

# Optical Jamming Enhances the Secrecy Performance of the Generalized Space Shift Keying Aided Visible Light Downlink

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**Abstract**—In order to enhance the secrecy performance of the Generalized Space Shift Keying (GSSK) visible light communication (VLC) system, in this paper, an optical jamming aided secrecy enhancement scheme is proposed, in which the source transmitter (S) simultaneously sends both the confidential desired signal and optical jamming signals under amplitude and power constraints. The optical jamming signals obey the truncated Gaussian distribution for satisfying the constraints. Given the discrete set of channel inputs, the optical jamming aided GSSK-VLC system's secrecy performance is analyzed. Explicitly, the average mutual information (AMI), the lower bound of AMI and its closed-form approximation as well as the achievable secrecy rate are formulated analytically. Furthermore, the optimal power sharing strategy of the proposed GSSK-VLC systems relying on optical jamming is derived. Closed-form expressions are provided for the optimal power sharing in both the low- and high-SNR regions. Finally, extensive simulation results are presented to validate our analytical results.

**Index Terms**—Visible light communication (VLC), generalized space shift keying (GSSK), physical layer security (PLS), optical jamming, achievable secrecy rate, power sharing.

## I. INTRODUCTION

IN the face of the limited radio-frequency (RF) spectrum [1], visible light communication (VLC) relying on light-emitting diodes (LEDs) both for illumination and for data communications [2] has gained considerable attention from both academia and industry [2]–[4]. However, given the broadcast nature of the VLC downlink, they are inherently vulnerable to eavesdroppers, which are located in the illumination area of the transmitter LEDs. Therefore, similar to their RF counterparts, information privacy and confidentiality to the legitimate VLC users is an important issue, particularly when

the communicating nodes are deployed in public areas [5]. Additionally, securing VLC transmissions is also necessary owing to floor-to-door gaps, keyholes and partially covered windows as well as due to non-line-of-sight (LoS) reflections inside a room [6].

Traditionally, most cryptographic encryption and decryption methods are part of the upper layers of wireless systems [7]. However, these classic encryption techniques may be decrypted in the face of the ever-increasing computational power [8]. As a complement to the conventional upper-layer solutions, such as cryptographic techniques, various physical layer secrecy (PLS) techniques have been proposed to provide perfect security in wireless communication systems [9]–[12], which were first studied in the context of wiretap channels by Wyner in [13] for point-to-point communication systems and later for broadcast channels by Csiszár and Körner in [14], and for Gaussian wiretap channels in [15]. These contributions focused on maximizing the data rate of secret communications, namely, the secrecy capacity of a wiretap channel where a source transmitter (S) is equipped with multiple antennas for confidentially communicating with a legitimate destination receiver (D), while a passive eavesdropper (E) is trying to wiretap the confidential information of S and D. PLS techniques have been applied to a wide range of RF wireless systems, which has improved the overall security by complementing existing classic cryptography-based techniques [10]. However, given the average optical power, the peak optical power and the electrical power constraints imposed on the VLC signals, these techniques developed for RF wireless communication systems cannot be directly transplanted into practical VLC scenarios. Thus particular designs have to be proposed to avoid new vulnerabilities.

At the time of writing, most of the PLS aided VLC treatises considered single-input single-output (SISO) and multiple-input single-output (MISO) Gaussian wiretap channels. In a little more detail, the upper and lower capacity bounds of the modulation and direct detection (IM/DD) aided SISO VLC channel was investigated in [16]–[18]. As an extension, Chaaban *et al.* [19] developed the upper and lower bounds for the multiple-input multiple-output (MIMO) channel capacity under the idealized simplifying assumption of having perfect channel state information (CSI) at the transmitter. As one of the key techniques of achieving secrecy, the multi-LED based wiretap channel enjoyed particular attention [5], [20]–[23], where the associated high degree of freedom may be

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exploited by beamforming schemes [5], [22], [23] and artificial noise injection [20], [21], [24]. Specifically, the authors of [5] derived the lower and upper bounds on the capacity of the SISO Gaussian wiretap channel under continuous input distribution and amplitude constraints. Additionally, when the eavesdropper's channel was exactly known by S, the closed-form secrecy rate expressions were obtained for both MMSE and zero-forcing beamforming strategies. Then, robust secrecy-enhancing beamformers were designed for the secrecy rate maximization of a MISO VLC system, when either the CSI or the CSI-statistics of the eavesdropper was assumed to be perfectly known or imperfectly known [22], [23]. Alternatively, when considering MISO VLC systems, and when the eavesdropper's instantaneous CSI was not available at the transmitter, a friendly jammer was relied upon for transmitting jamming signals by a beamformer to maximize the secrecy rate [20], [21]. Furthermore, to enhance the communication secrecy, a scheme that jointly relies on both transmit beamforming and friendly jamming was proposed for a MISO VLC system operating in the face of multiple eavesdroppers [25]. Additionally, Zou et.al. in [26] have analyzed the secrecy rate of MISO optical wireless scattering communication systems, under certain constraints. Explicitly, a pair of secure communication protocols, termed as a non-jamming protocol and a cooperative jamming protocol, were proposed and tractable solutions were obtained based on the alternating optimization approach. A novel secure MIMO-VLC system has been designed using the modified Rivest-Shamir-Adleman (RSA) technique of [26] to encrypt the transmitted data in the MAC layer based on the location of the user supported by MIMO-VLC systems [26]. In [27], the closed-form analytical expressions of both secrecy outage probability and of the average secrecy capacity of a downlink VLC system have been derived by ignoring any amplitude constraints and by assuming Gaussian input distribution, while considering random positions for both the one and only legitimate user as well as for the multiple illegitimate users. By contrast, the analytical expressions of the exact and asymptotic secrecy outage probability have been derived in [28]. Light energy harvesting and random positions were assumed for the one and only legitimate user and for an illegitimate user in the context of a hybrid VLC-RF system. A secure system has been proposed in [29] for barcode-based VLC smartphones communicating over a full-duplex smartphone VLC channel.

A common assumption of these contributions is that the distributions of both the channel's input signals and of the jamming signals are continuous. In [5], [22] and [23], a continuous uniform distribution across a given interval was considered, while in [21] and [30] a truncated Gaussian distribution or its generalized form was assumed, where the design-objective was to increase the secrecy rate under a specific amplitude constraint of the input signal. However, the Gaussian input signal assumption is impractical due to its infinite peak power and excessive detection complexity. Additionally, it was found for RF communication that for the SISO Gaussian wiretap channel having both amplitude and power constraints, the channel's input distributions capable of achieving the secrecy capacity rely on finite support sets [31].

However, for the MISO Gaussian wiretap channel, subject to an amplitude constraint, the optimal input signal distributions capable of approaching the secrecy capacity are unknown. For the friendly jammer aided secrecy strategy of classic RF schemes it has been proven that the secrecy rate of any discrete channel inputs associated with finite alphabets is a non-concave function with respect to (w.r.t.) the total transmit power [32]. However, the associated secrecy rate of Gaussian channel inputs is concave and the optimal power allocation has a closed-form expression [33]. Another common feature of the above-mentioned contributions on artificial noise based friendly jamming is that the power sharing between the confidential desired signal and jamming signals has not as yet been considered in the literature, even though it is a critical issue.

The generalized space shift keying (GSSK) modulation technique [34]–[37] is eminently suited for the utilization of multiple LEDs relying on IM/DD techniques in VLC systems. As a result, GSSK aided VLC systems have been extensively studied. However, to the best of our knowledge, there are no achievable secrecy rate performance results in the literature for GSSK-VLC systems. Hence we propose a friendly jammer aided secrecy enhancement scheme and conceive the optimal power sharing between the desired confidential signals and jamming signals.

Explicitly, we consider the PLS issues of GSSK-VLC systems and improve their secrecy performance enhancement with the aid of friendly optical jamming. The contributions of this paper can be summarized as follows.

- A GSSK-VLC system communicating over a Gaussian wiretap channel is conceived, where S transmits its modulated signal jointly with optical jamming signals generated in the null space of the desired channel. The proposed secrecy enhancement scheme does not impose any change on the signals received by the destination, but the eavesdropper will suffer from the intentional optical jamming, regardless of its position. We adopt the truncated Gaussian distribution for the optical jamming signals to satisfy the amplitude constraint.
- The secrecy performance of the GSSK-VLC system relying on optical jamming is analyzed, which includes both the average mutual information (AMI), as well as the lower bound of AMI and its closed-form approximation and finally, the achievable secrecy rate. Additionally, closed-form approximations of the AMI of the Source-to-Destination (S-D) and the Source-to-Eavesdropper (S-E) links are derived, respectively. Furthermore, the influence of the optical jamming signal's amplitude constraint is discussed, which involves the derivation of the probability density function (PDF) of the sum of independent identical truncated Gaussian distributions, where an approximation is proposed to obtain a closed-form expression.
- The optimal power sharing between the confidential information and the jamming is determined for the sake of maximizing the achievable secrecy rate of the system under amplitude and power constraints. It is indeed plausible that by degrading the performance of E by jamming improves the secrecy, but naturally, the reduced

power assigned to the desired link degrades its BER. Specifically, the closed-form expressions of the optimal power sharing factors and the corresponding achievable secrecy rates are derived in both the low- and high signal-to-noise ratio (SNR) regions.

The remainder of this paper is organized as follows. The system and channel models are described in Section II. In Section III, our secrecy enhancement relying on optical jamming is proposed and the corresponding secrecy performance is analyzed. Then, in Section IV, the optimal power sharing is derived. Our simulation results and discussions are provided in Section V. Finally, we conclude in Section VI.

