2 e_x - (1 - e_x) \log_2 (1 - e_x),$$

where $e_x$ is the crossover probability. By substituting (16) into (15) and exploiting $H(z) = 1$ we have:

$$I(z; \tilde{z}) = 1 + e_x \log_2 e_x + (1 - e_x) \log_2 (1 - e_x).$$

Since the input bit in our GPSM scheme may be mapped either to a spatial symbol or to a conventional modulated symbol with a probability of $k_{ant}/k_{eff}$ and $N_a k_{mod}/k_{eff}$ respectively, the achievable rate becomes

$$R = k_{ant} I (e_x = e_{ant}^b) + N_a k_{mod} I (e_x = e_{mod}^b),$$

where $e_{ant}^b$ represents the BER of the spatial symbol, while $e_{mod}^b$ represents the BER of the conventional modulated symbol in the presence of spatial symbol errors due to the detection of (8).

252 B. Error Probability

253 1) The Expression of $e_{eff}^s$ and $e_{eff}^b$: Let us first let $e_{ant}^s$ represent the SER of the spatial symbol, while $e_{ant}^b$ represent the SER of the conventional modulated symbols in the presence of spatial symbol errors. Let further $N_{ant}^c$ and $N_{ant}^s$ represent the number of symbol errors in the spatial symbols and in the conventional modulated symbols, respectively. Then we have

$$e_{ant}^s = N_{ant}^c / N_s$$

and $e_{ant}^b = N_{ant}^c / N_{ant}^s$, where $N_s$ is the total number of GPSM symbols. Hence, the average SER $e_{eff}^s$ of our GPSM scheme is given by:

$$e_{eff}^s = \frac{(N_{ant}^c + N_{ant}^s)}{(1 + N_a) N_s}.$$

Similarly, the average BER $e_{eff}^b$ of our GPSM scheme may be written as:

$$e_{eff}^b = \frac{k_{ant} e_{ant}^b + N_a k_{mod} e_{mod}^b}{k_{eff}} \approx \frac{(\delta_{ant} e_{ant}^b + N_a e_{mod}^b)}{k_{eff}}.$$

254 where the second equation of (20) follows from the relation

$$e_{mod}^b \approx \frac{e_{mod}^b}{k_{mod}} e_{ant}^b \approx \frac{\delta_{ant} e_{ant}^b}{k_{ant}}.$$
301 3) Upper Bound of $\bar{e}^s_{\text{mod}}$: Considering a general case of $N_r$ as well as $N_a$ and assuming that the RA pattern $\mathcal{C}(k)$ was activated, after substituting (3) into (6), we have:

$$y_{vi} = \sqrt{\beta/N_a}b_{mi} + w_{vi}, \quad \forall v_i \in \mathcal{C}(k),$$

$$y_{vi} = w_{vi}, \quad \forall u_i \in \mathcal{C}(k),$$

304 where $\bar{C}(k)$ denotes the complementary set of the activated RA pattern $\mathcal{C}(k)$ in $\mathcal{C}$. Hence, we have the signal to noise ratio (SNR) given as

$$\gamma = \gamma_{vi} = \frac{\beta}{N_a\sigma^2} = \frac{\lambda}{2}, \quad \forall v_i$$

307 and for the remaining deactivated RAs in $\bar{C}(k)$, we have only 308 random noises of zero mean and variance of $\sigma^2$.

309 The SER $e^s_{\text{mod}}$ of the conventional modulated symbol $b_{mi} \in \mathcal{A}$ in the absence of spatial symbol errors may be upper 311 bounded by [24]:

$$e^s_{\text{mod}} < N_{\text{min}} \int_0^\infty Q(d_{\text{min}}\sqrt{\gamma/2}) f_\gamma(\gamma) d\gamma = e^{s,\text{ub}}_{\text{mod}},$$

312 where in general $f_\gamma(\gamma)$ has to be acquired by the empirical 313 histogram based method. When Lemma III.3 or Lemma III.4 is exploited, $f_\gamma(\gamma)$ is a scaled version of $f_\lambda(\lambda)$, i.e. we have

$$f_\gamma(\gamma) = 2f_\lambda(2\gamma).$$

315 Moreover, $d_{\text{min}}$ is the minimum Euclidean distance in the conventional modulated symbol constellation, $N_{\text{min}}$ is the average number of the nearest neighbours separated by $d_{\text{min}}$ in the constellation and $Q(\cdot)$ denotes the Gaussian Q-function.

