# Refinement in Practise

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### Abstraction

- Abstraction can be viewed as a process of simplifying our understanding of a system.
- The simplification should
  - focus on the intended purpose of the system
  - ignore details of how that purpose is achieved.
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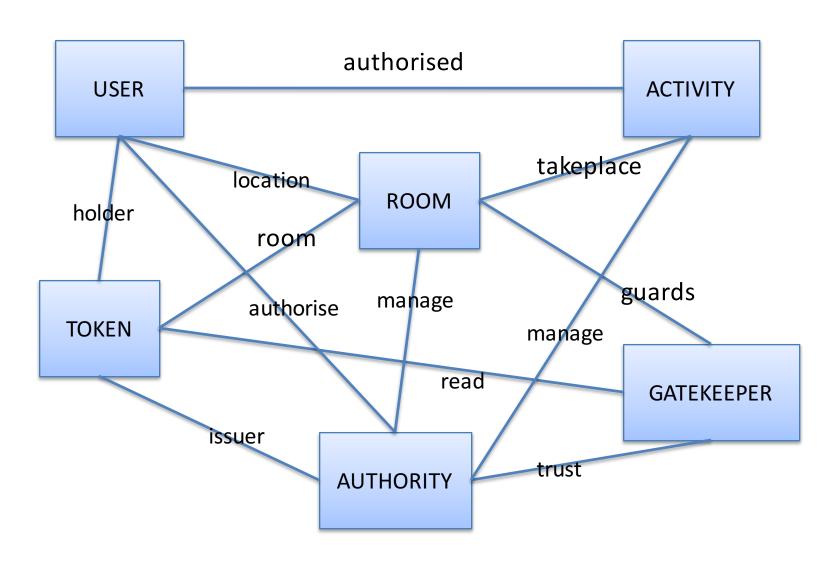
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- ▶ If the purpose is to provide some service, then
  - model what a system does from the perspective of the service users
  - 'users' might be computing agents as well as humans.
- ▶ If the purpose is to control, monitor or protect some phenomenon, then
  - the abstraction should focus on those phenomenon
  - in what way should they be controlled or protected?



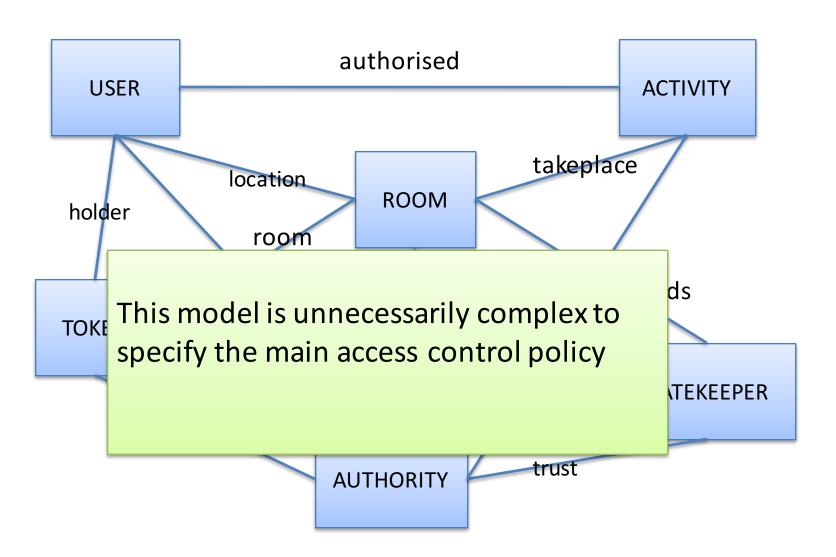
# Access control system

- Users are authorised to engage in activities
- User authorisation may be added or revoked
- Activities take place in rooms
- Users gain access to a room using a one-time token provided they have authority to engage in the room activities
- Tokens are issued by a central authority
- Tokens are time stamped
- A room gateway allows access with a token provided the token is valid

# Class diagram



# Class diagram



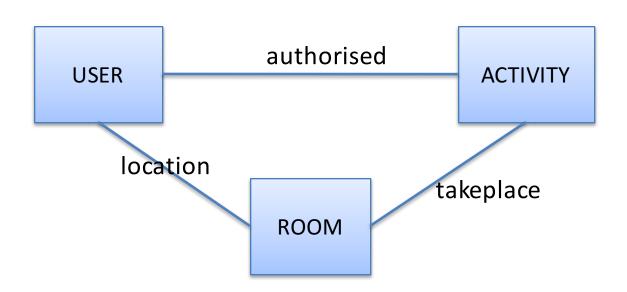
# Extracting the essence

 Access Control Policy: Users may be in a room only if they are authorised to engage in all activities that may take place in that room

