

Asynchronous updates can promote the evolution of cooperation on multiplex networks

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Abstract

We study the importance to the frequency of cooperation of the choice of updating strategies in a game played asynchronously or synchronously across layers in a multiplex network. Updating asynchronously in the public goods game leads to higher frequencies of cooperation compared to synchronous updates. How large this effect is depends on the sensitivity of the game dynamics to changes in the number of cooperators surrounding a player, with the largest effect observed when players payoffs are small. The discovery of this effect enhances understanding of cooperation on multiplex networks, and demonstrates a new way to maintain cooperation in these systems.

Keywords:

Cooperation, multiplex networks, public goods game

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1. Introduction

In many social systems a social dilemma exists where what is good for the individual is not necessarily good for the group, and so explaining why in both laboratory [1, 2] and field [3, 4] experiments participants are found to cooperate (act in the interests of the group) more frequently than would be expected for purely rational self-interested players has proved a challenge to evolutionary biology and sociology. The prevalence of cooperation between unrelated members of a population in both biological and human social systems has been the subject of a long history of study, with a number of key insights over the past few decades [5].

One of the central tools used to understand the resolution of social dilemmas and the evolution of cooperation is game theory [6]. Here each player can choose from a number of strategies with a payoff, dependent on their own strategy and those chosen by their opponents, indicating to each player how well they are performing. Playing the prisoner's dilemma (PD) on a lattice network (where the network defines those that each player plays against) cooperation has been shown to increase compared to the unstructured

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30 case [7] by cooperators forming clusters that are resistant to exploitation by surrounding defectors, an effect
31 known as network reciprocity. Developing these ideas the evolution of cooperation on networks received
32 intense subsequent study, with the effect on cooperation of network structure [8, 9, 10, 11], the type of
33 game [12], the rules used to update the strategy [13], additional strategies such as punishment and loners
34 [14, 15, 16] and diversity between players [17] investigated. In each case it was found that added heterogeneity
35 introduces higher levels of cooperation, as cooperators are able to seed clusters more easily and boost the
36 effect of network reciprocity.

37 In scenarios such as electricity power networks and networks of air travel it is more accurate to describe
38 a system as a combination of networks than a single one [18, 19, 20, 21, 22]. These multilayered systems
39 can take two forms: either interdependent networks, where two separate networks are connected through a
40 number of between-network edges, or multiplex networks, where the nodes on each layer represent the same
41 entity, but each layer represents a different aspect of the system [23].

42 How the amount of cooperation on a multilayered system varies has been studied in a number of cases,
43 from interdependent networks of different topologies [24, 25] to lattices with weighted payoffs between the
44 layers [26, 27] and weighted fitnesses [28], probabilistic connections between layers [29, 30], different imitation
45 rules between layers [31] and memory [32]. In each of these cases cooperation is enhanced by the addition of
46 extra networks caused by strategies on each layer supporting each other against exploitation by defectors,
47 and there is a peak interdependence between the two layers due to maximum heterogeneity in the system,
48 even when the games played on each network are different [33].

49 The dynamics of games have also been studied on multiplex networks [34]. These are networks with
50 more than one layer, where the nodes on each layer represent the same player, but the connections between
51 them are not necessarily the same. On these networks players can play different strategies on each layer,
52 but their payoffs are summed across all layers. Cooperation was found to be increased compared to the
53 single layer when games are played on random network multiplexes due to the ability of players to play a
54 differing strategy on each layer [35, 36]. These “incoherent” players are able to maintain cooperation when
55 it would usually have disappeared in the single layer system. The effect of separating the networks on which
56 players calculate their payoffs and those on which they update their strategies onto different layers has been
57 explored [37], as has how the assortativity of the degrees between the networks affects the final frequency
58 of cooperation, what happens if the game played on each layer of the multiplex is different [38], and if
59 players imitate those that they are most similar to [39]. Zhang *et al.* [40] found the analytical conditions
60 for cooperation on a multiplex by studying structured groups across a two-layer network, where cooperation
61 was found to be maximised when intermediate levels of migration between groups was implemented.

62 Despite the existing work discussing cooperation in multilayered networks, more needs to be done in
63 order to understand exactly what is important in these systems for cooperation to be maintained, and how
64 cooperation can be encouraged in order to offer solutions to social dilemmas. One assumption that is made

65 in a number of previous models is about how frequently the nodes on each layer update themselves. It
66 has been found [13, 41] that the proportion of nodes that update their strategies on a network before the
67 payoff of each player is recalculated can alter the amount of cooperation in the system. What has not
68 been considered in the case of multilayered networks is whether updating each of the nodes across all layers
69 compared to just a single layer at a time affects the final amount of cooperation in the system. Whether a
70 player updates their strategy on all layers of the system at the same time or separately will depend strongly
71 on the system that is to be modelled. It is therefore important to know whether this has any effect on the
72 final frequency of cooperation in the system. This is the question that we will address in this article.

73 2. Mathematical Model

74 In order to study the spread of cooperation on the multiplex network the public goods game (PGG) is
75 played by each of the nodes. In this game there are two possible strategies, either cooperation or defection.
76 In the PGG the players are divided into groups before donating a certain amount to the groups that they
77 are members of. The total that has been donated is then multiplied up by an enhancement factor, before
78 being divided between all of the members of that group. When playing the PGG on a network the groups
79 are defined to be those players that are connected to a common node. So, on each layer the player plays
80 the game in $k + 1$ groups, where k is the degree of the player on that particular layer. In the PGG, player i
81 donates an amount c_i , and so the payoffs are given by

$$P_D^i = r \frac{\sum_{j=1}^{N_C} c_j}{G} \quad (1)$$

$$P_C^i = r \frac{\sum_{j=1}^{N_C} c_j + c_i}{G} - c_i \quad (2)$$

82 where $P_{C,D}^i$ is the payoff of player i playing as a cooperator or a defector respectively, G is the number of
83 players in the group, N_C is the number of cooperators in the group, r is the enhancement factor (the return
84 on the investment to the group) and the last term in equation (2) represents the donation by the cooperators
85 to the public good. The final payoff for each player is found by summing across all of the groups in which
86 they play across all of the layers. We will study two possible versions of the PGG: fixed cost per group
87 (FCG) when each player donates one unit to whichever group they are playing the game in e.g. $c_i = 1$ for
88 all groups, or fixed cost per individual (FCI) where each player has only a single unit to donate and divides
89 this between all of the groups in which it plays e.g. $c_i = \frac{1}{k_i+1}$ where k_i is the degree of node i .

90 Initially each player on each layer is assigned a strategy (cooperate or defect) at random, which need
91 not be the same on every layer. Each node plays the PGG against their neighbours on each layer, with each
92 player's payoff accumulated across all layers. Once the payoffs have been calculated the players have an
93 opportunity to update their strategy through imitation. In general, it makes sense for a player to imitate

94 those that are performing better. Therefore on selected layers each player chooses a neighbour at random
 95 and compares payoffs, imitating the neighbour's strategy with probability calculated below.

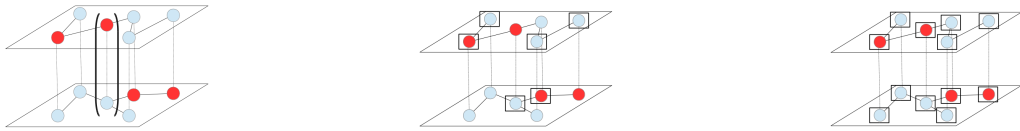
96 We study the game on a multiplex consisting of two Erdős-Rényi random networks of size N nodes. These
 97 networks are formed independently on each layer by creating an edge between two nodes with probability
 98 $\frac{\langle k \rangle}{N-1}$, where $\langle k \rangle$ is the mean degree of each node. The final payoff of player i is calculated by summing the
 99 payoffs for each layer e.g. $P_i = \sum_{j=1}^L p_{ij}$ where L is the total number of layers and p_{ij} is the payoff of player
 100 i on layer j . A player is more likely to imitate another's strategy if that player has a higher payoff, and the
 101 probability of imitation should be proportional to the payoff difference. Therefore, in common with many
 102 other studies of cooperation on networks [8, 9, 10, 11] player i will imitate the strategy of neighbour j with
 103 probability

$$P(s_i \rightarrow s_j) = \frac{1}{1 + e^{-\frac{P_j - P_i}{\beta}}} \quad (3)$$

104 where β is a parameter that defines how responsive the player is to their neighbour's payoff. Here $\beta = 0$
 105 leads to deterministic dynamics where a player will always imitate the strategy of a player with a larger
 106 payoff, and $\beta \rightarrow \infty$ leads to players imitating their neighbours strategy at random, with no consideration of
 107 their respective payoffs.