320 When taking into account of the spatial symbol errors, we have Lemma III.5 for the upper bound of $\bar{e}^s_{\text{mod}}$:

322 Lemma III.5. (Proof in Appendix E): Given the k-th activated RA pattern, the SER of the conventional modulated symbols in the presence of spatial symbol errors can be upper bounded by:

$$\bar{e}^s_{\text{mod}} < \left(1 - e^{s,\text{ub}}_{\text{ant}}\right) e^{s,\text{ub}}_{\text{mod}} + e^{s,\text{ub}}_{\text{ant}} \sum_{\delta \neq k} N_{\delta} e^{s,\text{mod}}_{\delta} + N_a e^s_{\text{ant}} = \bar{e}^{s,\text{ub}}_{\text{mod}},$$

325 where $N_\ell$ and $N_{\text{ant}} = (N_a - N_\ell)$ represent the number of common and different RA between $\mathcal{C}(\ell)$ and $\mathcal{C}(k)$, respectively.

327 Mathematically we have $N_\ell = \sum_{i=1}^{N_a} \mathbb{I}[\mathcal{C}(\ell,i) \in \mathcal{C}(k)]$. Moreover, over $e^{s}_{\text{ant}} = (M - 1)/M$ is SER as a result of random guess.

329 4) Upper Bound of $e_{\text{eff}}^s$ and $e_{\text{eff}}^{b,ub}$: By substituting (25) and (33) into (19) and (20), we arrive at the upper bound of the 331 average symbol and bit error probability as

$$e_{\text{eff}}^s = \frac{e^{s,\text{ub}}_{\text{ant}} + N_a e^s_{\text{mod}}}{1 + N_a},$$

$$e_{\text{eff}}^{b,\text{ub}} = \frac{\delta_{\text{ant}} e^{s,\text{ub}}_{\text{ant}} + N_a e^s_{\text{mod}}}{k_{\text{eff}}},$$

Fig. 1. DCMC capacity versus the SNR of the CI TPC aided GPSM scheme based on the loose power-normalisation factor of (4) under $\{N_i, N_a\} = \{8, 4\}$ and employing QPSK, while having $N_a = \{1, 2, 3, 4\}$ activated RAs.

Similarly, by substituting (25) and (33) into (24), we obtain the lower bound of the achievable rate as

$$R_{\text{lb}} = k_{\text{ant}} I \left(\delta_{\text{ant}} e^{s,\text{ub}}_{\text{ant}}, \frac{e^s_{\text{mod}}}{k_{\text{mod}}}\right).$$

IV. NUMERICAL RESULTS

334 We now provide numerical results for characterizing both the 335 DCMC capacity of our GPSM scheme and for demonstrating 336 the accuracy of our analytical error probability results.

337 A. DCMC Capacity

338 1) Effect of the Number of Activated RAs: Fig. 1 characterises the DCMC capacity versus the SNR of the CI TPC aided GPSM scheme based on the loose power-normalisation factor of (4) under $\{N_i, N_a\} = \{8, 4\}$ and employing QPSK, while having $N_a = \{1, 2, 3, 4\}$ activated RAs. It can be observed in Fig. 1 that the larger $N_a$, the higher the capacity of our GPSM scheme. Importantly, both the GPSM scheme of $N_a = 3$ marked by the diamonds and its conventional MIMO 346 counterpart of $N_a = 4$ marked by the triangles attain the same 347 ultimate DCMC capacity of 8 bits/symbol at a sufficiently high 348 SNR, albeit the former exhibits a slightly higher capacity before 349 reaching the 8 bits/symbol value. Furthermore, the DCMC capacity of the conventional Maximal Eigen-Beamforming (Max 351 EB) scheme is also included as a benchmark under $\{N_i, N_a\} = \{32, 4\}$ and employing QPSK, which exhibits a higher DCMC 353 capacity at low SNRs, while only supporting 2 bits/symbol 354 at most.

355 We further investigate the attainable bandwidth efficiency by 356 replacing the SNR used in Fig. 1 by the SNR per bit in Fig. 2, 357 where we have $\text{SNR}_{\text{dB}} = \text{SNR}_{\text{dB}} - 10 \log_{10}(C/N_a)$. It can be seen from Fig. 2 that the lower $N_a$, the higher the 359 bandwidth efficiency attained in the low range of $\text{SNR}_{\text{dB}}$. Importantly, the achievable bandwidth efficiency of $N_a = 3$ is 361
Fig. 2. Bandwidth efficiency versus the SNR of CI TPC aided GPSM scheme with the loose power-normalisation factor of (4) under \{N_r, N_t\} = \{8, 4\} and employing QPSK, while having \(N_a = \{1, 2, 3, 4\}\) activated RAs.