*Notation:* Vectors (matrices) are denoted by boldface lower-case (uppercase) letters. The set of  $N$ -dimensional real-valued (non-negative) numbers is denoted by  $\mathcal{R}^N (\mathcal{R}_+^N)$ .  $|\cdot|$ ,  $\|\cdot\|$ ,  $(\cdot)^T$ ,  $\lfloor \cdot \rfloor$ ,  $\odot$ ,  $\mathbb{E}\{\cdot\}$ ,  $\mathbb{I}(\cdot; \cdot)$ ,  $\binom{\cdot}{\cdot}$  denote absolute value, Euclidean norm, transposition, floor operation, Hadamard product, expected value, mutual information, and binomial coefficient, respectively. We use  $\mathbf{I}_N$  and  $\mathbf{1}$  to denote the  $N$ -dimensional identity matrix and the all-one column vector of length  $N$ , respectively. The curled inequality symbol  $\preceq$  between two vectors denotes componentwise inequality. Superscript  $[x]^+$  denotes  $\max\{x, 0\}$ . A lowercase letter  $x, y, z$  denotes one realization of the random variable  $X, Y, Z$ , respectively. We use the subscripts  $(\cdot)_D$  and  $(\cdot)_E$  to denote relevance to destination (D) and eavesdropper (E), respectively.

## II. SYSTEM DESCRIPTION AND CHANNEL MODELS

In this paper, we assume that the S is installed on the ceiling has  $N_t$  down-facing LEDs and communicates privately with D, which has a single up-facing photo-detector (PD). There is an eavesdropper E, who is also equipped with a single PD<sub>E</sub>. For simplicity, the PD<sub>E</sub> is also assumed to be facing vertically upwards. The parameters of all LEDs and PDs are assumed to be identical in this paper.

Therefore, the system considered represents a typical multiple-input single-output single-E (MISOSE) VLC Gaussian wiretap channel model. The signals received by D and E are expressed, respectively, as

$$y = \mathbf{h}_D^T \mathbf{x} + w_D, \quad (1)$$

$$z = \mathbf{h}_E^T \mathbf{x} + w_E, \quad (2)$$

where, by definition, we have  $\mathbf{h}_D = [h_{D,1}, h_{D,2}, \dots, h_{D,N_t}]^T \in \mathcal{R}_+^{N_t}$  and  $\mathbf{h}_E = [h_{E,1}, h_{E,2}, \dots, h_{E,N_t}]^T \in \mathcal{R}_+^{N_t}$ , which represent the channel gains of the S-D link and S-E link, respectively. We assume that S exploits the full knowledge of  $\mathbf{h}_D$  but no knowledge about  $\mathbf{h}_E$ . We assume that E can estimate its own channel gains  $\mathbf{h}_E$ . In (1) and (2),  $\mathbf{x} = [x_1, x_2, \dots, x_{N_t}]^T \in \mathcal{R}^{N_t}$  is an information-bearing signal vector sent by S, which is assumed to be superimposed on an identical direct current (DC) bias  $I_{DC} \in \mathcal{R}_+$  for adjusting the illumination level of LEDs. For safety reason and to avoid clipping distortion, we assume that the total current  $I_{DC} + x_{N_t}$  is restricted within the range of  $(I_{DC} - \alpha I_{DC}, I_{DC} + \alpha I_{DC})$ , where  $\alpha \in [0, 1]$  is termed as the modulation index [5], [16]. In other words,  $\mathbf{x}$  is subject to the amplitude constraint termed as  $|\mathbf{x}| \preceq A \mathbf{1}$ , where  $A = \alpha I_{DC}$ . Finally, in (1) and (2),  $w_D \sim \mathcal{N}(0, \sigma_D^2)$  and  $w_E \sim \mathcal{N}(0, \sigma_E^2)$  are zero-mean additive

white Gaussian noise (AWGN) processes, received by D and E, respectively.

When the LEDs are installed on the ceiling facing down, the channel model may neglect all non-LoS components, hence we only consider the LoS signal component for obtaining tractable analytical results [5], [22], [23], [38]. Then, assuming a generalized Lambertian emission pattern, the path gain  $G_{D,n_t}$  between the  $n_t$ -th LED and the PD<sub>D</sub> can be represented as [5], [39],

$$G_{D,n_t} = \begin{cases} \frac{1}{2\pi d_{n_t}^2} (\tilde{m} + 1) A_{Rx} \cos^{\tilde{m}}(\phi) \cos \psi_{n_t}, & |\psi_{n_t}| \leq \Psi_{FoV}, \\ 0, & |\psi_{n_t}| > \Psi_{FoV}, \end{cases} \quad (3)$$

where  $d_{n_t}$  is the LoS distance between the  $n_t$ -th LED and the PD<sub>D</sub>,  $\phi$  is the angle of irradiance of the LED,  $\psi_{n_t}$  is the angle of incidence of the  $n_t$ -th optical link, which is measured from the axis perpendicular to the receiver surface,  $\tilde{m} = -1/\log_2(\cos \Phi_{1/2})$  is the order of Lambertian emission with half irradiance at semi-angle  $\Phi_{1/2}$ , which is measured from the optical axis of the LED, and  $\Psi_{FoV}$  is the receiver's field-of-view (FoV) semi-angle. Finally,  $A_{Rx}$  is the effective detection area of the PD, which is given by [39]

$$A_{Rx} = \frac{\beta^2}{\sin^2(\Psi_{FoV})} A_{PD}, \quad (4)$$

where  $\beta$  is the refractive index of the optical concentrator and  $A_{PD}$  is the PD area. Then, the VLC channel gain between the  $n_t$ -th LED and the PD<sub>D</sub> can be expressed as  $h_{D,n_t} = TRG_{D,n_t}\eta$ ,  $n_t = 1, 2, \dots, N_t$ , where  $T$  is the gain of a transimpedance amplifier,  $R$  is the responsivity of the PD and  $\eta$  is the current-to-light conversion efficiency of the LEDs, respectively. Similarly, the channel gain between the  $n_t$ -th LED and the PD<sub>E</sub> can be expressed as  $h_{E,n_t} = TRG_{E,n_t}\eta$ , for  $n_t = 1, 2, \dots, N_t$ .

Note that the channel gain of a VLC link depends on the specific position of the transmitter LED and the receiver PD. If a receiver PD is not in a transmitter's FoV, the channel gain of the link will be zero [5].

## III. OPTICAL JAMMING AIDED GSSK-VLC SYSTEM AND ITS PERFORMANCE ANALYSIS

In this section, the performance of the GSSK-VLC system relying on optical jamming is considered. Firstly, the GSSK-VLC system and signal model are introduced. Then, the optical jamming aided GSSK-VLC scheme is designed for improving the secrecy performance. Finally, the secrecy performance of the optical jamming aided GSSK-VLC system is analyzed.

### A. GSSK-VLC System and Signals Modelling

We assume that there are  $N$  LEDs in the room considered, among which a subset of  $N_t$  LEDs are utilized for communication. Specifically, we assume that  $N_t$  out of  $N$  LEDs are employed for GSSK modulation. Then, during a specific symbol period,  $n_t$  ( $1 \leq n_t < N_t$ ) LEDs are activated to transmit an information symbol, while the remaining  $(N_t - n_t)$  LEDs are only used for illumination. Hence, there are in total  $M' = \binom{N_t}{n_t}$  possible combinations, among which  $M = 2^m$  with  $m = \lfloor \log_2 M' \rfloor$  are used transmitting  $m$  bits per symbol.

In our ensuring analysis, we explicitly assume that the first  $M$  combinations are used for conveying the information.

Let a GSSK symbols be expressed as  $x \in \mathcal{X}$ , where  $\mathcal{X}$  is the set of  $M$  possible GSSK symbols. Based on  $x$ ,  $n_t$  out of  $N_t$  LEDs are activated to transmit, with each having a constant intensity of  $I = s/\sqrt{n_t}$ . Correspondingly, the transmitted signal vector  $\mathbf{x}$  can be expressed as

$$\begin{aligned} \mathbf{x} &= \frac{s}{\sqrt{n_t}} \sum_{n_i=1}^{n_t} \mathbf{e}_{\omega_{n_i}} \\ &= \frac{s}{\sqrt{n_t}} \underbrace{[1 \cdots 0 \ 1 \ 0 \cdots 1 \cdots 0]^T}_{n_t \text{ out of } N_t \text{ non-zero elements}} \\ &= \frac{s}{\sqrt{n_t}} \mathbf{e}_\omega, \end{aligned} \quad (5)$$

where  $\mathbf{e}_{\omega_{n_i}}$ ,  $\omega_{n_i} \in \{1, 2, \dots, N_t\}$ , represents a specific column of an identity matrix  $\mathbf{I}_{N_t}$ , with the  $\omega_{n_i}$ -th element being one. Hence,  $\mathbf{e}_\omega = \sum_{n_i=1}^{n_t} \mathbf{e}_{\omega_{n_i}}$  is a  $N_t$ -length vector with  $n_t$  non-zero elements corresponding to the  $n_t$  activated LEDs, and  $\omega \in \Omega = \{1, 2, \dots, M\}$ . Without loss of generality, We assume that the average power constraint of  $\mathbf{x}$  is  $\mathbb{E}\{\|\mathbf{x}\|^2\} = \mathcal{P}$ , as well as that we have  $\mathbb{E}\{\|\mathbf{x}\|\} \leq A1$  for our peak amplitude constraint, where  $\mathcal{P}$  is the total power per transmission.