108 In this article we investigate the effect of updating the players strategies either asynchronously or syn-
 109 chronously across layers on the final frequency of cooperation. In the asynchronous case, after the payoffs of
 110 every player have been calculated each player selects one layer at random on which to update their strategy,
 111 and then a random neighbour on that layer to compare total payoffs. When updating synchronously, each
 112 player selects a neighbour to imitate on each layer (not necessarily the same neighbour on each). On this
 113 multiplex, each player is allowed to play a different strategy on each layer even though they are the same
 114 individual. Examples of a multiplex network, where a single player and its strategy on each layer are high-
 115 lighted, and examples of each rule are shown in figure 1, where the nodes in the hollow squares are those
 116 updating their strategy. In order to compare the different rules the asynchronous update rule will update
 117 twice per time step and the synchronous rule just once so that each node updates the same number of times
 118 on average.

119 We study how updating the player strategies either asynchronously or synchronously affects the final
 120 frequency of cooperation for both the PGG and the PD. In the results presented below for each iteration
 121 and for each rule a new network and initial conditions are generated. In order to ensure that the system
 122 has reached equilibrium each node is updated 50,000 times on average for each run. The final frequency of
 123 cooperation is then averaged over the last 5000 complete updates of the system. Results for each simulation
 124 are plotted as points and the mean of all of the simulations is shown by a line. All results presented in this
 125 article (unless otherwise stated) are for $m = 2$ layer multiplex networks with $N = 1000$ nodes on each layer.



(a) An example of a multiplex network (b) Asynchronous update (c) Synchronous update

Figure 1: An example of a multiplex network (1a) with a single player and their strategy on each layer indicated, the asynchronous update (1b) and the synchronous update (1c) rules. The nodes that are contained in squares are those selected to update their strategy in that round. Red dots denote defectors and blue dots denote cooperators.

126 **3. Results**

127 *3.1. The PGG on a two-layer multiplex*

128 *3.1.1. The fixed cost per individual PGG*

129 The first set of results shows the mean frequency of cooperation across both layers plotted against the
 130 scaled enhancement factor $\eta = \frac{r}{\langle k \rangle + 1}$ in the case of the fixed cost per individual (FCI) PGG for $m = 2, \langle k \rangle =$
 131 $3, \beta = 0.5$ (figure 2).

132 Updating the strategies on each layer of the multiplex asynchronously leads to a higher frequency of
 133 players choosing the cooperative strategy within the network compared to updating synchronously (figure
 134 2). At $\eta = 0.76$ the difference is very large, with almost no cooperation present in the synchronous update
 135 system, and between 30 and 100% in the asynchronous update, with a similarly large difference seen for
 136 $\eta = 0.78$. The value of η at which the system changes from complete defection to mixed defection and
 137 cooperation is higher in the synchronous update case. Similarly, the value of η at which the system changes
 138 from a mixture of strategies to complete cooperation is also higher in the case of the synchronous update.
 139 This demonstrates that asynchronous updating can act as a mechanism to maintain cooperation at lower
 140 enhancement factors that has not previously been considered.

141 To explain why there is a difference in the frequency of cooperation between the two update rules a
 142 cartoon of three players on the multiplex, s, t and u , is shown in figures 3a and 3b. The strategies for player
 143 s are s_i on the top layer and s_j on the bottom, and similarly for players t and u .

144 On the top layer, player s has selected player u to imitate, and on the lower layer player s will imitate
 145 player t . Assuming initially that $P_u > P_s$ and $P_t > P_s$, and that β is not large, figure 3a demonstrates the
 146 likely strategies after a single time step, with player s imitating player u on the top layer and player t on
 147 the lower layer. When updating using the asynchronous rule the neighbours (and the payoffs) of the focal
 148 cooperator and defector may change before the cooperator has a chance to imitate the defector, resulting
 149 in a difference in imitation probability. If this probability changes by a large amount then player s will no
 150 longer imitate player t on the lower layer. Therefore updating the multiplex asynchronously rather than
 151 synchronously may lead to different frequencies of cooperation. A similar mechanism was found in reference
 152 [42], which found increased cooperation through strategy correlation across different network layers.

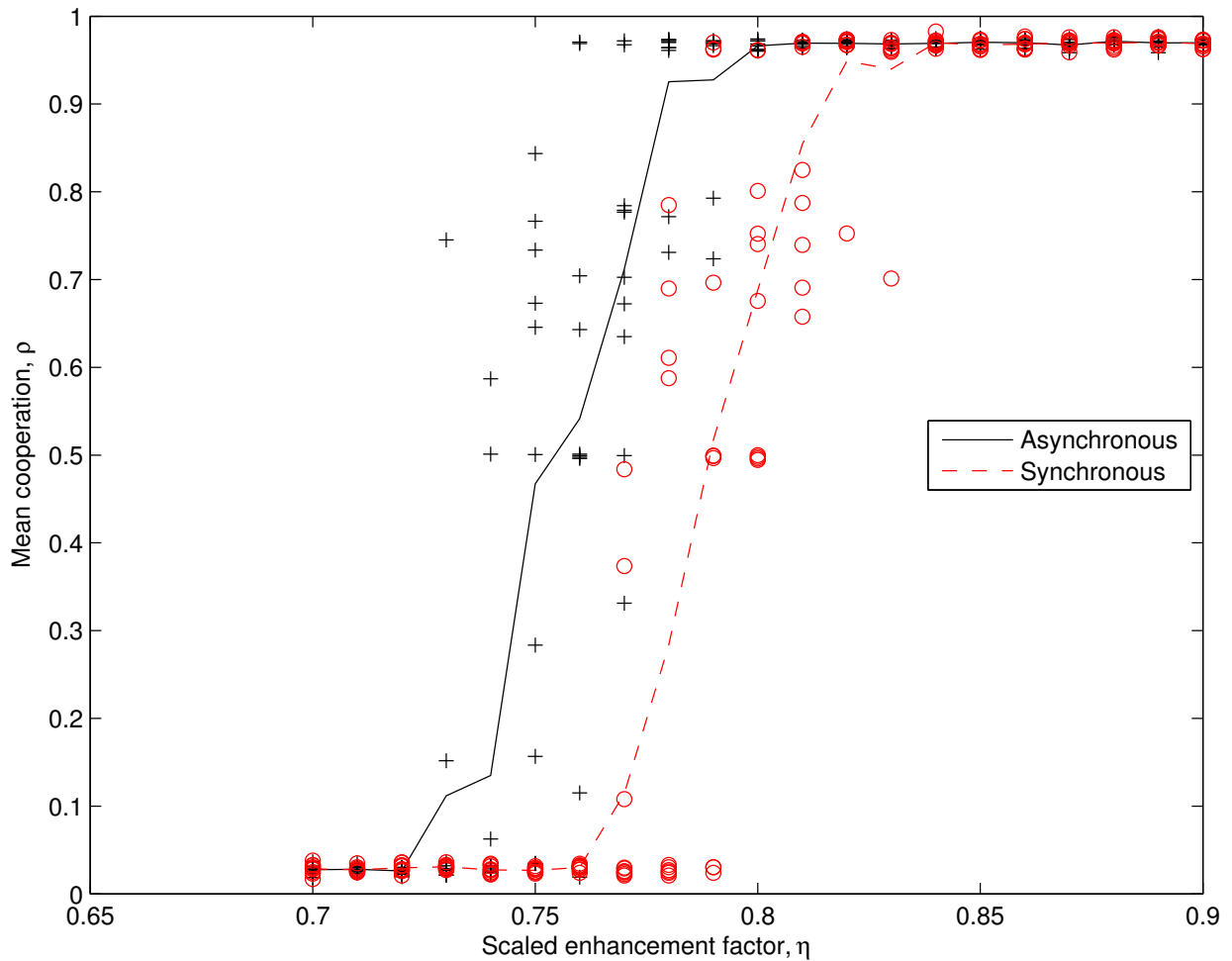
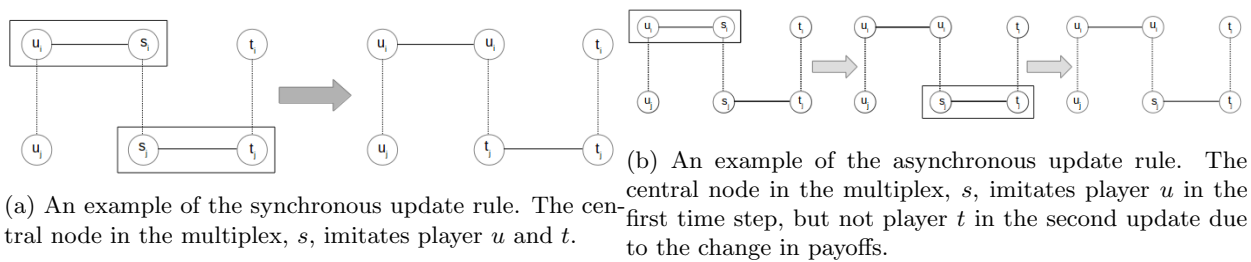