Fig. 3. The effect of imperfect CSIT with \(\sigma_i = 0.4\) on the DCMC capacity versus the SNR of CI TPC aided GPSM scheme with the loose power-normalisation factor of (4) under \{N_t, N_r\} = \{8, 4\} and employing QPSK having \(N_a = \{1, 2, 3, 4\}\) activated RAs.

Fig. 4. The effect of antenna correlation with \(\rho_t = \rho_r = 0.3\) on the DCMC capacity versus the SNR of CI TPC aided GPSM scheme with the loose power-normalisation factor of (4) under \{N_t, N_r\} = \{8, 4\} and employing QPSK having \(N_a = \{1, 2, 3, 4\}\) activated RAs.

Fig. 5 characterises the DCMC capacity versus the SNR degradation is observed owing to antenna correlation. On the other hand, as seen in Fig. 4, roughly the same level of degradation is observed owing to antenna correlation.

3) Effect of Modulation Order and MIMO Configuration: Fig. 5 characterises the DCMC capacity versus the SNR of our CI TPC aided GPSM scheme with the loose power-normalisation factor of (4) under \{N_t, N_r\} = \{8, 4\} and employing QPSK having \(N_a = \{1, 2, 3, 4\}\) activated RAs. It can be seen in Fig. 6 that for a fixed 412 MIMO setting, the higher the value of \(N_a\), the higher the 413

of our CI TPC aided GPSM scheme relying on the loose 399 power-normalisation factor of (4) under \{N_t, N_r\} = \{8, 4\} and 400 employing various conventional modulation schemes having 401 \(N_a = \{1, 2, 3, 4\}\) activated RAs. It can be seen that the higher the 402 modulation order \(M\), the higher the achievable DCMC capacity. Furthermore, for a fixed modulation order \(M\), the higher the value of \(N_a\), the higher the achievable DCMC capacity becomes as a result of the information embedded in the spatial 406 symbol.

Fig. 6 characterises the DCMC capacity versus the SNR 408 for our CI TPC aided GPSM scheme for the loose power- 409 normalisation factor of (4) under different settings of \{N_t, N_r\} = \{410 1, 2\} and employing QPSK, while having \(N_a = 411 \{1, 2\}\) activated RAs.

2) Robustness to Impairments: Like in all TPC schemes, an important aspect related to GPSM is its resilience to CSI inaccuracies. In this paper, we let \(H = H_a + H_t\), where \(H_a\) represents the matrix hosting the average CSI, with each entry obeying the complex Gaussian distribution of \(h_a \sim CN(0, \sigma_a^2)\) and \(H_t\) is the instantaneous CSI error matrix obeying the complex Gaussian distribution of \(h_t \sim CN(0, \sigma_t^2)\), where we have \(\sigma_a^2 + \sigma_t^2 = 1\). As a result, only \(H_a\) is available at the transmitter for pre-processing.

Another typical impairment is antenna correlation. The correlated MIMO channel is modelled by the widely-used Kronecker model, which is written as \(H = (R_t^{1/2})^T G (R_r^{1/2})\), with \(G\) representing the original MIMO channel imposing no correlation, while \(R_t\) and \(R_r\) represents the correlations at the transmitter and receiver side, respectively, with the correlation entries given by \(R_t(i, j) = \rho_t^{i-j}\) and \(R_r(i, j) = \rho_r^{i-j}\). Figs. 3 and 4 characterise the effect of imperfect CSIT associated with \(\sigma_i = 0.4\) and of antenna correlation of \(\rho_t = 0.3\) on the attainable DCMC capacity versus the SNR for our CI TPC aided GPSM scheme with the loose power-normalisation factor of (4), respectively, under \{N_t, N_r\} = \{8, 4\} and employing QPSK having \(N_a = \{1, 2, 3, 4\}\) activated RAs. It can be seen that as expected, both impairments result into a degraded DCMC capacity. Observe in Fig. 3 for imperfect CSIT that the degradation of the conventional MIMO associated with \(N_a = 4\) and marked by the triangle is larger than that of our GPSM scheme corresponding \(N_a = \{1, 2, 3\}\). On the other hand, as seen in Fig. 4, roughly the same level of degradation is observed owing to antenna correlation.