Note that, the GSSK-VLC system is reduced to a space shift keying (SSK) VLC system, when  $n_t = 1$ . Therefore, the SSK-VLC scheme is a special case of our GSSK-VLC. Hence, all the following analysis can be applied to SSK-VLC by letting  $n_t = 1$ .

When the signal of (5) is transmitted over the VLC wiretap channel, following (1) and (2), we have

$$y = \mathbf{h}_D^T \mathbf{x} + w_D = \frac{s}{\sqrt{n_t}} \mathbf{h}_D^T \mathbf{e}_\omega + w_D = h_{D(\omega)} s + w_D, \quad (6)$$

$$z = \mathbf{h}_E^T \mathbf{x} + w_E = \frac{s}{\sqrt{n_t}} \mathbf{h}_E^T \mathbf{e}_\omega + w_E = h_{E(\omega)} s + w_E, \quad (7)$$

where by definition, we have  $h_{D(\omega)} = \frac{\mathbf{h}_D^T \mathbf{e}_\omega}{\sqrt{n_t}} \in \mathcal{H}_{D(\omega)}$  and

$h_{E(\omega)} = \frac{\mathbf{h}_E^T \mathbf{e}_\omega}{\sqrt{n_t}} \in \mathcal{H}_{E(\omega)}$ , with  $\mathcal{H}_{D(\omega)}$  and  $\mathcal{H}_{E(\omega)}$  being the two sets collecting all the  $M$  possible channel states observed at D and E, respectively, when S transmits one of the  $M$  legitimate symbols.

### B. Optical Jamming Aided GSSK-VLC System

In this subsection, we enhance the secrecy performance of the S-D link by the optical jamming of E's reception, hence degrading its signal-to-interference-plus-noise ratio (S-INR). In a practical GSSK-VLC system, it is reasonable to assume that S does not have the CSI of the S-E link for a passive eavesdropper. However, S may transmit an optical jamming signal along with the modulated GSSK signal in the nullspace of the S-D channel. In principle, whilst S transmits a symbol using a set of  $n_t$  approximately selected LEDs, all the  $N_t$  LEDs can additionally be utilized to emit jamming signals without degrading the reception of D. In this case, D is capable of receiving its information as in the conventional GSSK system, while E experiences intentional interference. Consequently, the secrecy performance can be enhanced without degrading the reception of D. However, the

secrecy performance enhancement is achieved at the cost of activating more LEDs and by assigning additional power to the jamming signals, whilst maintaining the same total power. It should be noted that, although Zou et.al. in [40] have analyzed the secrecy rate for the MISO optical wireless scattering communication systems under certain constraints, there are three main differences between [40] and the present paper. 1) Instead of considering the non-LoS ultra-violet optical wireless scattering communication in [40], LoS indoor VLC is considered in this paper.; 2) Instead of considering the capacity of a Poisson channel based on the OOK modulation of [40], in this paper, we determine the capacity of the Gaussian wiretap channel for GSSK, which is eminently suited for multiple LEDs relying on IM/DD techniques in VLC systems; 3) Instead of knowing the accuracy of CSI of the S-E link as stipulated in [40], in this paper we assume that E is a passive eavesdropper and we have no CSI about the S-E link, which makes the considered scheme to be employable universally.

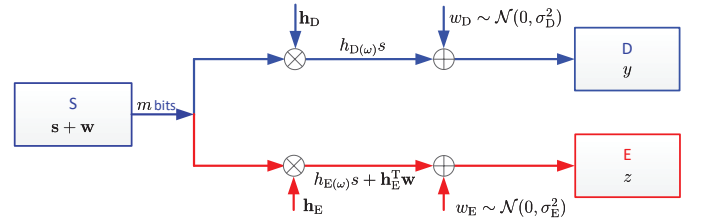


Fig. 1. System model of the GSSK-VLC Gaussian wiretap channel relying on optical jamming.

The system model of the GSSK-VLC system relying on optical jamming is illustrated in Fig. 1. With the aid of the classic singular value decomposition (SVD) [41], we can express  $\mathbf{h}_D$  as

$$\mathbf{h}_D^T = [\lambda, \mathbf{0}^T] [\mathbf{v}_s, \mathbf{V}_n]^T, \quad (8)$$

where  $\lambda$  is the singular value. From (8), we obtain a null space  $\mathbf{V}_n = [\mathbf{v}_1, \mathbf{v}_2, \dots, \mathbf{v}_{N_t-1}] \in \mathcal{R}^{N_t \times (N_t-1)}$ , since  $\mathbf{h}_D$  has  $\text{rank}(\mathbf{h}_D) = 1$ . Then, the optical jamming signals emitted by S can be designed to obey

$$\mathbf{w} = \mathbf{V}_n \mathbf{u}, \quad (9)$$

where  $\mathbf{u} = [u_1, u_2, \dots, u_{N_t-1}]^T \in \mathcal{R}^{N_t-1}$  is a time-varying jamming signal vector, whose entries are from a real-valued truncated Gaussian distribution [42] confined to the interval of  $[-\frac{A_2}{N_t-1}, \frac{A_2}{N_t-1}]$ . Thus, the peak amplitude constraint of the  $(N_t - 1)$  jamming codewords in every transmission is in the interval  $[-A_2, +A_2]$ . Then, we denote the peak amplitude constraint of the confidential signal as  $A_1$ , which is a constant for our GSSK-VLC system and can be expressed as  $A_1 = s\sqrt{n_t}$ . Thus, the whole peak amplitude constraint of the output intensity of the LEDs can be expressed as  $A = A_1 + A_2$ .

For a real-valued random variable defined as  $\tilde{P} \sim \mathcal{N}(\mu, \sigma^2)$ , we denote the double-sided truncated Gaussian random variable of  $\tilde{P}$  as  $P$ , whose PDF is expressed as

$\mathcal{TN}(\mu, \sigma^2, -B, B)$ , and is given by

$$f_P(p) = \frac{\frac{1}{\sqrt{2\pi\sigma^2}} \exp\left(-\frac{(p-\mu)^2}{2\sigma^2}\right)}{\Phi\left(\frac{B-\mu}{\sigma}\right) - \Phi\left(\frac{-B-\mu}{\sigma}\right)} I_{[-B, B]}(p), \quad (10)$$

where  $\Phi(\cdot)$  is the Gaussian cumulative distribution function (CDF),  $I_{[-B, B]}(p)$  is the indicator function satisfying  $I_{[-B, B]}(p) = 1$  if  $p \in [-B, B]$  and  $I_{[-B, B]}(p) = 0$ , otherwise. In our scheme, we set  $u_{n_i} \sim \mathcal{TN}(0, \frac{\sigma_J^2}{N_t-1}, -\frac{A_2}{N_t-1}, \frac{A_2}{N_t-1})$ ,  $n_i = 1, 2, \dots, N_t - 1$ , as seen in Fig. 1.

Upon involving the jamming signals, the signals transmitted from the  $N_t$  LEDs of the system can now be expressed as

$$\mathbf{x} = \mathbf{s} + \mathbf{w} = \frac{s}{\sqrt{n_t}} \mathbf{e}_\omega + \mathbf{V}_n \mathbf{u}. \quad (11)$$

In the proposed GSSK-VLC system, the total transmit power is constrained by

$$\mathbb{E}\{\|\mathbf{x}\|^2\} = \mathbb{E}\{\mathbf{s}^T \mathbf{s} + \mathbf{w}^T \mathbf{w}\} = s^2 + \sigma_J^2 = 1. \quad (12)$$

Hence, the optimal power can be assigned to the desired signals and to jamming by appropriately adjusting  $s^2$  and  $\sigma_J^2$  in (12). Consequently, the observations attained by D and E can be respectively expressed as

$$y = \mathbf{h}_D^T \mathbf{x} + w_D = \frac{s}{\sqrt{n_t}} \mathbf{h}_D^T \mathbf{e}_\omega + \mathbf{h}_D^T \mathbf{w} + w_D \\ = h_{D(\omega)} s + w_D, \quad (13)$$

$$z = \mathbf{h}_E^T \mathbf{x} + w_E = \frac{s}{\sqrt{n_t}} \mathbf{h}_E^T \mathbf{e}_\omega + \mathbf{h}_E^T \mathbf{w} + w_E \\ = h_{E(\omega)} s + \mathbf{h}_E^T \mathbf{V}_n \mathbf{u} + w_E, \quad (14)$$

where (13) is valid, because  $\mathbf{h}_D$  is orthogonal to  $\mathbf{w}$ . As (13) and (14) indicate, the secrecy performance of D and E is the same in the GSSK-VLC system considered, if no optical jamming is imposed on E's reception. However, when jamming is provided by  $\mathbf{w}$ , D's reception is not affected, apart from the allocation of a portion of the power to jamming, while E's reception may be severely degraded owing to this jamming, hence resulting in an improved secrecy performance for D, which is experimentally verified in the simulations of Section V-C.