Figure 2: The frequency of cooperation on the two-layer PGG plotted against the scaled enhancement factor $\eta = \frac{r}{\langle k \rangle + 1}$ for the asynchronous (solid black line) and synchronous (dashed red line) update rules for fixed cost per individual (FCI). $m = 2, N = 1000, \langle k \rangle = 3, \beta = 0.5$.



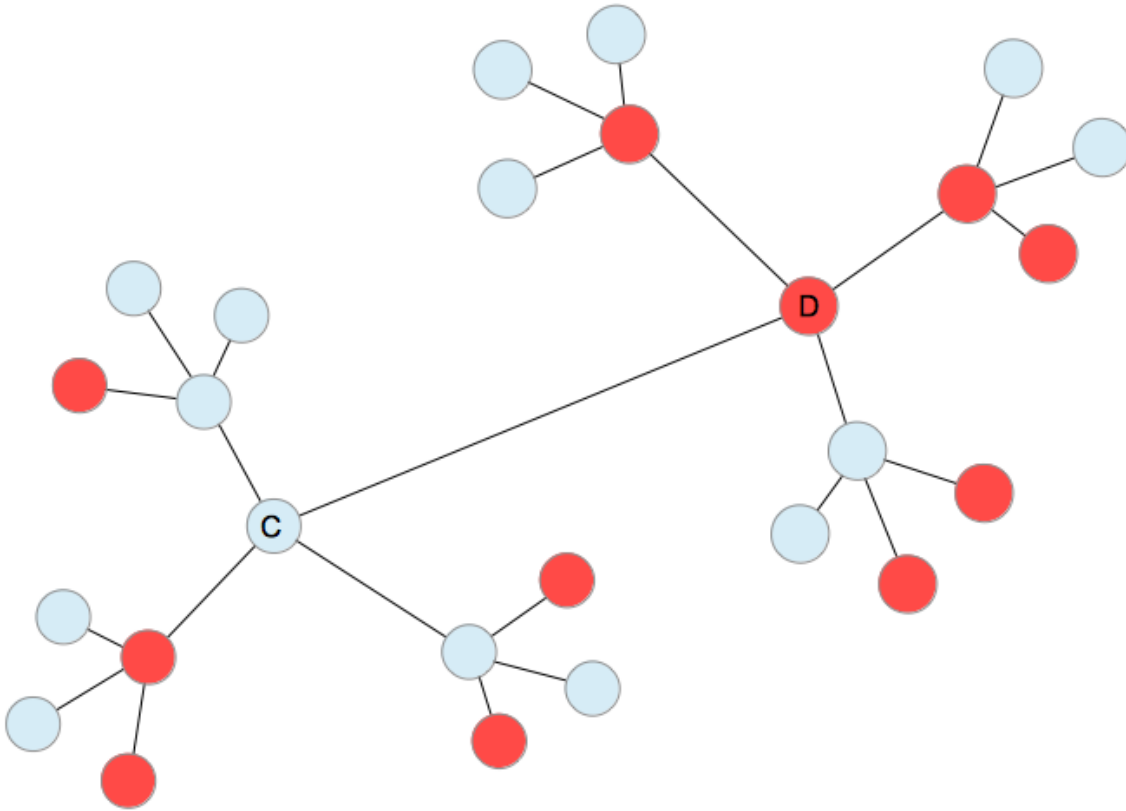


Figure 3: An example of a neighbouring cooperator and defector (labelled C and D respectively), and the composition of the neighbouring groups in which they accumulate payoffs (other cooperators coloured light blue, other defectors coloured dark red, colour online).

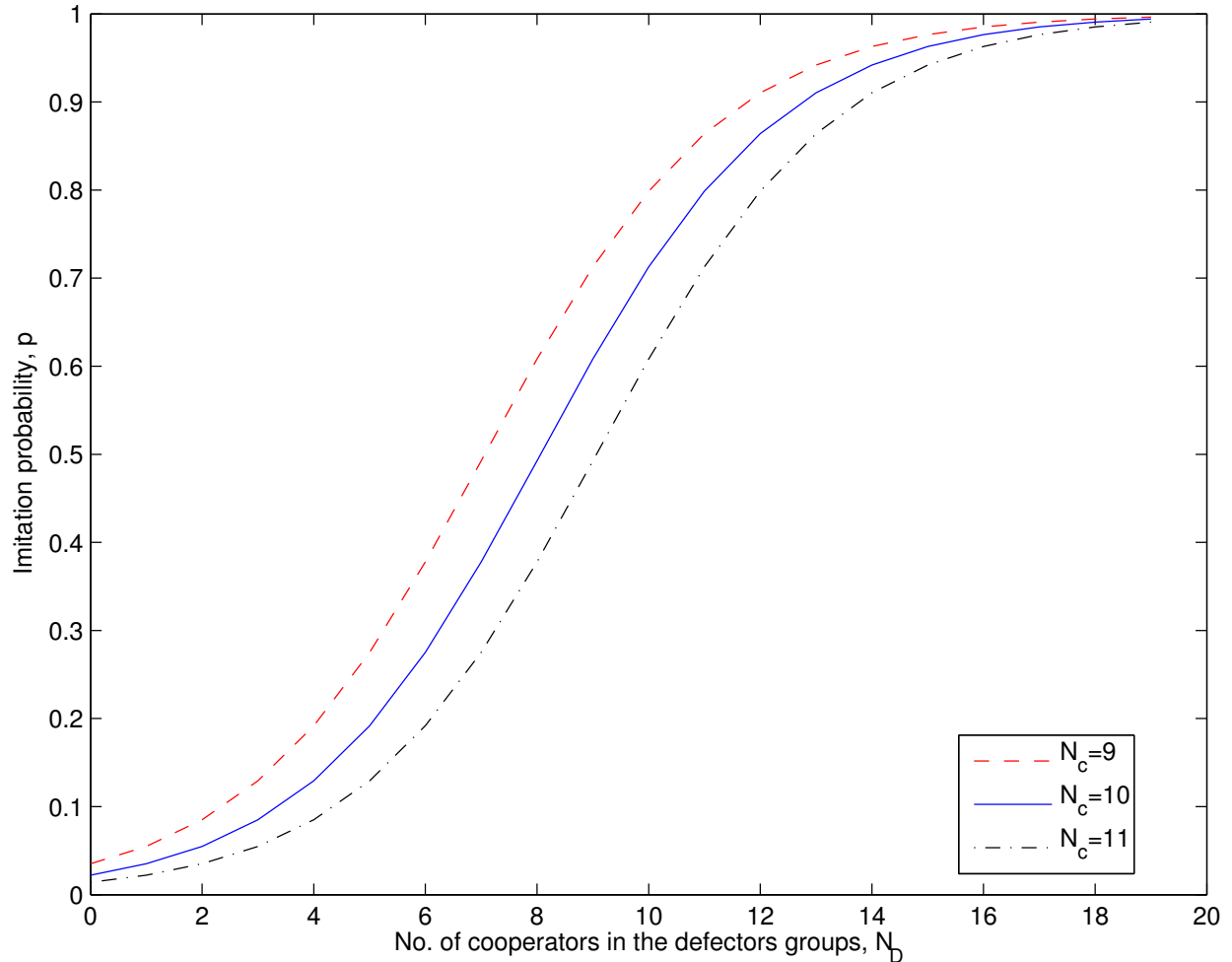


Figure 4: The probability of a cooperator in a neighbourhood composed of N_C cooperators imitating a defector plotted against number of cooperators in the defector's neighbourhood, N_D for the FCI PGG. $r = 0.78, \beta = 0.5, \langle k \rangle = 3$.