3) Effect of Modulation Order and MIMO Configuration: Fig. 5 characterises the DCMC capacity versus the SNR consistently and significantly higher than that achieved by \(N_a = 4\), before they both converge to 8 bits/symbol/Hz at their maximum. Overall, there is always a beneficial configuration for our GPSM scheme that offers the same bandwidth efficiency as that of its conventional MIMO counterpart, which is achieved at a lower SNR per bit.
Fig. 5. DCMC capacity versus the SNR of our CI TPC aided GPSM scheme relying on the loose power-normalisation factor of (4) under \( \{N_t, N_r\} = \{8, 4\} \) and employing various conventional modulation schemes having \( N_a = \{1, 2\} \) activated RAs.

Fig. 6. DCMC capacity versus the SNR for our CI TPC aided GPSM scheme for the loose power-normalisation factor of (4) under different settings of \( \{N_t, N_r\} \) with \( N_t/N_r = 2 \) and employing QPSK, while having \( N_a = \{1, 2\} \) activated RAs.

DCMC capacity becomes. Importantly, for a fixed \( N_a \), the larger the size of the MIMO antenna configuration, the higher DCMC capacity.

B. Achievable Rate

1) Error Probability: Figs. 7–10 characterize the GPSM scheme’s SER as well as the BER under both the loose power-normalisation factor of (4) and the stringent power-normalisation factor of (5) for \( \{N_t, N_r\} = \{16, 8\} \) and employing QPSK, respectively. From Figs. 7–10, we recorded the curves from left to right corresponding to \( N_a = \{1, 2, 4, 6\} \). For reasons of space-economy and to avoid crowded figures, our results for \( N_a = \{3, 5, 7\} \) were not shown here, but they obey the same trends.

It can be seen from Figs. 7 and 9 that our analytical SER results of (34) form tight upper bounds for the empirical simulation results. Hence they are explicitly referred to as ‘tight upper bound’ in both figures. Additionally, a loose upper bound of the GPSM scheme’s SER is also included, which may be written as

\[
e_{s,\text{eff}} = 1 - (1 - e_{\text{ant}}^{s,\text{ub}}) (1 - e_{\text{mod}}^{s,\text{ub}}).
\]

(37)

Note that in this loose upper bound expression, \( e_{\text{mod}}^{s,\text{ub}} \) of (32) is required rather than \( e_{\text{mod}}^{s,\text{ub}} \) of (33). This expression implicitly assumes that the detection of (8) and (9) are independent. However, the first-step detection of (8) significantly affects the second-step detection of (9). Hence, the loose upper bound shown by the dash-dot line is only tight for \( N_a = 1 \) and becomes much looser upon increasing \( N_a \), when compared to the tight upper bound of (34).
Similarly, when the GPSM scheme’s BER is considered in Figs. 8 and 10, our the analytical results of (35) again form tight upper bounds for the empirical results.

2) Separability: To access the inner nature of first-step detection of (8), Fig. 11 reveals the separability between the activated RAs and deactivated RAs in our GPSM scheme, where the PDF of (44) and (45) were recorded both for SNR $=-5$ dB (left subplot) and for SNR $=0$ dB (right subplot) respectively for the same snapshot of MIMO channel realisation with the aid of CI TPC and the loose power-normalisation factor of (4) under $\{N_t, N_r\} = \{16, 8\}$ and employing QPSK. By comparing the left subplot to the right subplot, it becomes clear that the higher the SNR, the better the separability between the activated and the deactivated RAs, since the mean of the solid curves representing (44) move further apart from that of the dashed curve representing (45). Furthermore, as expected, the lower $N_a$, the better the separability becomes, as demonstrated in both subplots of Fig. 11.

3) Comparison: Finally, Fig. 12 characterizes the comparison between the DCMC capacity (14) of our GPSM scheme relying implicitly on the ML-based joint detection and its lower bound of the achievable rate in (36) relying on the low-complexity decoupled detection, where we use CI TPC with the loose power-normalisation factor of (4) under $\{N_t, N_r\} = \{16, 8\}$ and employing QPSK having $N_a = \{1, 2, 3\}$.

It is clear that the DCMC capacity is higher than the achievable rate for each $N_a$ considered, although both of them converge to the same value, when the SNR is sufficiently high. Noticeably, the discrepancy between the two quantities before their convergence is wider, when $N_a$ is higher. This is because the higher $N_a$, the lower the achievable rate at low SNRs.
which is shown by comparing the solid curves. This echoes our observations of Fig. 11, namely that a higher $N_a$ leads to a reduced separability and consequently both to a higher overall error probability and to a lower achievable rate. In fact, the achievable rate becomes especially insightful after being compared to the DCMC capacity, where we may tell how a realistic decoupled detection performs and how far its performance is from the DCMC capacity.