### C. Secrecy Performance of Optical Jamming Aided GSSK-VLC System

In the following, we derive the AMI between S and E. Based on the results obtained, we can then evaluate the AMI between S and D by letting  $\mathbf{w} = \mathbf{0}$ . Let  $\check{w}_E$  denote the equivalent noise observed by E. Then, from (14), we have

$$\check{w}_E = \mathbf{h}_E^T \mathbf{w} + w_E = \mathbf{h}_E^T \mathbf{V}_n \mathbf{u} + w_E = \sum_{i=1}^{N_t-1} u_i \mathbf{h}_E^T \mathbf{v}_i + w_E, \quad (15)$$

which is the sum of  $N_t$  independent identically distributed (i.i.d.) variables. Since the resultant aggregate distribution is given by the convolution of this set of double-sided truncated Gaussian variables, the distribution of  $\check{w}_E$  is difficult to derive in a closed form. Even for the sum of two variables obeying the double-sided truncated Gaussian distribution, the resultant PDF becomes complicated. Specifically, let  $X_1 \sim \mathcal{TN}(\mu_{x_1}, \sigma_{x_1}, a_{x_1}, b_{x_1})$ ,  $X_2 \sim \mathcal{TN}(\mu_{x_2}, \sigma_{x_2}, a_{x_2}, b_{x_2})$

and  $f(x_i, \mu, \sigma) = \frac{1}{\sqrt{2\pi}\sigma} e^{-\frac{(x_i-\mu)^2}{2\sigma^2}}$ , as well as denote  $Z_1 = X_1 + X_2$ . Then, the PDF of  $Z_1$  is in the form of

$$f_{Z_1}(z_1) = \frac{\int_{\delta_1}^{\delta_2} f_1 f_2 dx_1}{\Phi_1 \Phi_2}, \quad (16)$$

where  $\delta_1 = \max\{a_{x_1}, z_1 - b_{x_2}\}$ ,  $\delta_2 = \min\{b_{x_1}, z_1 - a_{x_2}\}$ ,  $f_1 = f(x_1, \mu_{x_1}, \sigma_{x_1})$ ,  $f_2 = f(z_1 - x_1, \mu_{x_2}, \sigma_{x_2})$ ,  $\Phi_1 = \Phi(b_{x_1}, \mu_{x_1}, \sigma_{x_1}) - \Phi(a_{x_1}, \mu_{x_1}, \sigma_{x_1})$ ,  $\Phi_2 = \Phi(b_{x_2}, \mu_{x_2}, \sigma_{x_2}) - \Phi(a_{x_2}, \mu_{x_2}, \sigma_{x_2})$ . Furthermore, for deriving the PDF of  $\check{w}_E$ , we have to compute the convolution of  $f_{Z_1}(z_1)$  and the PDF of a Gaussian distributed  $w_E$ , which is excessively complex. Hence we invoke the central limit theorem, which states that the sum of  $n$  i.i.d. signals tends to the Gaussian distribution, as  $n$  becomes large. In practice, the number of transmit LEDs is sufficiently high, hence we can be confident that  $\check{w}_E$  has a near-Gaussian distribution.

It can be shown that the mean of  $\check{w}_E$  is zero, while its covariance obeys

$$\Omega_E = \mathbb{E}\{\check{w}_E \check{w}_E^T\} = \frac{\sigma_J^2}{N_t - 1} \mathbf{h}_E^T \mathbf{V}_n \mathbf{V}_n^T \mathbf{h}_E + \sigma_E^2, \quad (17)$$

Consequently, the received signal of E can be expressed with the aid of the Gaussian approximation as  $z = h_{E(\omega)} s + \check{w}_E$ . Furthermore, after normalizing  $z$  by multiplying  $\Omega_E^{-1/2}$ , we arrive at:

$$\tilde{z} = \Omega_E^{-1/2} h_{E(\omega)} s + \Omega_E^{-1/2} \check{w}_E = \tilde{h}_{E(\omega)} s + \tilde{w}_E, \quad (18)$$

where  $\tilde{w}_E$  has a zero mean and unit variance, and  $\tilde{h}_{E(\omega)} = \Omega_E^{-1/2} h_{E(\omega)}$ . Consequently, the conditional and unconditional PDFs of E's received signal  $\tilde{z}$  can be expressed as

$$f_{\tilde{z}|\tilde{h}}(\tilde{z}|\tilde{h} = \tilde{h}_{E(\omega)}) = \frac{1}{\sqrt{2\pi}} \exp\left(-\frac{(\tilde{z} - \tilde{h}_{E(\omega)} s)^2}{2}\right), \quad (19)$$

$$f_{\tilde{z}}(\tilde{z}) = \sum_{\omega \in \Omega} \frac{1}{\sqrt{2\pi} M} \exp\left(-\frac{(\tilde{z} - \tilde{h}_{E(\omega)} s)^2}{2}\right), \quad (20)$$

respectively. Similarly, the conditional and unconditional PDFs of  $y$  can be respectively expressed as

$$f_{Y|h_D}(y|h = h_{D(\omega)}) = \frac{1}{\sqrt{2\pi}\sigma_D} \exp\left(-\frac{(y - h_{D(\omega)} s)^2}{2\sigma_D^2}\right), \quad (21)$$

$$p_Y(y) = \sum_{\omega \in \Omega} p_{Y|h_D}(y|h_{D(\omega)}) P_{h_D}(h_{D(\omega)}) \\ = \sum_{\omega \in \Omega} \frac{1}{\sqrt{2\pi}\sigma_D M} \exp\left(-\frac{(y - h_{D(\omega)} s)^2}{2\sigma_D^2}\right). \quad (22)$$

According to [43], the AMI between two probability spaces remains invariant to a reversible transformation, implying that the transformation of (18) does not change the AMI. Consequently, given the PDF expressions of (19)-(22), the AMIs of both the S-D link and of the S-E link can be obtained, which are stated in Theorem 1.

**Theorem 1:** For the GSSK-VLC system relying on both optical jamming and on finite discrete inputs, the AMI of the S-D channel is

$$\mathbb{I}^J(h_D; Y) = \log_2 M - \frac{1}{M} \\ \times \sum_{\omega=1}^M \mathbb{E}_{w_D} \left[ \log_2 \sum_{\varpi=1}^M \exp\left(\frac{1}{2} \varrho_D (w_D^2 - (w_D + \zeta_{\omega, \varpi} s)^2)\right) \right], \quad (23)$$

where  $\zeta_{\omega,\varpi} = h_{D(\omega)} - h_{D(\varpi)}$ ,  $\varrho_D = 1/\sigma_D^2$ , and the AMI of S-E channel is

$$\mathbb{I}^J(h_E; Z) = \log_2 M - \frac{1}{M} \sum_{\omega=1}^M \mathbb{E}_{\tilde{w}_E} \left[ \log_2 \sum_{\varpi=1}^M \exp \left( \frac{1}{2} \left( \tilde{w}_E^2 - (\tilde{w}_E + \Omega_E^{-1/2} \xi_{\omega,\varpi} s)^2 \right) \right) \right] \quad (24)$$

where  $\xi_{\omega,\varpi} = h_{E(\omega)} - h_{E(\varpi)}$ .

*Proof:* Please refer to Appendix A. ■

Given  $\mathbb{I}^J(h_D; Y)$  and  $\mathbb{I}^J(h_E; Z)$ , as shown in (23) and (24), the secrecy rate achievable by D can be expressed as  $R_{\text{sec}}^J = [\mathbb{I}^J(h_D; Y) - \mathbb{I}^J(h_E; Z)]^+$ .

In order to simplify the results of (23) and (24), the lower bounds for the AMI of the S-D and S-E links can be derived, which are given by the following theorem.

**Theorem 2:** The AMI of the S-D channel can be lower bounded by

$$\mathbb{I}_L^J(h_D; Y) = \log_2 M - \frac{1}{2} (\log_2 e - 1) - \frac{1}{M} \sum_{\omega=1}^M \log_2 \sum_{\varpi=1}^M \exp \left( -\frac{(\zeta_{\omega,\varpi} s)^2}{4\sigma_D^2} \right). \quad (25)$$

The AMI of the S-E channel can be lower-bounded by

$$\mathbb{I}_L^J(h_E; Z) = \log_2 M - \frac{1}{2} (\log_2 e - 1) - \frac{1}{M} \sum_{\omega=1}^M \log_2 \sum_{\varpi=1}^M \exp \left( -\frac{(\xi_{\omega,\varpi} s)^2}{4\Omega_E} \right). \quad (26)$$

*Proof:* Please refer to Appendix B. ■

Furthermore, based on Theorem 1 and Theorem 2, below we derive the approximate expressions for  $\mathbb{I}^J(h_D; Y)$  and  $\mathbb{I}^J(h_E; Z)$ , respectively. Firstly, when  $\varrho_D \rightarrow \infty$  and  $\varrho_D \rightarrow 0$ , we can derive the limits of (23)  $\mathbb{I}^J(h_D; Y)$ , which are

$$\lim_{\varrho_D \rightarrow \infty} \mathbb{I}^J(h_D; Y) = \log_2 M \quad \text{and} \quad \lim_{\varrho_D \rightarrow 0} \mathbb{I}^J(h_D; Y) = 0. \quad (27)$$

Secondly, we can obtain the limits of  $\mathbb{I}_L^J(h_D; Y)$  from (25) for  $\varrho_D \rightarrow \infty$  and  $\varrho_D \rightarrow 0$  as

$$\lim_{\varrho_D \rightarrow \infty} \mathbb{I}_L^J(h_D; Y) = \log_2 M - \frac{1}{2} (\log_2 e - 1),$$

$$\lim_{\varrho_D \rightarrow 0} \mathbb{I}_L^J(h_D; Y) = -\frac{1}{2} (\log_2 e - 1). \quad (28)$$

When we compare the results in (27) and (28), we can see that there is a constant gap of  $\frac{1}{2} (\log_2 e - 1)$  between the AMI and its corresponding lower bound in both the high- and low-SNR regions. Furthermore, it can be shown that both  $\mathbb{I}^J(h_D; Y)$  and  $\mathbb{I}_L^J(h_D; Y)$  are monotonically increasing functions w.r.t.  $\varrho_D$ . It can be readily inferred from these observations that for any given SNR, especially for an SNR located in the high- or low-SNR region, the difference between  $\mathbb{I}^J(h_D; Y)$  and  $\mathbb{I}_L^J(h_D; Y)$  should approximately be a constant of  $\frac{1}{2} (\log_2 e - 1)$ . Based on similar arguments, we can also quantify the difference between  $\mathbb{I}^J(h_E; Z)$  and  $\mathbb{I}_L^J(h_E; Z)$ , which is also approximately a constant of  $\frac{1}{2} (\log_2 e - 1)$ .