153 We numerically calculate the expected change in probability by artificially forming neighbourhoods where
154 a central cooperator and defector (labelled 'C' and 'D' in figure 3) play in groups consisting of N_C and N_D
155 cooperators respectively. The probability of a cooperator in a neighbourhood consisting of N_C cooperators
156 imitating a defector surrounded by N_D cooperators is calculated using equation (3). Averaging these proba-
157 bilities over all possible arrangements of N_C and N_D an expected imitation probability is calculated. Figure
158 4 plots this expected probability against N_D for $N_C = 9, 10, 11$ for the FCI PGG ($r = 0.78, \beta = 0.5, \langle k \rangle = 3$).
159 This plot shows that if the group composition changes by a large amount, then the probability of a coop-
160 erator imitating a neighbouring defector will be significantly affected. Therefore, there is likely to be a
161 large difference in the probability of a cooperator imitating a neighbouring defector under synchronous and
162 asynchronous updating.

163 These results are confirmed empirically by running a single iteration of the dynamics. Recording that
 164 player i selects player j to compare payoffs at time step t enables the payoffs of the two players to be
 165 measured at both t and $t - 1$ (for $t > 1$). Writing these payoffs as $P_i(t), P_j(t), P_i(t - 1)$ and $P_j(t - 1)$
 166 and substituting them into equation (3) for both t and $t - 1$, the change in the probability of imitating
 167 a neighbouring strategy over two payoff updates can be measured. A histogram plotting the difference in
 168 imitation probabilities between two neighbours for consecutive time steps for the asynchronous rule is shown
 169 in figure 5 for $r = 0.78, N = 1000, \beta = 0.5$.

170 In this histogram approximately 45% of the update probabilities change between each payoff update (in
 171 the synchronous case this number would be 0%). Therefore, starting from identical networks with identical
 172 strategies, the probability of a given cooperator imitating a neighbouring defector (and vice versa) is altered
 173 when comparing the asynchronous and synchronous update rules. Because the largest change in group
 174 compositions occurs when the fraction of cooperation is approximately 50%, this explains why the largest
 175 difference between the two rules is observed at this frequency of cooperation (figure 2).

176 To explain why cooperation is higher for the asynchronous update rule the probability of player i im-
 177 itating the strategy of neighbour j for each payoff difference $P_j - P_i$ is empirically measured for both the
 178 asynchronous and synchronous update rules. We collected results from a single simulation for 1000 updates
 179 of the entire network, and calculated the probabilities by placing each measured payoff difference into a bin
 180 of width 0.1. Figure 6 shows the probability that a player imitates a neighbour's strategy plotted against
 181 the payoff difference, with the plot on the left showing the probability of a cooperator imitating a defector
 182 and the plot on the right showing the probability that a defector imitates a cooperator.

183 What is observed for both update rules (figure 6) is that for positive payoff differences the probability
 184 of strategy imitation is high, and increases as the difference increases, the reverse being true for negative
 185 payoff differences. In the case of the synchronous update rule (the dashed red line) the curve obeys equation
 186 (3) for most payoff differences, as expected. The probabilities become noisier at very low payoff differences,
 187 but this is due to the small number of empirical data points. In contrast the asynchronous rule (the solid
 188 black line) does not follow equation (3) particularly well for positive payoff differences. For most of the
 189 range of the positive payoff differences both the probability of a cooperator imitating a defector, and vice
 190 versa, is lower for each payoff difference compared to the results for the synchronous update rule (although
 191 at negative and very small positive payoff differences the two rules almost agree). In the asynchronous
 192 case a cooperator is less likely to imitate a defector than in the synchronous case, and this promotes the
 193 maintenance of cooperation throughout the network. Hence, the asynchronous and synchronous update
 194 rules lead to different frequencies of cooperation in the population due to the additional payoff update
 195 in each time step in the asynchronous case, and this leads to a reduced probability of a player imitating
 196 a better performing neighbour. In the parameter range investigated here ($\eta < 1$) in the PGG a better
 197 performing neighbour is likely to be a defector, and so the frequency of cooperation is likely to be higher in

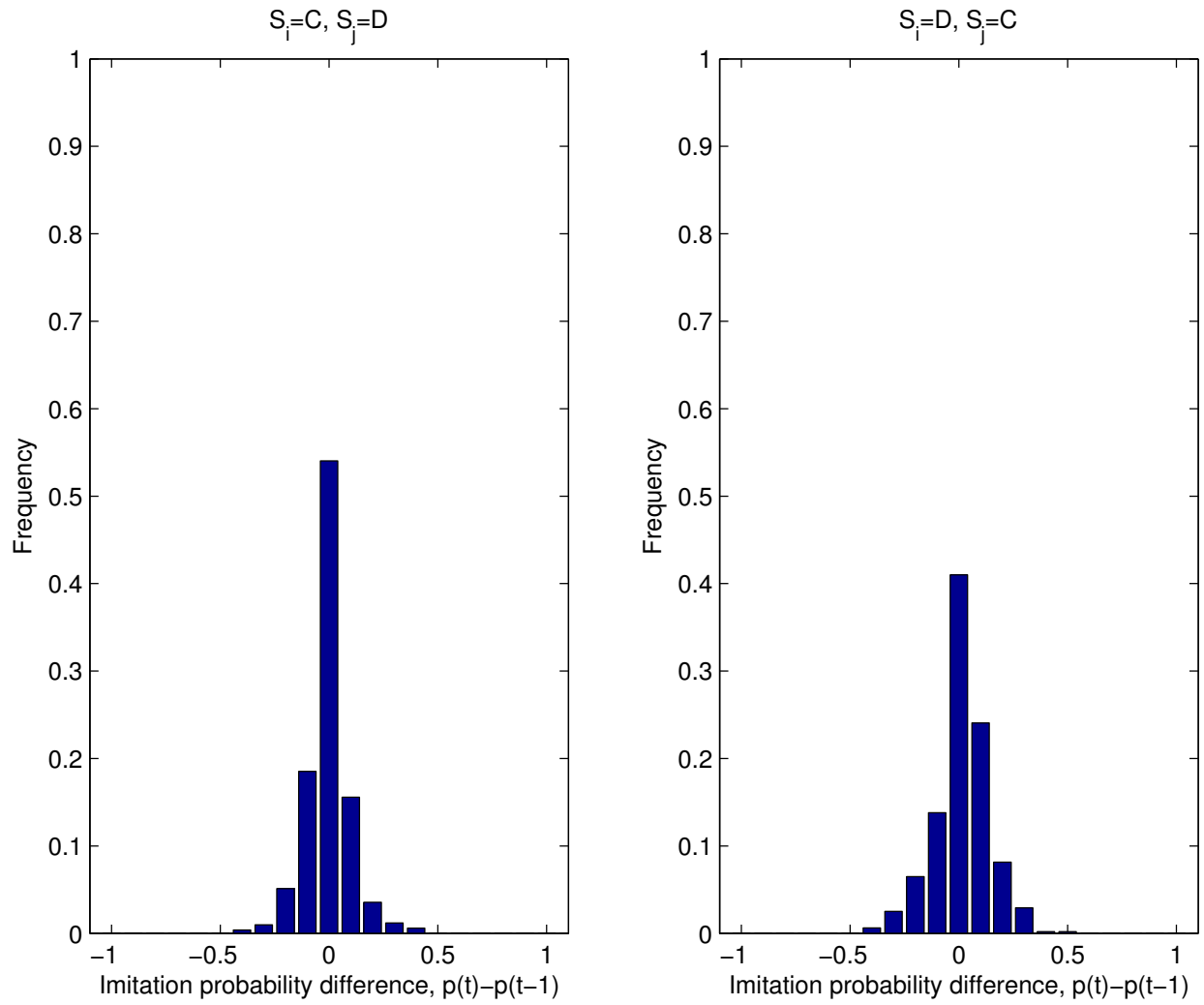


Figure 5: Histogram showing the change in imitation probabilities between players i and j after a single update of the payoffs for $s_i = C, s_j = D$ (left) and $s_i = D, s_j = C$ (right) for the FCI PGG. $r = 0.78, m = 2, N = 1000, \beta = 0.5$