V. CONCLUSION

In this paper, we introduced the concept of our GPSM scheme and carried out its theoretical analysis in terms of both its DCMC capacity as well as its achievable rate relying on our analytical upper bound of the SER and the BER expressions, when a low-complexity decoupled detector is employed. Our numerical results demonstrate that the upper bound introduced is tight and the DCMC capacity analysis indicates that our GPSM scheme constitutes a flexible MIMO arrangement. Our future work will consider a range of other low-complexity MIMO schemes, such as the receive antenna selection and the classic SM, in the context of large-scale MIMOs.

Furthermore, the insights of our error probability and capacity analysis are multi-folds:

- It can be seen that there is a gap between the DCMC capacity relying on ML detection and the achievable rate of decoupled detection. Thus, a novel detection method is desired for closing this gap and for striking a better trade-off between the performance attained and the complexity imposed.

- The error probability derived serves as a tight upper bound of the alphabet represented by 0 and 1, respectively. We may thereafter refer to the zero-alphabet (one-alphabet). Hence, the number of pair-wise bit errors $\epsilon_1$ in the one-alphabet is simply $\epsilon_1 = \epsilon_0 + A$, where $A = 2^{k_{ant}t}$ accounts for the difference in the first preceding bit. Hence the total number of pair-wise bit errors is $\epsilon = 2\epsilon_0 + 2^{k_{ant}t}$. Taking into account an equal probability of $1/(2^{k_{ant}t} - 1)$ for each possible spatial symbol, we arrive at the correction factor given by $\delta_{k_{ant}} = (2\epsilon_0 + 2^{k_{ant}t} - 1)$. Since $\epsilon_0$ represents the total number of pair-wise bit errors $\epsilon$ corresponding to case of $(k_{ant}t) - 1$ bits of information, we have $\epsilon_0 = (2^{k_{ant}t} - 1)\delta_{k_{ant}} - 1$. Hence the resultant expression of the correction factor may be calculated recursively according to (23) after some further manipulations.\footnote{By assuming equal-probability erroneously detected patterns, a spatial symbol may be mistakenly detected as any of the other spatial symbols with equal probability. Let us now give an example for highlighting the rationale of introducing the correction factor. For example, spatial symbol ‘0’ carrying bits $[0,0]$ was transmitted, it would result into a one-bit difference when the spatial symbol ‘1’ carrying $[0,1]$ or ‘2’ carrying $[1,0]$ was erroneously detected. However, it would result into a two-bits difference when spatial symbol ‘3’ carrying $[1,1]$ was erroneously detected. This corresponds to four bit errors in total for these three cases, thus a correction factor of $4/3$ is needed when converting the symbol error ratio to bit error ratio.}

\begin{app}
\section{APPENDIX A
PROOF OF LEMMA III.1

Let $A_{k_{ant}}$ denote the alphabet of the spatial symbol having $k_{ant}$ bits of information. Then the cardinality of the alphabet $A_{k_{ant}}$ is twice higher compared to that of $A_{k_{ant}-1}$. Thus, $A_{k_{ant}}$ may be constructed by two sub-alphabets of $A_{k_{ant}-1}$, represented by 0 and 1, respectively. We may thereafter refer to the alphabet of $A_{k_{ant}-1}$ preceded by the above-mentioned with 0 (1) as zero-alphabet (one-alphabet).

Assuming that the spatial symbol representing $k_{ant}$ zeros was transmitted, we may then calculate the total number of pair-wise bit errors $\epsilon_0$ in the above zero-alphabet. Hence, the number of pair-wise bit errors $\epsilon_1$ in the one-alphabet is simply $\epsilon_1 = \epsilon_0 + A$, where $A = 2^{k_{ant}t}$ accounts for the difference in the first preceding bit. Hence the total number of pair-wise bit errors is $\epsilon = 2\epsilon_0 + 2^{k_{ant}t}$. Taking into account an equal probability of $1/(2^{k_{ant}t} - 1)$ for each possible spatial symbol, we arrive at the correction factor given by $\delta_{k_{ant}} = (2\epsilon_0 + 2^{k_{ant}t} - 1)$. Since $\epsilon_0$ represents the total number of pair-wise bit errors $\epsilon$ corresponding to case of $(k_{ant}t) - 1$ bits of information, we have $\epsilon_0 = (2^{k_{ant}t} - 1)\delta_{k_{ant}} - 1$. Hence the resultant expression of the correction factor may be calculated recursively according to (23) after some further manipulations.\footnote{By assuming equal-probability erroneously detected patterns, a spatial symbol may be mistakenly detected as any of the other spatial symbols with equal probability. Let us now give an example for highlighting the rationale of introducing the correction factor. For example, spatial symbol ‘0’ carrying bits $[0,0]$ was transmitted, it would result into a one-bit difference when the spatial symbol ‘1’ carrying $[0,1]$ or ‘2’ carrying $[1,0]$ was erroneously detected. However, it would result into a two-bits difference when spatial symbol ‘3’ carrying $[1,1]$ was erroneously detected. This corresponds to four bit errors in total for these three cases, thus a correction factor of $4/3$ is needed when converting the symbol error ratio to bit error ratio.}