Therefore, we may conclude that  $\mathbb{I}^J(h_D; Y)$  can be closely approximated as  $\mathbb{I}^J(h_D; Y) \approx \mathbb{I}_L^J(h_D; Y) + \frac{1}{2} (\log_2 e - 1)$ , and from Theorem 2, we have

$$\mathbb{I}_S^J(h_D; Y) \approx \log_2 M - \frac{1}{M} \sum_{\omega=1}^M \log_2 \sum_{\varpi=1}^M \exp \left( -\frac{(\zeta_{\omega,\varpi} s)^2}{4\sigma_D^2} \right). \quad (29)$$

Similarly, we have

$$\mathbb{I}_S^J(h_E; Z) \approx \log_2 M - \frac{1}{M} \sum_{\omega=1}^M \log_2 \sum_{\varpi=1}^M \exp \left( -\frac{(\xi_{\omega,\varpi} s)^2}{4\Omega_E} \right). \quad (30)$$

Consequently, by subtracting the approximation of the AMI between S and E  $\mathbb{I}_S^J(h_E; Z)$  from that of  $\mathbb{I}_S^J(h_D; Y)$  between S and D, the approximately achievable secrecy rate of D in our GSSK-VLC system relying on optical jamming can be expressed as

$$R_{S,\text{sec}}^J = [\mathbb{I}_S^J(h_D; Y) - \mathbb{I}_S^J(h_E; Z)]^+ = \left[ \frac{1}{M} \sum_{\omega=1}^M \log_2 \left[ \frac{\sum_{\varpi=1}^M \exp \left( -\frac{(\xi_{\omega,\varpi} s)^2}{4\Omega_E} \right)}{\sum_{\varpi=1}^M \exp \left( -\frac{(\zeta_{\omega,\varpi} s)^2}{4\sigma_D^2} \right)} \right] \right]^+ = \frac{1}{M} (\Upsilon_2 - \Upsilon_1), \quad (31)$$

$$\text{where, by definition, } \Upsilon_1 = \sum_{\omega=1}^M \log_2 \sum_{\varpi=1}^M \exp \left( -\frac{(\zeta_{\omega,\varpi} s)^2}{4\sigma_D^2} \right) \quad \text{and} \quad \Upsilon_2 = \sum_{\omega=1}^M \log_2 \sum_{\varpi=1}^M \exp \left( -\frac{(\xi_{\omega,\varpi} s)^2}{4\Omega_E} \right).$$

#### IV. OPTIMAL POWER SHARING

In order to achieve the highest secrecy rate possible for the proposed optical jamming aided GSSK-VLC system, there is an optimal power sharing between the information-bearing signal and jamming signals. Let the total power per transmission be denoted by  $\mathcal{P} = \mathcal{P}_1 + \mathcal{P}_2$ , where  $\mathcal{P}_1$  and  $\mathcal{P}_2$  are the average power constraint of the confidential signal and of the jamming signals, respectively. For the confidential signal, the average power constraint  $\mathcal{P}_1$  can be expressed as  $\mathcal{P}_1 = s^2$ . Then, for the jamming signals associated with a certain  $\sigma_J$ ,  $\mathcal{P}_2$  can be represented by  $A_2$  in a closed form as

$$\mathcal{P}_2 = (N_t - 1) \left[ 1 - \frac{\frac{2A_2^*}{\sigma_J^2} f \left( \frac{A_2^*}{\sigma_J^2} \right)}{2\Phi \left( \frac{A_2^*}{\sigma_J^2} \right) - 1} \right] \sigma_J^{*2}$$

$$= \left[ 1 - \frac{\frac{2A_2^*}{\sigma_J^2} f \left( \frac{A_2^*}{\sigma_J^2} \right)}{2\Phi \left( \frac{A_2^*}{\sigma_J^2} \right) - 1} \right] \sigma_J^2$$

$$= g(A_2), \quad (32)$$

where we have  $\sigma_J^{*2} = \frac{\sigma_J^2}{N_t - 1}$  and  $A_2^* = \frac{A_2}{N_t - 1}$ , while  $f$  and  $\Phi$  are defined as in Section III. Thus, for a certain  $s$  and  $\sigma_J$ , the relationship between  $\mathcal{P}$  and  $A$  can be expressed as

$$\mathcal{P} = \mathcal{P}_1 + \mathcal{P}_2 = s^2 + g(A_2) \leq s^2 + \sigma_J^2, \quad (33)$$

where (33) implies that the maximum average power per symbol required for the GSSK-VLC system considered is  $s^2 + \sigma_J^2$ , regardless of what the peak amplitude constraint is. Hence, we do not consider the peak amplitude constraint, when discussing the optimal power allocation problem. Without loss of generality, we assume that  $\sigma_D^2 = \sigma_E^2 = \sigma^2 = 1/\varrho$ . Then, if the specific portion of power allocated to the information-bearing signal is denoted as  $s^2 = \kappa \mathcal{P}$ , the associated jamming power becomes  $\sigma_J^2 = (1 - \kappa) \mathcal{P}$ .

In the proposed GSSK-VLC system, we assume that the channel  $h_{D(\omega)}$  is known to S. Hence, according to  $h_{D(\omega)}$ , S can

adapt the value of  $\kappa$  to achieve the optimal secrecy rate, which is hence referred to as the adaptive power sharing. Specifically, with the aid of the approximate closed-form expressions of (31) derived for the secrecy rate, the optimal value of  $\kappa$  set for achieving the highest secrecy rate at different SNRs can be found by a one-dimensional search, albeit at the cost of a relatively high complexity. Hence, below we analyze the behavior of the optimal value of  $\kappa$ , when the system is operated either in the low- or in the high-SNR region.

For convenience, we define  $d_\omega = \frac{1}{M} \sum_{\varpi=1}^M \zeta_{\omega,\varpi}^2$ , which represents the average value of the squared difference of the channel gain between S and D, when activating  $n_t$  LEDs. Similarly, we define  $g_\omega = \frac{1}{M} \sum_{\varpi=1}^M \xi_{\omega,\varpi}^2$ , representing the average value of the squared difference of the channel gain between S and E, when activating  $n_t$  LEDs. Furthermore, we define  $\bar{d} = \frac{1}{M} \sum_{\omega=1}^M d_\omega = \frac{1}{M^2} \sum_{\omega=1}^M \sum_{\varpi=1}^M \zeta_{\omega,\varpi}^2$  and  $\bar{g} = \frac{1}{M} \sum_{\omega=1}^M g_\omega = \frac{1}{M^2} \sum_{\omega=1}^M \sum_{\varpi=1}^M \xi_{\omega,\varpi}^2$ . Additionally, we define  $\chi = \mathbf{h}_{E,\omega}^T \mathbf{V}_n \mathbf{V}_n^T \mathbf{h}_{E,\omega}$ . With the aid of the above definitions, we first make the following observations.

**Theorem 3:** When the proposed optical jamming aided GSSK-VLC system is operated in the low-SNR region of  $\varrho \ll 1$  and when  $\bar{d} > \bar{g}$ , the value of  $\kappa$  maximizing  $R_{S,\text{sec}}^J$  is  $\kappa = 1$ . When  $\varrho \ll 1$  and  $\bar{d} < \bar{g}$ , we should maximize the jamming power by letting  $\kappa = 0$ , which results in a secrecy rate  $R_{S,\text{sec}}^J(\kappa)$  of 0.

*Proof:* Please refer to Appendix C. ■

Note that, when  $\bar{d} = \bar{g}$  and when operating in the low-SNR region, we also have  $R_{S,\text{sec}}^J = 0$ . In this case, the spatial distributions of D and E are identical when viewed from S.

From the proof of Theorem 3, we know that  $(1 - \frac{1}{4}\varrho PM\bar{g}\kappa) \geq 0$  and  $(1 - \frac{1}{4}\varrho PM\bar{d}\kappa) \geq 0$  should be satisfied. With the aid of these conditions, we can determine the region of low SNR. Based on our simulations, we found that provided that for  $\text{SNR} \leq 3$  dB, we satisfy the conditions required for Theorem 3.

As for the power-sharing in the high-SNR region, for an arbitrary  $\omega$ , we define  $\zeta_\omega^{\min} = \min_{\varpi \neq \omega, \varpi=1, \dots, M} \{\zeta_{\omega,\varpi}^2\}$ , and denote by  $\ell_\omega$  the total number of this kind of minimum elements. Furthermore, we define  $\zeta^{\min} = \min_{\omega=1, \dots, M} \{\zeta_\omega^{\min}\}$  and denote by  $\ell$  the total number of the minimum of  $\zeta^{\min}$ . Then, the optimum power-sharing is given by the following Theorem.