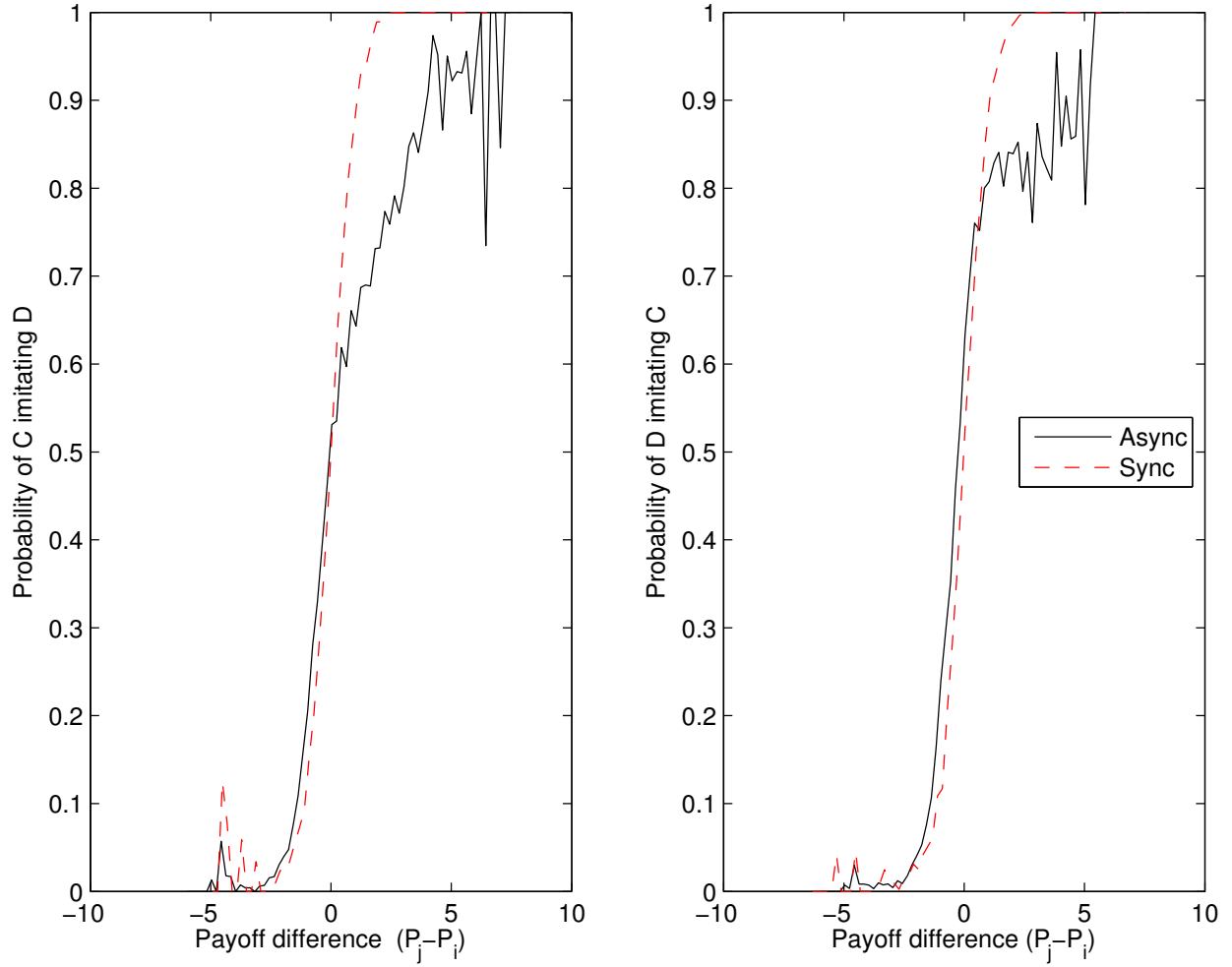


Figure 6: The probability of strategy imitation on the two-layer multiplex for the asynchronous (solid black line) and synchronous (dashed red line) update rules for fixed cost per individual (FCI). $N = 1000, m = 2, \eta = 0.76, \langle k \rangle = 3, \beta = 0.5$.