\end{app}

\begin{proof}

Considering a general case of $N_a$, as well as $N_a$ and assuming that the RA pattern $C(k)$ was activated, after substituting (3) into (6), we have:

\begin{equation}
y_v_i \sim \mathcal{N}(\sqrt{\beta/N_a} R(b_{m_i}), \sigma_0^2),
\end{equation}

where $C(k)$ denotes the complementary set of the activated RA pattern $C(k)$ in $C$. Furthermore, upon introducing $\sigma_0^2 = \sigma^2/2$, we have:

\begin{align}
|y_v_i|^2 &= \mathcal{R}(y_v_i)^2 + \mathcal{I}(y_v_i)^2 \\
&\sim \mathcal{N}(\sqrt{\beta/N_a} R(b_{m_i}), \sigma_0^2) + \mathcal{N}(0, \sigma_0^2),
\end{align}

\begin{align}
|y_u_i|^2 &= \mathcal{R}(u_{m_i})^2 + \mathcal{I}(u_{m_i})^2 \\
&\sim \mathcal{N}(0, \sigma_0^2) + \mathcal{N}(0, \sigma_0^2),
\end{align}

where $\mathcal{R}(\cdot)$ and $\mathcal{I}(\cdot)$ represent the real and imaginary operators, respectively. As a result, by normalization with respect to $\sigma_0^2$, we have the following observations:

\begin{align}
|y_v_i|^2 &\sim \chi^2(g; \lambda_v), \quad \forall v_i \in C(k), \\
|y_u_i|^2 &\sim \chi^2(g), \quad \forall u_i \in \bar{C}(k),
\end{align}

where the non-centrality is given by $\lambda_v = \beta |b_{m_i}|^2/|N_a\sigma_0^2|$. Exploiting the fact that $\mathbb{E}[|b_{m_i}|^2] = 1$, $\forall i$ (or $|b_{m_i}|^2 = 1$, $\forall i$ for 551 PSK modulation), we have $\lambda = \lambda_v$, $\forall v_i$. Note that $\lambda$ is also a $552$ random variable obeying the distribution of $f_\lambda$. The non-centrality of the spatial symbols given the non-centrality $\lambda$, $557$

\end{proof}

\begin{section}{APPENDIX B
PROOF OF LEMMA III.2

Consider a general case of $N_a$, as well as $N_a$ and assuming that the RA pattern $C(k)$ was activated, after substituting (3) into (6), we have:

\begin{equation}
y_v_i \sim \mathcal{N}(\sqrt{\beta/N_a} R(b_{m_i}), \sigma_0^2),
\end{equation}

where $C(k)$ denotes the complementary set of the activated RA pattern $C(k)$ in $C$. Furthermore, upon introducing $\sigma_0^2 = \sigma^2/2$, we have:

\begin{align}
|y_v_i|^2 &= \mathcal{R}(y_v_i)^2 + \mathcal{I}(y_v_i)^2 \\
&\sim \mathcal{N}(\sqrt{\beta/N_a} R(b_{m_i}), \sigma_0^2) + \mathcal{N}(0, \sigma_0^2),
\end{align}

\begin{align}
|y_u_i|^2 &= \mathcal{R}(u_{m_i})^2 + \mathcal{I}(u_{m_i})^2 \\
&\sim \mathcal{N}(0, \sigma_0^2) + \mathcal{N}(0, \sigma_0^2),
\end{align}

where $\mathcal{R}(\cdot)$ and $\mathcal{I}(\cdot)$ represent the real and imaginary operators, respectively. As a result, by normalization with respect to $\sigma_0^2$, we have the following observations:

\begin{align}
|y_v_i|^2 &\sim \chi^2(g; \lambda_v), \quad \forall v_i \in C(k), \\
|y_u_i|^2 &\sim \chi^2(g), \quad \forall u_i \in \bar{C}(k),
\end{align}

where the non-centrality is given by $\lambda_v = \beta |b_{m_i}|^2/|N_a\sigma_0^2|$. Exploiting the fact that $\mathbb{E}[|b_{m_i}|^2] = 1$, $\forall i$ (or $|b_{m_i}|^2 = 1$, $\forall i$ for 551 PSK modulation), we have $\lambda = \lambda_v$, $\forall v_i$. Note that $\lambda$ is also a $552$ random variable obeying the distribution of $f_\lambda$. The non-centrality of the spatial symbols given the non-centrality $\lambda$, $557$