 To express this we only require Users, Rooms, Activities and relationships between them

Abstraction: focus on key entities in the problem domain

# Diagrammatic representation of an abstract model



# Variables and invariants of Event-B model

# Variables of Event-B model

```
    @inv1 authorised ∈ User ↔ Activity // relation
    @inv2 takeplace ∈ Room ↔ Activity // relation
    @inv3 location ∈ User → Room // partial function
```

# Access control invariant:

**if** user u is in room r, **then** u must be authorised to engaged in all activities that can take place in  $\underline{r}$ 

```
@inv4 \forall u,r. u \in dom(location) \land location(u) = r \Rightarrow takeplace[r] \subseteq authorised[u]
```

- ► **Refinement** is a process of enriching or modifying a model in order to
  - augment the functionality being modelled, or
  - explain how some purpose is achieved

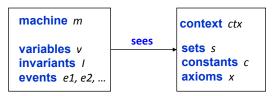
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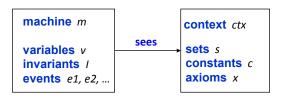
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  - ▶ We use proof to verify the consistency of a refinement step
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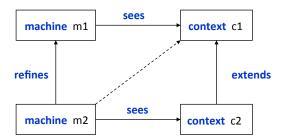
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- ► Abstraction and refinement together should allow us to manage system complexity in the design process

# Modelling Components and Refinement



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## Extension Refinement in Event-B

A refined machine has the following form:

```
machine M2 refines M1 variables ... invariants ... events...
```

Extension refinement can be used to extend or add new features to a model.

- Add variables and invariants
- Extend existing events to act on additional variables
- Add new events to act on additional variables

All events must maintain the new invariants.

# **Event Extension**

Event E1 of M1 may be extended by E2 in M2:

```
E2 extends E1 =
any < additional parameters > where
< additional guards >
then
< additional actions >
end
```

#### **Event Extension**

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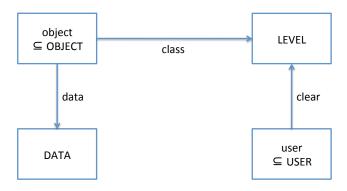
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E2 extends E1 =
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end
```

Extending an event means

- adding parameters
- adding guards
- adding actions

Extension example: add ownership to secure database

# Class diagram for secure database



# Types and variables for Secure DB

```
context c1
sets OBJECT DATA USER
constants LEVEL
axioms IEVEI = 1..10
machine SecureDB1
sees c1
variables object, user, data, class, clear
invariants
                   object ⊂ OBJECT
                   user \subseteq USER
                   data \in object \rightarrow DATA
                   class \in object \rightarrow LEVEL
                   clear \in user \rightarrow LEVEL
```

# Adding object ownership

Extend the database specification so that each object has an owner.

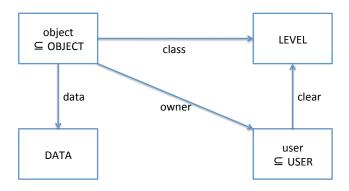
The clearance associated with that owner must be at least as high as the classification of the object.

Only the owner of an object is allowed to delete it.

What additional variables are required?

What events are affected?

# Class diagram with ownership



```
machine SecureDB2

refines SecureDB1

variables object, user, data, class, clear, owner invariants

owner \in object \rightarrow user
```

Note we must list all the variables: those from M1 that we wish to retain as well as new ones

Here owner is a new variable.

We do not repeat invariants of M1 in M2

# Adding users

```
\begin{array}{ccc} \textit{AddUser} & \triangleq \\ & \textbf{any} \ \textit{u}, \textit{c} \ \textbf{where} \\ & \textit{u} \in \textit{USER} \\ & \textit{u} \notin \textit{user} \\ & \textit{c} \in \textit{LEVEL} \\ & \textbf{then} \\ & \textit{user} := \textit{user} \cup \{\textit{u}\} \\ & \textit{clear}(\textit{u}) := \textit{c} \\ & \textbf{end} \end{array}
```

Do we need to modify this?

# Adding objects