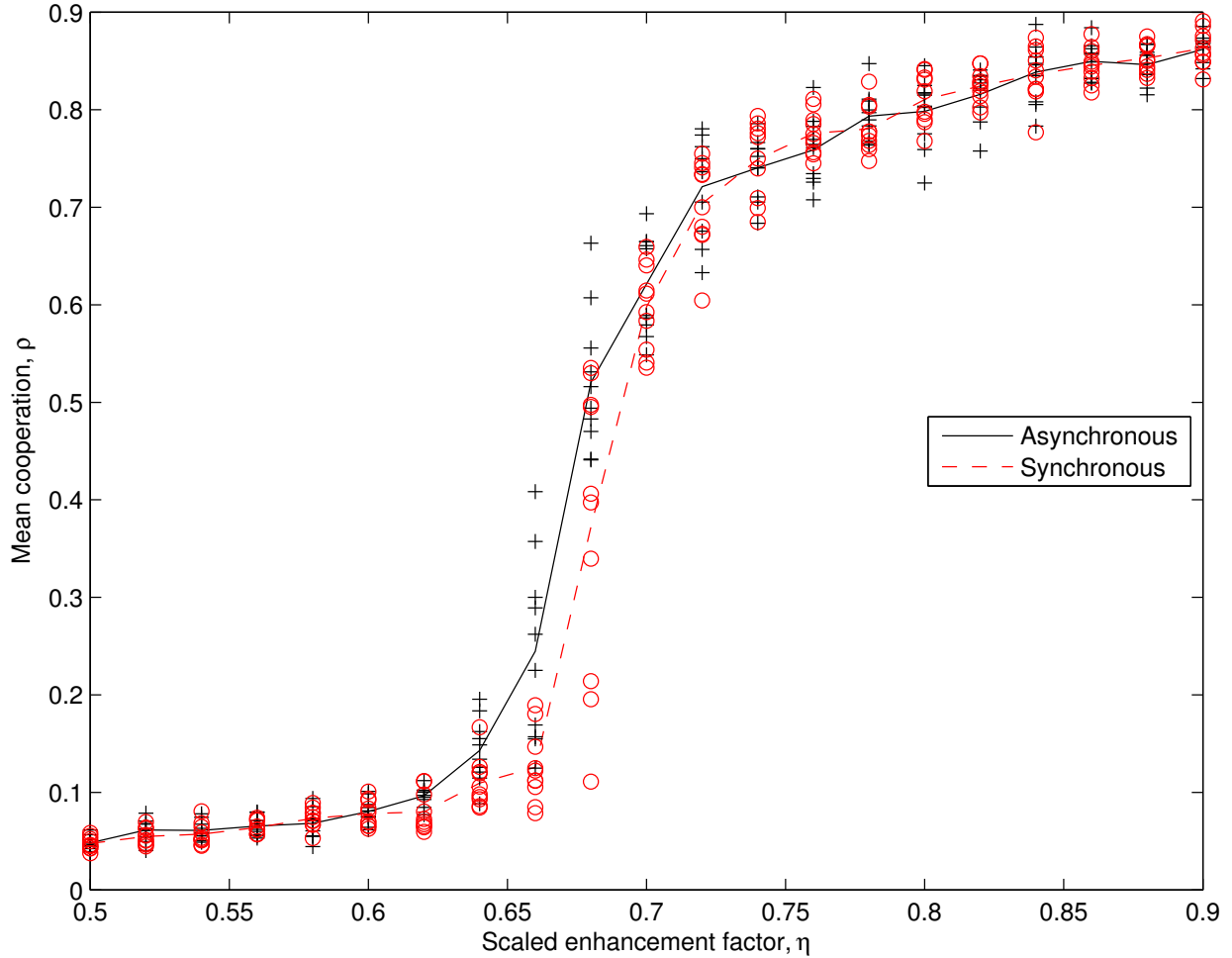


Figure 7: The frequency of cooperation on the two-layer PGG plotted against the scaled enhancement factor $\eta = \frac{r}{\langle k \rangle + 1}$ for the asynchronous (solid black line) and synchronous (dashed red line) update rules for fixed cost per individual (FCI). $N = 1000, m = 2, \langle k \rangle = 3, \beta = 0$.

198 the asynchronous case. This also explains the asymmetry between the asynchronous results in each plot, as
 199 cooperators are less likely to imitate better performing defectors.

200 3.2. Additional results on the PGG

201 3.2.1. Deterministic ($\beta = 0$) dynamics for the fixed cost per individual PGG

202 Testing the importance of the rule for other values of β , figure 7 shows that for FCI deterministic updates
 203 ($\beta = 0$) the asynchronous update rule once again leads to a higher final frequency of cooperation for a range
 204 of enhancement factors. However, the differences in figure 7 are much smaller than in figure 2, with some
 205 asynchronous and synchronous results overlapping for the entire range of η .

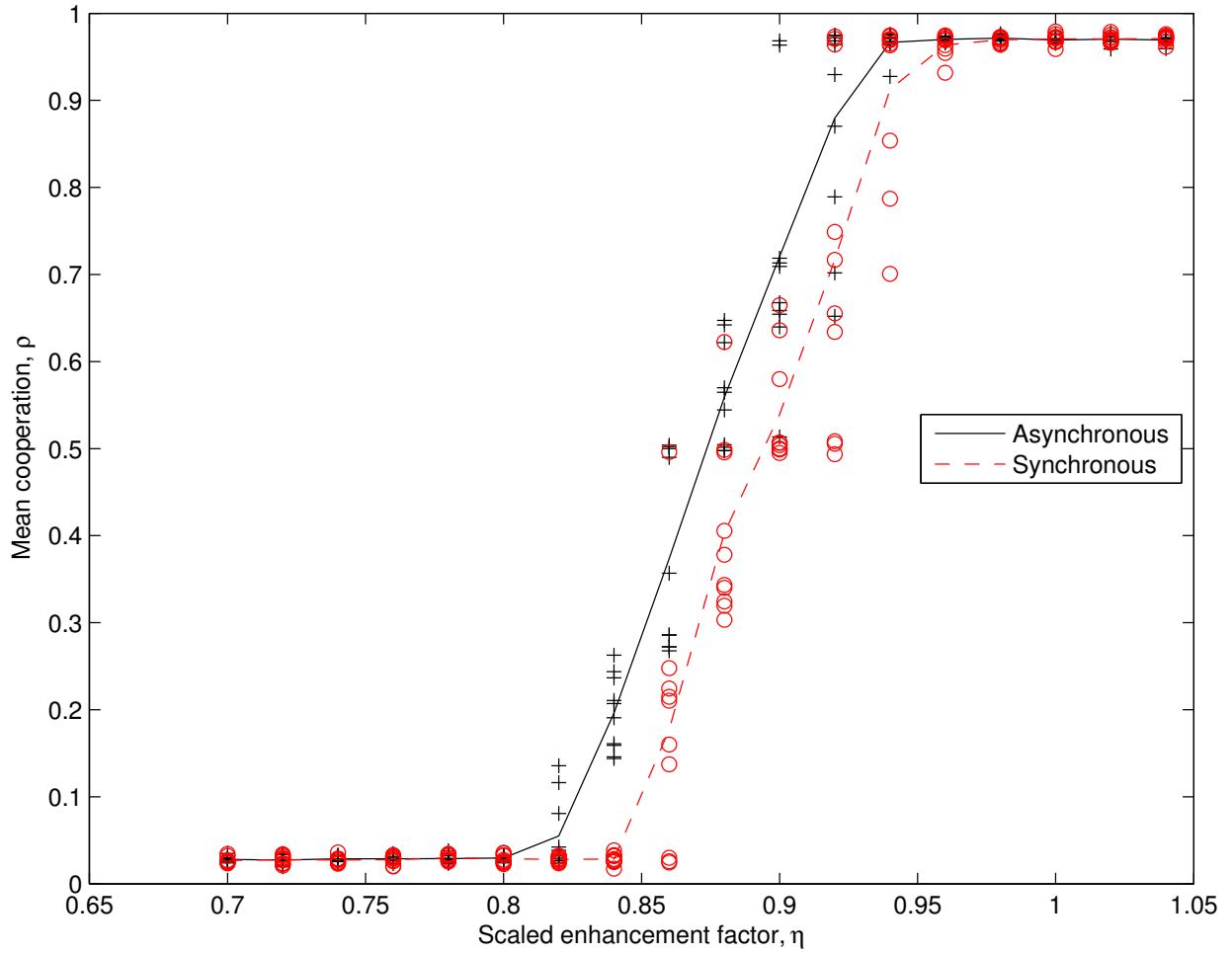


Figure 8: The frequency of cooperation on the two-layer PGG plotted against the scaled enhancement factor $\eta = \frac{r}{\langle k \rangle + 1}$ for the asynchronous (solid black line) and synchronous (dashed red line) update rules for fixed cost per group (FCG). $N = 1000, m = 2, \langle k \rangle = 3, \beta = 0.5$.

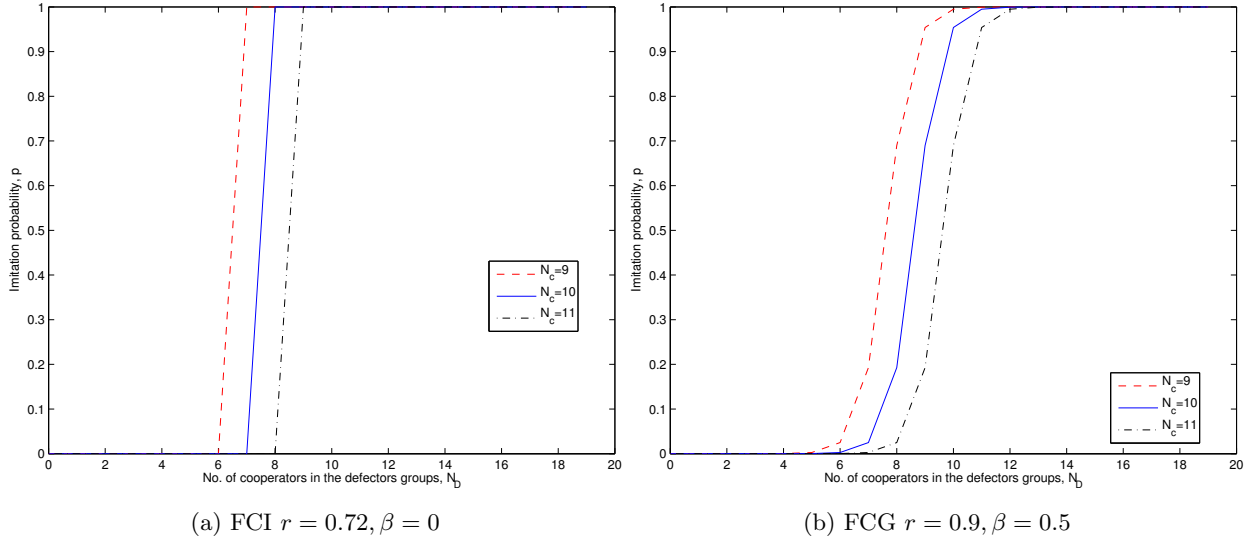


Figure 9: The probability of a cooperator in a neighbourhood composed of N_C cooperators imitating a defector plotted against number of cooperators in the defector’s neighbourhood, N_D for the FCI PGG ($r = 0.72, \beta = 0$) and the FCG PGG ($r = 0.9, \beta = 0.5$). $\langle k \rangle = 3$.