**Theorem 4:** When the proposed optical jamming aided GSSK-VLC system is operated in the high-SNR region of  $\varrho \gg 1$ , the optimal value of  $\kappa$  is

$$\kappa = \frac{4}{\varrho P \zeta^{\min}} \ln \left( \frac{\varrho P \zeta^{\min} \chi \ell}{M \bar{g}} - \ell \right), \quad (34)$$

which yields the maximum secrecy rate  $R_{S,\text{sec}}^J$  of

$$R_{S,\text{sec}}^J(\kappa) = \log_2 M + \frac{1}{M} \left[ \log_2 \left( 1 - \frac{1}{4\chi} \kappa M \bar{g} \right) - \log_2 \left[ 1 + \ell \exp \left( -\frac{\varrho P \zeta^{\min}}{4} \right) \right] \right]. \quad (35)$$

*Proof:* Please refer to Appendix D. ■

## V. SIMULATIONS AND NUMERICAL RESULTS

In this section, to validate the analytical secrecy performance and to demonstrate the efficiency of the proposed

TABLE I  
THE DISTRIBUTIONS OF THE LEDs' LOCATIONS

2 LEDs		8 LEDs	
LED	$(O_X, O_Y, O_Z)$		
1	(1.25, 2.50, 3.0) m	2	(3.75, 0.625, 3.0) m
2	(3.75, 2.50, 3.0) m	3	(1.25, 1.875, 3.0) m
4 LEDs		4	(3.75, 1.875, 3.0) m
1	(1.25, 1.25, 3.0) m	5	(1.25, 3.125, 3.0) m
2	(3.75, 1.25, 3.0) m	6	(3.75, 3.125, 3.0) m
3	(1.25, 3.75, 3.0) m	7	(1.25, 4.375, 3.0) m
4	(3.75, 3.75, 3.0) m	8	(3.75, 4.375, 3.0) m

TABLE II  
SIMULATION PARAMETERS

Simulation setup	
Room size ( $L \times W \times H$ )	$5 \times 5 \times 3$ m <sup>3</sup>
Number of LEDs	2, 4, 8
LEDs (D) height	3 m
Receivers (Bob and Eve) height	0.85 m
Transmitter parameters	
Semi-angle at half power ( $\Phi_{1/2}$ )	60°
Optical power/ electric conversion efficiency ( $\eta$ )	813.6 $\mu$ W/mA
Modulation index ( $\alpha$ )	0.1
Receiver parameters	
Refractive index ( $\beta$ )	1.5
Physical area of a PD ( $A_{\text{PD}}$ )	1.0 cm <sup>2</sup>
Receiver FoV semi-angle ( $\Psi_{\text{FoV}}$ )	60°
PD responsivity ( $R$ )	100 $\mu$ A/mW

optical jamming secrecy enhancement strategy, we provide numerical results for an indoor VLC environment having the dimensions of  $[5 \times 5 \times 3]$  m<sup>3</sup>, represented by a three-dimensional (3D) Cartesian coordinate system  $[O_X, O_Y, O_Z]$  with the origin being in one corner of the room. Again, the transmit LEDs are assumed to be perpendicular to the ceiling and down-facing to the floor. Similarly, the receivers (D and E) are located on the desks at the height of 0.85 m from the floor, which are assumed to be perpendicular to the desk and facing the ceiling. Unless specially noted, we assume that the positions of LEDs are those presented in Table I.

The half-illuminance semi-angle of LED  $\Phi_{1/2}$  is set to be 60°, which is a typical value for commercially-available high-brightness LEDs. Both D and E have a 60° FoV (semi-angle), the area of each PD is  $A_{\text{PD}} = 1.0$  cm<sup>2</sup> and the responsivity is  $R = 100$   $\mu$ A/mW/cm<sup>2</sup> [5]. For convenience, all the parameters involved in our simulations are summarized in Table II.

### A. Secrecy Performance of the Proposed GSSK-VLC Systems

To investigate the secrecy performance of the proposed optical jamming aided secrecy-enhancing scheme, a typical VLC scenario is considered, where we assume that D is located at (2.15, 1.28, 0.85) m. Unless specifically noted, we assume that the power allocated to the information signal and jamming signals is equal, i.e.,  $\kappa = 0.5$ . It should be noted that in practical applications, it is unreasonable to constrain D to a specific location. Actually, in our simulations, the D's location is selected randomly, hence the corresponding results are valid for all the areas that D can reach.

Fig. 2 depicts the AMI between S and E from (24) and its lower bound from (26) for the GSSK-VLC systems operating both with and without optical jamming, where  $N_t = 2, 4, 8$  for

SSK-VLC systems and  $N_t = 8, n_t = 2$  for GSSK-VLC system. E is located at (2.60, 0.88, 0.85) m, all other parameters involved in this simulation are taken from Table I and Table II. Observe from the simulation results that upon increasing the SNR, both  $\mathbb{I}(h_E; Z)$  and  $\mathbb{I}_L(h_E; Z)$  tend to constant values. Moreover,  $\mathbb{I}(h_E; Z)$  and  $\mathbb{I}_L(h_E; Z)$  also increase, as the number of LEDs  $N_t$  and that of the activated LEDs  $n_t$  is increased. Furthermore, the gap between  $\mathbb{I}(h_E; Z)$  and  $\mathbb{I}_L(h_E; Z)$  in the low- and high-SNR regions is approximately a constant of  $\frac{1}{2}(\log_2 e - 1)$ , which coincides with the theoretical analysis. It should be noted that the gap between  $\mathbb{I}(h_E; Z)$  and  $\mathbb{I}_L(h_E; Z)$  in the high-SNR region of all four cases is slightly higher than  $\frac{1}{2}(\log_2 e - 1)$ . This is due to the approximation error of using  $\mathcal{N}(0, 1)$  to estimate the distribution of  $\tilde{w}_E$ . Additionally, Fig. 2 reveals that the proposed optical jamming strategy is capable of dramatically decreasing the AMI between S and E for all the cases considered. In particular, for the GSSK scenario associated with  $N_t = 8, n_t = 2$ , when SNR = 40 dB, the AMI between S and E is  $\mathbb{I}(h_E; Z) = 3.55$  bits/symbol. After applying optical jamming, we have  $\mathbb{I}(h_E; Z) = 1.06$  bits/symbol, which is reduced substantially. Hence, all the proposed GSSK-VLC systems are capable of achieving an improved secrecy performance, when employing the optical jamming advocated.

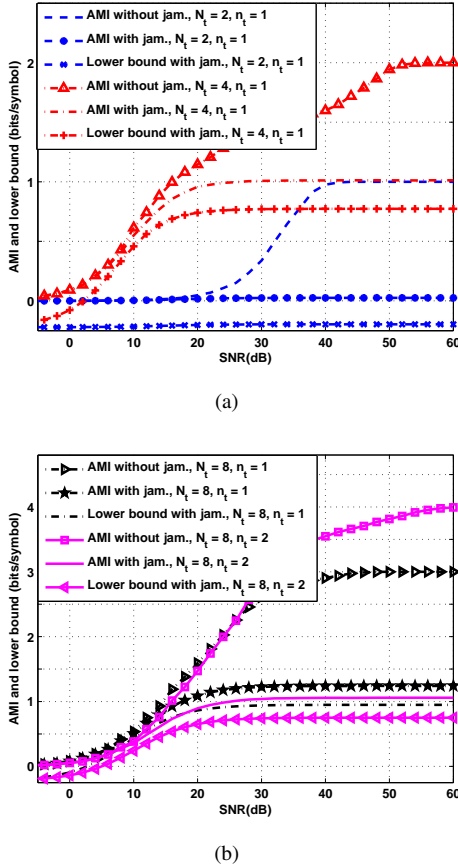


Fig. 2. AMI between S and E as well as its lower bound performance (a)  $N_t = 2, 4, n_t = 1$  in the SSK-VLC and (b)  $N_t = 8, n_t = 1, 2$  in the SSK- and GSSK-VLC systems with and without optical jamming. The results were calculated from (24) and (26).

Fig. 3 shows both the upper and lower bound on the achievable secrecy rate both in the low and high SNR regions. Explicitly, the upper bound is obtained as the upper bound of  $\mathbb{I}_B$  minus the lower bound of  $\mathbb{I}_E$ , while the lower bound is obtained as the lower bound of  $\mathbb{I}_B$  minus the upper bound of  $\mathbb{I}_E$ . Following this idea, when we rely on the proposed optical jamming approach and set  $\kappa = 0.5$ , the comparison of the achievable secrecy rate approximation to its upper and lower bounds is portrayed in Fig. 3, where D and E are located at (2.15, 1.28, 0.85) m and (2.60, 0.88, 0.85) m, respectively. Observe from Fig. 3 that in all the four cases, the achievable secrecy rate approximation is close to the upper and lower bound on the achievable secrecy rate except for the constant discrepancy of  $\frac{1}{2}(\log_2 e - 1)$ , especially in the low and high SNR regions, which coincides with the theoretical result.