\end{section}
558 when the RA pattern $C(k)$ was activated may be lower bounded
559 as in (46). (See equation at bottom of page) More explicitly,
560 • equation (a) serves as the lower bound, since it sets the
561 most strict condition for the correct detection, when each
562 metric $y_{u_j}$ of the inactivated RA indices in $C(k)$ is lower
563 than each metric $y_{v_i}$ of the activated RA indices in $C(k)$.
564 Note that, equality holds when $N_a = 1$;
565 • equation (b) follows from the fact that the $N_a$ random
566 variables $|y_{u_j}|^2$ are independent of each other;
567 • equation (c) follows from the fact that the $(N_r - N_a)$
568 random variables $|y_{v_i}|^2$ are independent and equation (d)
569 follows from the fact that the $N_a$ independent variables of
570 $|y_{v_i}|^2$ and the $(N_r - N_a)$ independent variables of $|y_{u_j}|^2$
571 are both identically distributed.
572 As a result, after averaging over the distribution of $f_\lambda(\lambda)$, the
573 analytical SER $c_{ant}^*$ of the spatial symbol in our GPSM scheme
574 may be upper bounded as in (25). In general, the expression
575 of $f_\lambda(\lambda)$ can be acquired with the aid of the empirical
576 histogram based method, while in case the loose/stringent
577 normalisation factor of (4)/(5) is used, the analytical expression
578 for $f_\lambda(\lambda)$ is given in Lemma III.3/Lemma III.4.

579 APPENDIX C
580 PROOF OF LEMMA III.3
581
582 Upon expanding the expression of $\lambda$ in (26) by taking into
583 account (4), we have:
584
$$\lambda = \frac{\beta_1}{N_a\sigma_0^2} = \frac{N_r}{N_a\sigma_0^2 \text{Tr}[HH^H]}. \quad (47)$$

585 Consider first the distribution of $\text{Tr}[(HH^H)^{-1}]$ and let $W = HH^H$. Since the entries of $H$ are i.i.d. zero-mean unit-
586 variance complex Gaussian random variables, $W$ obeys a 587 complex Wishart distribution. Hence the joint PDF of its eigen- 586
587 values $\{\lambda_{W_i}\}_{i=1}^{N_r}$ is given by [26], [27]
588
$$f_W(\{\lambda_{W_i}\}_{i=1}^{N_r}) = \frac{K^{-1}}{N_r!} \prod_{i} e^{-\lambda_{W_i}} \lambda_{W_i}^{N_r-N_i} \prod_{i<j} (\lambda_{W_i} - \lambda_{W_j})^2, \quad (48)$$

where $K$ is a normalising factor. Thus for its inverse $U = W^{-1}$, we have
589
$$f_U(\{\lambda_{W_i}\}_{i=1}^{N_r}) = \prod_{i} \lambda_{W_i}^{-2} f_W \left(\left\{\frac{1}{\lambda_{W_i}}\right\}_{i=1}^{N_r}\right). \quad (49)$$

Furthermore, since $\text{Tr}[U] = \sum \lambda_{U_i}$, where $\{\lambda_{U_i}\}_{i=1}^{N_r}$ is the 590
591 eigenvalues of $U$, we have the CDF of $\text{Tr}[U]$ given by (50), 591
592 where $T_1 = T$ and $t_j = 1/T$, while $\forall j > 1$
593
\begin{align*}
T_j & = T - \sum_{i=1}^{j-1} \lambda_{U_i}, \quad T - \sum_{i=1}^{j-1} \lambda_{U_i}. \quad (50)
\end{align*}