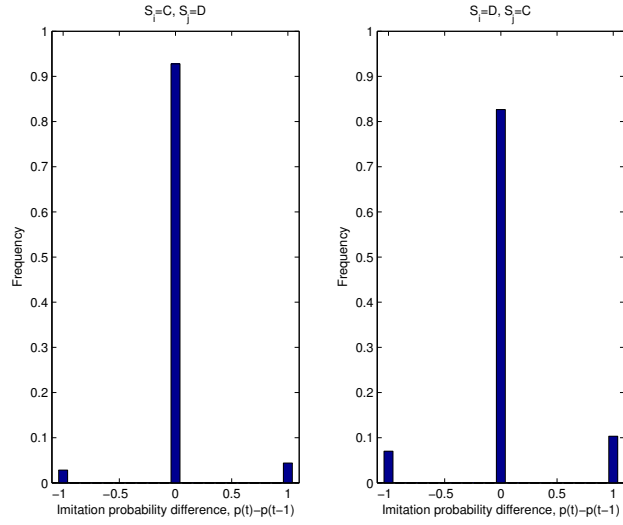
206 *3.2.2. The fixed cost per group PGG*

207 Results for the fixed cost per group (FCG) PGG are presented in this section. Figure 8 shows the effect
 208 of the choice of update rule when the players contribute the same to each group (FCG) in the PGG. Once
 209 again, in general the asynchronous update rule leads to higher frequencies of cooperation for a wide range
 210 of η (figure 8). However, comparing these results to the FCI case in figure 2 the choice of update rule makes
 211 less of a difference to both the final frequency of cooperation and the enhancement factor at which the
 212 system produces a mixed state of cooperators and defectors or fixates to a single strategy.

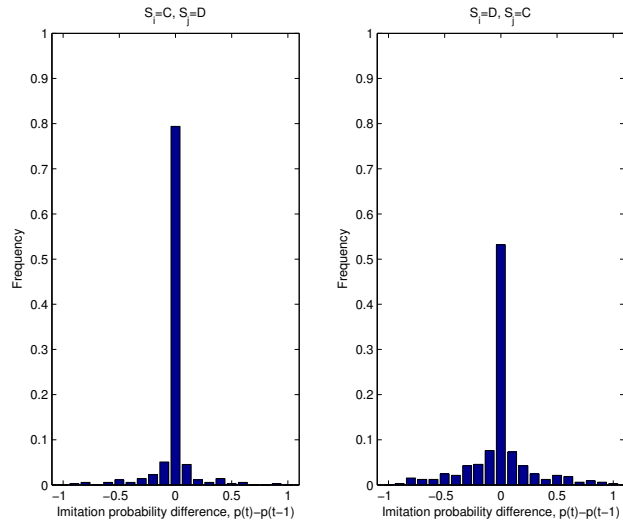
213 *3.3. Comparison of imitation probabilities for FCI $\beta = 0$ and FCG $\beta = 0.5$*

214 We repeated the calculations for imitation probabilities performed in section 3.1 for the FCI with no
 215 noise ($\beta = 0$) and present the results in figure 9a. The imitation probabilities are also calculated for the
 216 FCG PGG ($\beta = 0.5$) and the results shown in figure 9b. Comparing these results with those for the FCI
 217 ($\beta = 0.5$), the imitation probability curves in figures 9a and 9b are much steeper compared to those in figure
 218 4. Therefore, if the number of cooperators in a group changes between payoff updates this has less of an
 219 effect on the probability of strategy imitation, and so the choice of update rule will have less of an effect on
 220 the final frequency of cooperation.

221 These results are confirmed by the empirical measurements of the change in imitation probability in figure
 222 10 for the FCI PGG ($\beta = 0$) and the FCG PGG ($\beta = 0.5$). Comparing the frequency of $p(t) - p(t - 1) \neq 0$
 223 for the FCI ($\beta = 0$) (figure 10a) and the FCG ($\beta = 0.5$) (figure 10b) with the results for the noisy FCI
 224 ($\beta = 0.5$) (figure 5), the scenario with the largest imitation probability change is the FCI PGG ($\beta = 0.5$),



(a) FCI $r = 0.72, N = 1000, \beta = 0$



(b) FCG $r = 0.9, N = 1000, \beta = 0.5$

Figure 10: Histogram showing the change in imitation probabilities between players i and j after a single update of the payoffs for $s_i = C, s_j = D$ and $s_i = D, s_j = C$ for the PGG. Figure 10a FCI $r = 0.72, N = 1000, \beta = 0$, figure 10b FCG $r = 0.9, m = 2, N = 1000, \beta = 0.5$.

225 followed by FCG ($\beta = 0.5$) and finally FCI ($\beta = 0$). Therefore, the choice of update rule will have the
 226 most impact for the FCI ($\beta = 0.5$), as here the probability of imitating a neighbouring defector changes
 227 the most between payoff updates. What these results show is that when imitation probabilities are strongly
 228 dependent on the composition of the groups in which the players outcomes will be significantly affected by
 229 the choice of update rule.

230 3.4. The PD on a two-layer multiplex

231 In this section we consider another common game used to model cooperation, namely the prisoner's
 232 dilemma (PD). In the PD the game is played in pairs rather than groups, and the payoff for each player is
 233 the sum of the payoffs in each of these pairwise interactions. As in the PGG the neighbours in the network
 234 denote who each player plays against, and the total payoff for each player is found by summing over each of
 235 the pairwise interactions across all of the layers. The standard way to represent the payoffs gained is in the
 236 form of a matrix, where the rows define the strategy chosen by player one, and the columns the strategy
 237 chosen by player two, with the values in the matrix giving the payoffs allocated to each player.

Table 1: The payoffs for the pairwise game.

| | | |
|---|-----|-----|
| | C | D |
| C | R,R | S,T |
| D | T,S | P,P |

238 In the PD the payoffs are ordered $T > R > P > S$, and often this is simplified to $T = b, R = 1, P = S = 0$,
 239 where $b > 1$ is the parameter of interest and is labelled as the ‘temptation to defect’.

240 Figure 11 shows results for the PD run on a two-layer Erdős-Rényi multiplex network, for parameters
 241 $N = 1000, \langle k \rangle = 3$. The dynamics in these simulations are very similar to those in section 3.1, with each
 242 player selecting a neighbour at random and comparing payoffs. When examining the PD on networks two
 243 standard update rules are investigated. In figure 11a the Fermi imitation rule (equation (3)) is used in
 244 order to decide if player i should imitate player j . In the original model of cooperation on a multiplex (see
 245 reference [35]) strategy imitation occurs with probability

$$p(s_i \rightarrow s_j) = \frac{P_j - P_i}{K_{max}b} \quad (4)$$

246 where P_i is the payoff of player i as usual, K_{max} is the largest sum of the degrees of player i and j across
 247 all of the layers, and b is the temptation to defect. We show that the choice of update rule also alters the
 248 frequency of cooperation if this different imitation rule is used (figure 11b)

249 Figure 11 shows the final frequencies of cooperation for the asynchronous and synchronous update rules
 250 for these two different imitation rules. For Fermi imitations (Figure 11a) results are very similar to those
 251 found in section 3.1: for $1.61 < b < 1.86$ the asynchronous update rule leads to a higher frequency of