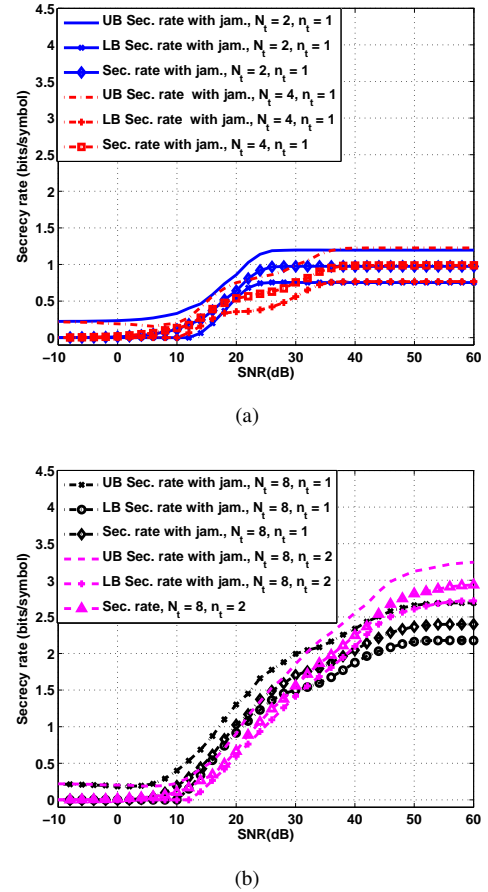


Fig. 3. The upper bound, lower bound and approximation of achievable secrecy rate performance between S and D (a)  $N_t = 2, 4, n_t = 1$  in the SSK-VLC system and (b)  $N_t = 8, n_t = 1, 2$  in the SSK- and GSSK-VLC systems with optical jamming. The results were calculated from (23), (25), (29) and (38).

Fig. 4 characterizes both the AMI and the achievable secrecy rate between S and D from (23) as well as that between S and E from (24) in Section III-C. The achievable secrecy rate of the VLC systems operating with and without optical jamming calculated from (31) is shown, where we have  $N_t = 2, 4, 8$  for our SSK-VLC systems as well as  $N_t = 8, n_t = 2$  for the GSSK-VLC system, and E is located at (2.60, 0.88, 0.85) m. All other parameters involved in this simulation are taken from

Table I and Table II. It is seen that in all the four cases, the achievable secrecy rate increases as the SNR increases, which is the explicit benefit of optical jamming.

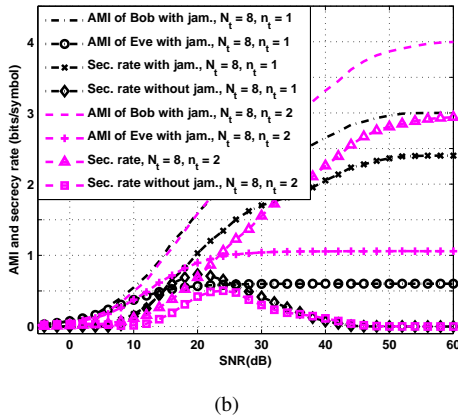
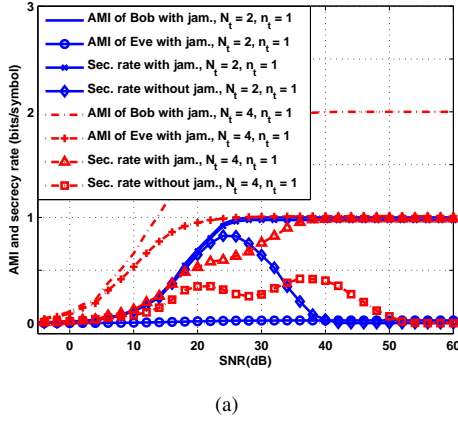


Fig. 4. AMI between S and D as well as that between S and E, and achievable secrecy rate performance (a)  $N_t = 2, n_t = 1$  in the SSK-VLC and (b)  $N_t = 8, n_t = 1, 2$  in the SSK- and GSSK-VLC systems with and without optical jamming. The results were calculated from (23), (24) and (31).

Fig. 5 demonstrates the achievable secrecy rate of the SSK-VLC system from (31) with optical jamming, where  $N_t = 4, n_t = 1$ . E's location is varied across the 2D plane at a height of 0.85 m, while D's position is fixed (2.15, 1.28, 0.85) m, all other parameters involved in this simulation are taken from Table I and Table II. The SNR is 30 dB. It can be observed from Fig. 5 that in most of the area considered, the SSK-VLC system achieves a relatively stable secrecy rate. However, as shown in [44], the achievable secrecy rate of the SSK-VLC system is zero in most of the scenarios considered, if no secrecy enhancement is utilized. The main reason behind this is that even for high SNRs, the detection performance of E can still be degraded by optical jamming, without affecting the reception performance of D, since the jamming signals are designed to lie in the null space of  $\mathbf{h}_D$ .

To make the above statement more convincing, in Fig. 6, we quantify the achievable secrecy rate of the GSSK-VLC system ( $N_t = 8, n_t = 2$ ) from (31) by letting SNR = 40 dB and by using the same parameters as in the above example. As shown in Fig. 6, the GSSK-VLC system achieves a relatively stable secrecy rate. As expected, when E moves close to D, the

achievable secrecy rate is significantly reduced. However, the achievable secrecy rate of the GSSK-VLC system increases rapidly, as E moves away from D. Additionally, when E is located at the symmetric regions of the projection of the transmitters, the achievable secrecy rate increases, which confirms the analytical results of [44].

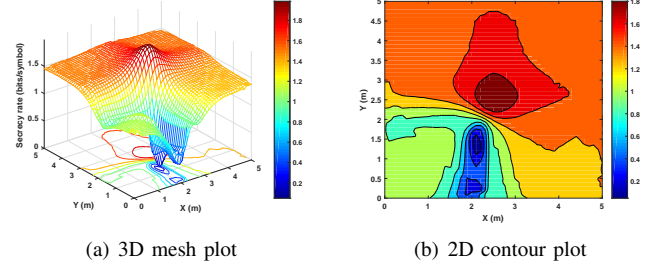


Fig. 5. Secrecy rate achieved by the SSK-VLC system with optical jamming. (a) 3D mesh plot; (b) 2D contour plot. The results were calculated from (31).

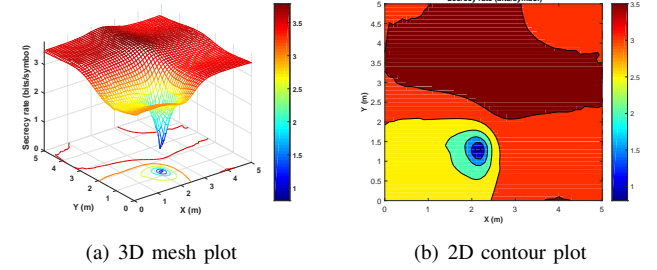


Fig. 6. Secrecy rate achieved by the GSSK-VLC system with optical jamming. (a) 3D mesh plot; (b) 2D contour plot. The results were calculated from (31).

In Fig. 7, the achievable secrecy rate of the GSSK-VLC system from (31) is depicted for the case, where  $N_t = 8, n_t = 2$ , and D's location is varied, while that of E is fixed at the location of (2.90, 1.88, 0.85) m. All other parameters involved in this simulation are taken from Table I and Table II. Observe from Fig. 7 that the secrecy rate of the GSSK-VLC system approaches zero, when D is close to E. Furthermore, we can also infer from the results of Fig. 7 that the achievable secrecy rate of the GSSK-VLC system has a relatively high correlation with D's location. When E's location is fixed, again there are symmetric regions exhibiting a low achievable secrecy rate.

### B. Optical Power Sharing

Fig. 8 shows the impact of different power sharing factors  $\kappa$  on the achievable secrecy rate of the GSSK-VLC system for low SNRs, the simulation results here are evaluated from Theorem 3, where  $N_t = 8, n_t = 1$ , while D is located at (2.15, 1.28, 0.85) m and E is located at (2.60, 1.88, 0.85) m, all other parameters involved in this simulation are taken from Table I and Table II. In this case, as implied by the theoretical results of Theorem 3, when  $\bar{d} > \bar{g}$ , the proposed GSSK-VLC system should allocate as much power to the information-bearing desired signal as possible, so as to achieve enhanced secrecy. Observe from Fig. 8(a) that for the low SNRs investigated, the achievable secrecy rate improves as the

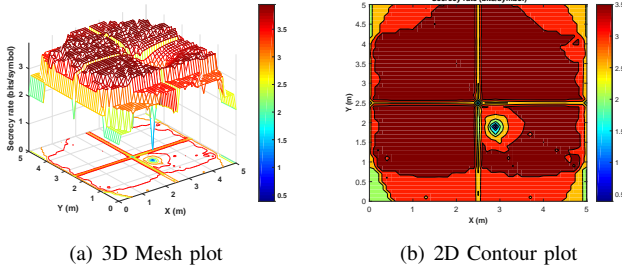


Fig. 7. Secrecy rate achieved by the GSSK-VLC system with optical jamming. (a) 3D mesh plot; (b) 2D contour plot. The results were calculated from (31).