\begin{align*}
\Delta & \geq \int_{0}^{\infty} P \left( |y_{u_1}|^2 < g_{v_1}, \ldots, |y_{u_{N_r-N_{a}}}|^2 < g_{v_1}, \ldots, |y_{u_{N_r-N_{a}}}|^2 < g_{v_{N_{a}}} \right) \times \\
& \times P \left( |y_{v_1}|^2 = g_{v_1}, \ldots, |y_{u_{N_{a}}}|^2 = g_{v_{N_{a}}} | \lambda_{v_1}, \ldots, \lambda_{v_{N_{a}}} \right) \; dg_{v_1}\cdots dg_{v_{N_{a}}} \\
& = \prod_{i=1}^{N_r} \int_{0}^{\infty} P \left( |y_{u_i}|^2 < g_{v_i}, \ldots, |y_{u_{N_r-N_{a}}}|^2 < g_{v_i} \right) P \left( |y_{v_i}|^2 = g_{v_i} | \lambda_{v_i} \right) \; dg_{v_i} \\
& = \prod_{i=1}^{N_r} \prod_{u_j \in C(k)} P \left( |y_{u_j}|^2 < g_{v_i} \right) P \left( |y_{v_i}|^2 = g_{v_i} | \lambda_{v_i} \right) \; dg_{v_i} \\
& = \left\{ \int_{0}^{\infty} f_{\lambda_2} (g) \lambda_{N_r-N_{a}} \lambda_{v_i} N_{a} \, dg \right\}
\end{align*}

(46)

\begin{align*}
F_{\text{Tr}[U]}(T) & = \int_{0}^{T_1} \int_{0}^{T_2} \cdots \int_{0}^{T_{N_r}} f_U(\{\lambda_{U_i}\}_{i=1}^{N_r}) \; d\lambda_{U_{N_r}} \cdots d\lambda_{U_1} = \int_{t_1}^{\infty} \int_{t_2}^{\infty} \cdots \int_{t_{N_r}}^{\infty} f_W(\{\lambda_{U_i}\}_{i=1}^{N_r}) \; d\lambda_{U_{N_r}}^{-1} \cdots d\lambda_{U_1}^{-1}.
\end{align*}

(50)
Since the entries of $\mathbf{H}$ are i.i.d. zero-mean unit-variance complex Gaussian random variables, $\mathbf{HH}^H$ obeys a complex Wishart distribution with $N_r$ dimensions and $2N_t$ degrees of freedom, where we have:

$$\mathbf{HH}^H \sim \mathcal{C}\mathcal{W}(\Sigma, N_r, 2N_t),$$

where $\Sigma = (1/2)I_{N_r}$ being the variance. By exploiting proposition 8.9 from [28] and letting $\lambda_0 = [s^H (\mathbf{HH}^H)^{-1} s]^{-1}$, we have:

$$\lambda_0 \sim \mathcal{C}\mathcal{W}\left([s^H \Sigma^{-1} s]^{-1}, 1, 2(N_t - N_r + 1)\right),$$

According to [28], the above one-dimensional complex-valued Wishart distribution is actually a chi-square distribution with $2(N_t - N_r + 1)$ degrees of freedom and scaling parameter of $[s^H \Sigma^{-1} s]^{-1} = 1/2N_a$. Thus, the PDF of $\lambda_0$ may be explicitly written as:

$$f_{\lambda_0}(\lambda_0) = f_{\chi^2}(2N_a\lambda_0; 2(N_t - N_r + 1))$$

$$= 2N_a e^{-\lambda_0 N_a} (2N_a\lambda_0)^{N_t - N_r - 1} / (N_t - N_r)!,$$

Finally, since $\lambda_0 = \sigma_0^2 / \lambda$, we have $f_{\lambda}(\lambda) = \sigma_0^2 f_{\lambda_0}(\sigma_0^2 / \lambda)$, which is (28).

**APPENDIX E**

**PROOF OF LEMMA III.5**

20 The SER of $\tilde{e}_{\text{mod}}^a$ is constituted by the SER of $e_{\text{mod}}^a$, when the detection of the spatial symbol is correct having a probability of $(1 - e_{\text{ant}}^a)$, plus the SER, when detection of the spatial symbol is erroneous having a probability of $e_{\text{ant}}^a$, which is expressed as:

$$\tilde{e}_{\text{mod}}^a = (1 - e_{\text{ant}}^a) e_{\text{mod}}^a + e_{\text{ant}}^a \sum_{\ell \neq k} P_{k \rightarrow \ell} \frac{N_c e_{\text{mod}}^a + N_d e_{\text{ant}}^a}{N_a} \lambda^a.$$


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AUTHOR QUERIES

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AQ1 = Please be informed that the capital letters were removed from the terms “multiple input multiple output,” “generalised pre-coded aided spatial modulation,” “symbol error ratio,” “bit error ratio,” “discrete-input continuous-output memoryless channel,” and “signal to noise ratio” in the Abstract per IEEE style and also in other occurrences of these terms in lines 88 to 91 and 305 for the sake of consistency. Please check if it is correct.

AQ2 = Please provide keywords.

AQ3 = Please check changes made in first footnote and the addition of an Acknowledgment Section.

AQ4 = Please check if “30 journals” should be “30 papers” instead.

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