```
any o, d, c where
      o ∈ OBJECT
      o ∉ object
      d \in DATA
      c \in LEVEL
   then
      object := object \cup \{o\}
      data(o) := d
      class(o) := c
   end
```

Do we need to modify this?

# **Event Extension**

```
\begin{array}{cccc} \textit{AddObject} & \textbf{extends} & \textit{AddObject} & \hat{=} \\ & \textbf{any} & u & \textbf{where} \\ & u \in \textit{user} \\ & \textit{clear}(u) \geq \textit{class}(o) \\ & \textbf{then} \\ & \textit{owner}(o) \ := \ u \\ & \textbf{end} \end{array}
```

# This is equivalent to

```
AddObject refines AddObject \hat{=}
    any o, d, c, u where
       o \in OBJECT
       o ∉ object
       d \in DATA
       c \in LEVEL
       u \in \mathit{user}
       clear(u) > class(o)
    then
       object := object \cup \{o\}
       data(o) := d
       class(o) := c
       owner(o) := u
    end
```

# Other events to consider

- Read
- Write
- ChangeClass
- ChangeClear
- RemoveUser, RemoveObject

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- Verification of 2, 3 and 4 requires gluing invariants that link abstract and concrete variables.

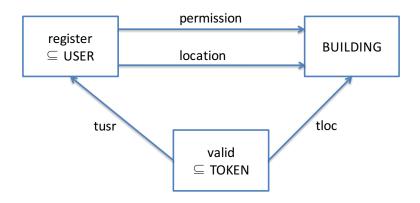


- Extension example: add ownership to secure database
- ► Extension with Guard Modification example: add tokens to buildings access system
- ▶ Variable replace example: simple data sampling system

## Abstract model of building access



# Refine this by introducing a token mechanism



#### Refinement of access control with tokens

```
EnterBuilding \hat{=} any u, b where
                             u \notin dom(location)
                             u \mapsto b \in permission
                         then
                             location(u) := b
                         end
RefinedEnterBuilding \hat{=} any u, b, t where
                                  u \notin dom(location)
                                  t \in valid
                                 tusr(t) = u
                                  tloc(t) = b
                              then
                                  location(u) := b
                              end
```

#### **GRD Proof Obligations**

We need to prove that the guard of a refined event is not weaker than the guard of the abstract event.

E.g., the refined enter event should not weaken the conditions under which a user may enter a room.

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Assume: guard(RefinedEnter) + invariants

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GRD Proof obligation:

Assume: guard(RefinedEnter) + invariants

Prove: guard(AbstractEnter)

For the access control refinement, we need this invariant:

 $\forall t \cdot t \in valid \implies tusr(t) \mapsto tloc(t) \in permission$ 

## Simple data sampling system

machine MaxSet1

```
variables samples
invariants samples \subseteq \mathbb{N}
initialisation samples := \{0\}
events
             Add = any \times where
                             x \in \mathbb{N}
                         then
                             samples := samples \cup \{x\}
                         end
             GetMax \triangleq any result where
                                 result = max(samples)
                             end
```

# Refine to a more optimal design

```
machine MaxSet2
refines MaxSet1
variables m we only need to store the maximum so far
invariants m \in \mathbb{N}
initialisation m := 0
events
               Add \triangleq any \times where
                             x \in \mathbb{N}
                          then
                              m := max(\{m, x\})
                          end
                GetMax = any result where
                                   result = m
```

end

## Gluing invariant

What is the relationship between m and samples?

#### Gluing invariant

What is the relationship between *m* and *samples*?

```
machine MaxSet2

refines MaxSet1

variables m

invariants m = max(samples)

events...
```

This is called a gluing invariant: it specifies the relationship between the abstract and refined variables.

```
Abstract Add: samples := samples \cup \{x\}
```

Refined Add:  $m := max(\{m, x\})$ 

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Assume: m = max(samples)

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This is valid since:

$$max(s \cup \{x\}) =$$

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Refined Add:  $m := max(\{m, x\})$ 

Assume: m = max(samples)

Prove:  $max(\{m,x\}) = max(samples \cup \{x\})$ 

This is valid since:

$$max(s \cup \{x\}) = max(\{max(s), x\})$$

## Closing Messages

- Role of formal modelling:
  - increase understanding
  - decrease errors
- Role of refinement:
  - manage complexity through multiple levels of abstraction
- Role of verification:
  - improve quality of models (consistency, invariants)
- ► Role of tools:
  - make verification as automatic as possible, pin-pointing errors