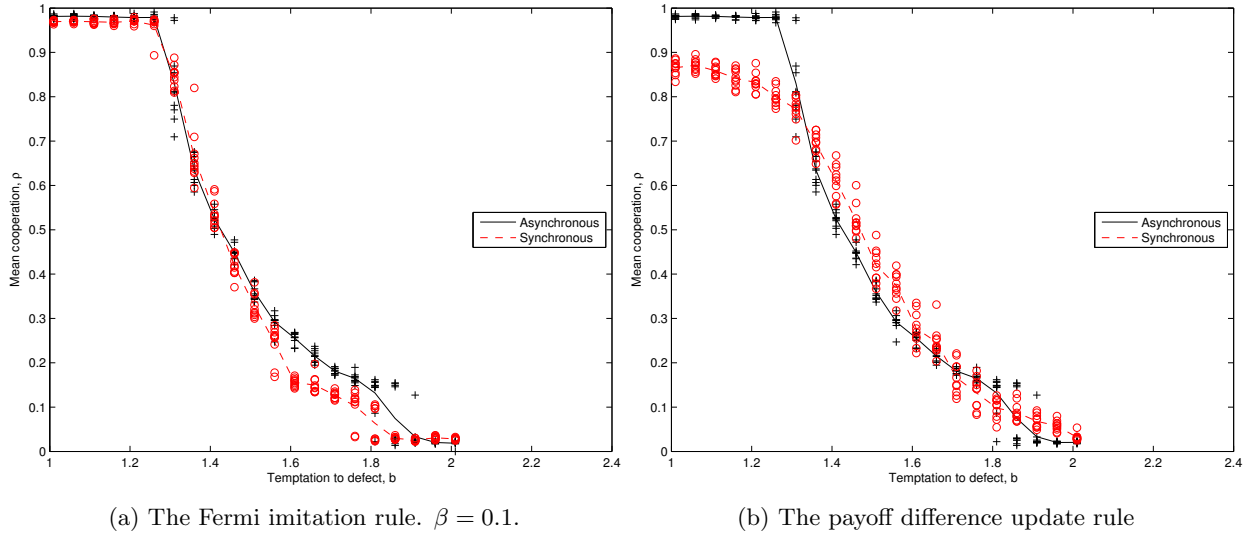


Figure 11: The frequency of cooperation on the two-layer PD plotted against the temptation to defect for the asynchronous (solid black line) and synchronous (dashed red line) update rules for the Fermi imitation rule ($\beta = 0.1$) 11a and the payoff difference update rule 11b. $N = 1000, m = 2, \langle k \rangle = 3$.

252 cooperation compared to the synchronous rule, although the effect is not large. When imitation probabilities
 253 follow equation (4) there is a significant difference between the two rules for low b , and then little difference
 254 as the temptation to defect increases.

255 This difference is explained by plotting the imitation probabilities for $b = 1.1$ for both the Fermi and the
 256 payoff difference update rules. For $b = 1.1$ the frequency of cooperation is high, and so the likely number of
 257 cooperators in the neighbouring defector's group is large. Figure 12a shows that in the Fermi imitation rule
 258 there is no difference between the different number of cooperators. Therefore, if the number of cooperators
 259 in the group changes through the additional strategy update in the asynchronous case, this will not lead
 260 to a large difference in imitation probability. This is not what is observed in figure 12b, where there is a
 261 difference in the imitation probability between each value of N_D . Therefore, if the number of cooperators
 262 in a group changes, this will lead to a different imitation probability.

263 4. Discussion and conclusion

264 We have shown that the asynchronous update rule consistently leads to higher frequencies of cooperation
 265 over a range of enhancement factors in the public goods game, and can also increase cooperation in the
 266 prisoner's dilemma. Through numerical simulations we find that the strength of the effect depends on a
 267 number of factors, including the type of game played, the noise and the strategy imitation rule. In each of
 268 these cases when the probability of a player of one strategy imitating another is highly sensitive to changes in
 269 the player's neighbourhood, the choice of asynchronous or synchronous update rule has more of an effect on
 270 the final frequency of cooperation. The public goods game is most sensitive to changes in the neighbourhood

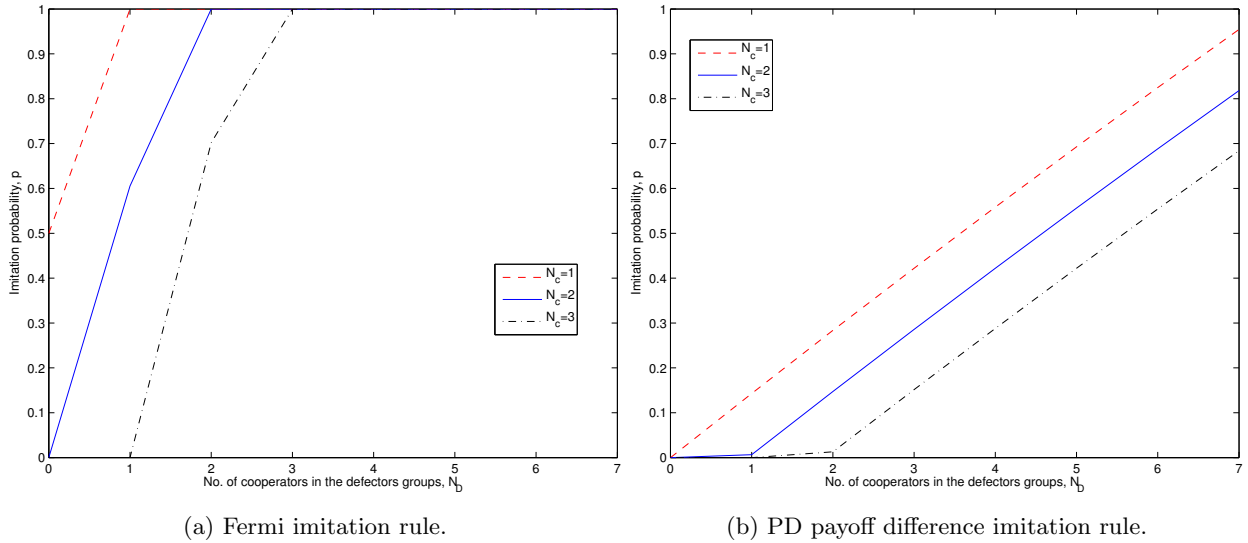


Figure 12: The probability of a cooperator in a neighbourhood composed of N_C cooperators imitating a defector plotted against number of cooperators in the defector's neighbourhood, N_D for the PD Fermi imitation and the PD payoff difference update. $b = 1.1$, $\langle k \rangle = 3$.

271 when a player divides a single unit between its neighbours (the FCI PGG), compared to when a single unit
 272 is donated to all (the FCG PGG). Therefore, the choice of rule is more significant for the FCI PGG. In
 273 contrast, the prisoner's dilemma is not as sensitive to changes in the neighbourhood, and so the effect is
 274 much less pronounced. A single change in a player's neighbourhood will alter the probability of strategy
 275 imitation by large or small amounts depending on the strategy imitation rule chosen. If the payoff difference
 276 imitation rule is used the probability is highly dependent on the neighbourhood, whereas the Fermi imitation
 277 rule is much less sensitive. If the Fermi rule has a large noise parameter, this makes it more sensitive to a
 278 change in the neighbourhood.

279 We conclude that if the strategy imitation rule is strongly payoff dependent, and if the average payoffs
 280 of the game on the multiplex are greatly affected by the number of cooperators in a group, then updating
 281 layers asynchronously will lead to an increase in cooperation for a wide parameter range. This suggests that
 282 in order to model the evolution of cooperation robustly on multiplex networks, we must choose an updating
 283 strategy that most accurately captures the dynamics of the real-world scenario we hope to model.

284 In general, updating layers asynchronously rather than synchronously increases the amount of coopera-
 285 tion on a multiplex network. This is because the probability of one strategy imitating another is sensitive
 286 to the composition of a player's group. The more sensitive the probability of one player imitating another
 287 is to the composition of the neighbourhood, the larger the difference in cooperation between the two rules.
 288 The choice of asynchronous or synchronous updates also suggests a method of increasing cooperation on
 289 multilayer systems, namely to slow down the rate of updating the strategies on each layer compared to the

290 rate at which the payoffs are calculated. Specifically, this can be achieved with a move to asynchronous
291 updates rather than synchronous.

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