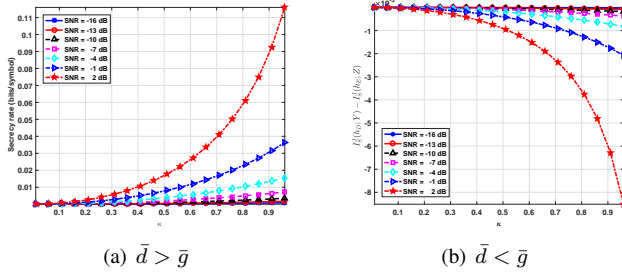


Fig. 8. (a) Achievable secrecy rate of the GSSK-VLC systems vs. the power sharing factor  $\kappa$  at low SNRs, when  $\bar{d} > \bar{g}$ ; (b)  $\mathbb{I}_S^J(h_D; Y) - \mathbb{I}_S^J(h_E; Z)$  results of the GSSK-VLC systems vs. the power sharing factor  $\kappa$  at low SNRs, where  $\bar{d} < \bar{g}$ . The results were calculated from Theorem 3 and (50).

power allocated to the information-bearing signal increases, which confirms the analytical results of Section IV. On the other hand, when  $\bar{d} < \bar{g}$ , Fig. 8(b) demonstrates the result of  $\mathbb{I}_S^J(h_D; Y) - \mathbb{I}_S^J(h_E; Z)$  for the GSSK-VLC system having different power sharing factors  $\kappa$  and low SNRs. As expected, in this case, even though an increased fraction of the power should be allocated to the jamming signals to enhance the system's secrecy, the achievable secrecy rate of the proposed GSSK-VLC system still remains zero. However, in practice, in order to avoid this unintended situation, we may adopt the user-centric LED allocation philosophy of [45] for ensuring that D always has better channel conditions than E, i.e.,  $\bar{d} > \bar{g}$ . From the results of Fig. 8(a) and Fig. 8(b), we conclude that in the low-SNR region, the power should be predominantly assigned to the information-bearing signal.

Fig. 9 characterizes the achievable secrecy rate of the proposed GSSK-VLC systems relying on the optimal power sharing obtained from (34) in the high-SNR region from Theorem 4. We assume that  $N_t = 2, 4, 8$  for our SSK-VLC systems and  $N_t = 8, n_t = 2$  for the GSSK-VLC system, where D is located at (2.15, 1.28, 0.85) m and E is located at (2.60, 1.88, 0.85) m, all other parameters involved in this simulation are taken from Table I and Table II. Observe from this figure that the optimal power sharing factor  $\kappa$  decreases almost logarithmically versus the SNR in the high-SNR region for all the scenarios investigated. From this figure, we can also observe that when a sufficiently high SNR is achieved by the proposed system,  $\kappa$  should be relatively small, so that the intended receiver D can successfully decode the confidential information, while E is unable to decode the confidential information due to the optical jamming. As shown in Fig.

9 for the same SNR, when the number of bits per GSSK symbol increases from 1 to 4, the optimal power sharing factor  $\kappa$  increases. This is because as  $N_t$  increases, the channel's correlation becomes higher, which will reduce the signal intensity received both by D and E. Consequently, the performance of the system will be degraded. In this case, in order to attain a better secrecy performance for the GSSK-VLC system, the transmitter has to allocate a higher proportion of the total power to the confidential information.

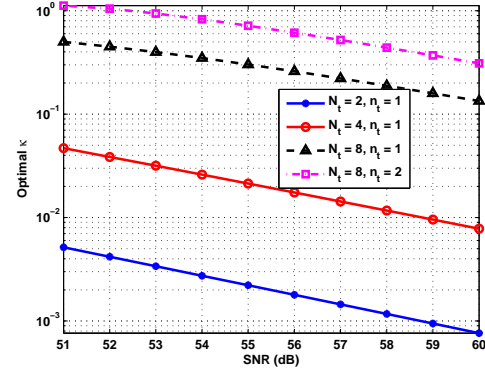


Fig. 9. Optimal power sharing factor  $\kappa$  for the optical jamming aided GSSK-VLC systems in the high-SNR region. The results were calculated from (34) and (35).

Fig. 10 depicts the optimal power sharing factor and the corresponding achievable secrecy rate obtained from (34) and (35) for our proposed GSSK-VLC systems at high SNRs from Theorem 4 vs. the SNR. Observe that the achievable secrecy rate  $R_{S,sec}^J$  reaches its maximum at the highest SNR and the smallest power sharing factor  $\kappa$  for the investigated four cases, where we have  $N_t = 2, 4, 8$  for the SSK-VLC systems and  $N_t = 8, n_t = 2$  for the GSSK-VLC system, while D is located at (2.15, 1.28, 0.85) m and E is located at (2.60, 1.88, 0.85) m, all other parameters involved in this simulation are taken from Table I and Table II. We can also observe that  $R_{S,sec}^J$  approaches its minimum, when the SNR is relatively low and a small proportion of the power is allocated to transmit jamming in our SSK-VLC system having  $N_t = 2, 4, 8$ . Especially, for our GSSK-VLC system associated with  $N_t = 8, n_t = 2$ ,  $R_{S,sec}^J$  approaches its minimum when  $\kappa = 1$  and when the SNR is relatively low. The reason behind this is that as both  $N_t$  and  $n_t$  increase, the channel's correlation becomes higher, which will reduce the signal intensity received at D when SNR is relatively low. Additionally, in all the above-mentioned four cases,  $R_{S,sec}^J$  reaches its maximum, when the SNR is relatively high and almost more than half of the power is allocated to transmit optical jamming, i.e. we have  $\kappa \leq 0.5$ .

From the results shown in Fig. 11, we can observe that for all the considered four GSSK-VLC systems operating in the high-SNR region, when  $\kappa = 1$ , the achievable secrecy rate approaches its minimum. By contrast, when  $\kappa = 0.1$ , the system attains the best secrecy performance within the SNR range considered. We observe furthermore that, 1) For the SSK-VLC system associated with  $N_t = 2, n_t = 1$ , when  $\kappa$  varies from 1 to 0.7, the achievable secrecy rate can be

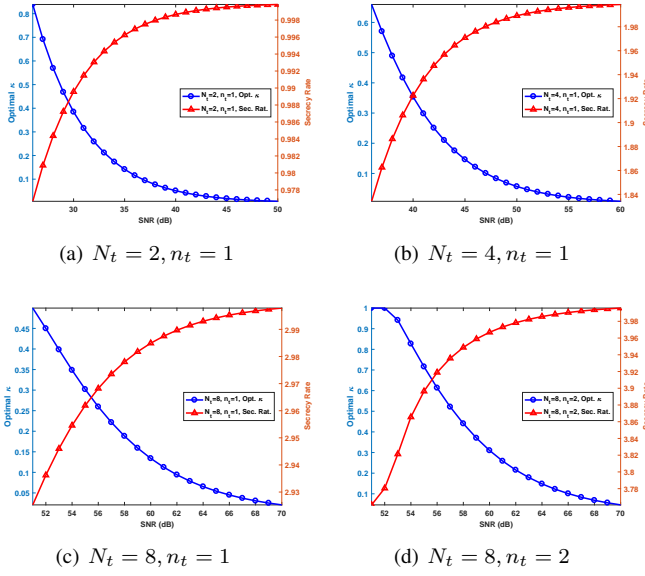


Fig. 10. Optimal power sharing factor  $\kappa$  and corresponding achievable secrecy rate versus SNR at high-SNR region. (a)  $N_t = 2, n_t = 1$ ; (b)  $N_t = 4, n_t = 1$ ; (c)  $N_t = 8, n_t = 1$ ; (d)  $N_t = 8, n_t = 2$ . The results were calculated from (34) and (35).

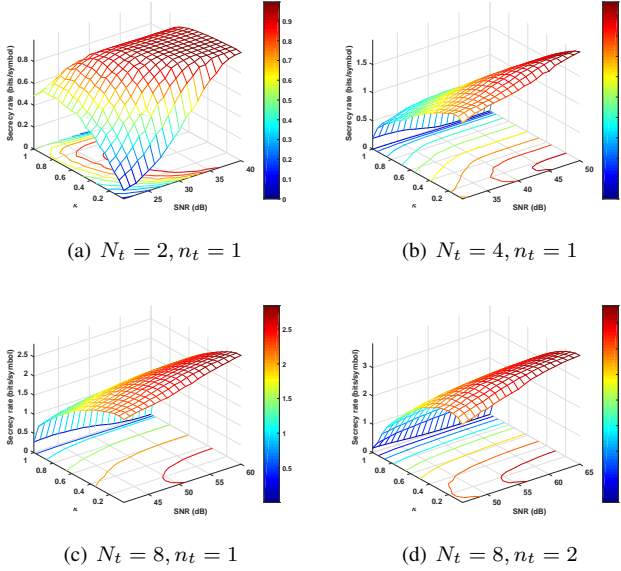


Fig. 11. Secrecy performance of the GSSK-VLC systems considered vs  $\kappa$  and SNR. (a)  $N_t = 2, n_t = 1$ ; (b)  $N_t = 4, n_t = 1$ ; (c)  $N_t = 8, n_t = 1$ ; (d)  $N_t = 8, n_t = 2$ . The results were calculated from (34) and (35).

increased by 1 bit/symbol; 2) For the SSK-VLC system having  $N_t = 4, n_t = 1$ , when  $\kappa$  varies from 1 to 0.5, the achievable secrecy rate can be increased by 1.5 bits/symbol; 3) For the SSK-VLC system having  $N_t = 8, n_t = 1$ , when  $\kappa$  varies from 1 to 0.5, the achievable secrecy rate can be increased by 2.5 bits/symbol; 4) For the GSSK-VLC system using  $N_t = 8, n_t = 2$ , when  $\kappa$  varies from 1 to 0.5, the achievable secrecy rate can be increased by 3 bits/symbol.

From Fig. 11, we can further conclude that the achievable secrecy rate increases, as more power is assigned to the jamming signals, *i.e.* for smaller  $\kappa$ . However, when a low power is allocated to the confidential information signal, the BER of the

S-D link will be degraded. The trade-off between the secrecy performance and the BER performance should be carefully considered for each application. Based on our analysis, a look-up table can be constructed to guide the system design of the optical jamming aided GSSK-VLC systems, so that the system parameters can be optimal by selected for the system considered. If the system is secrecy-critical, we may opt for a relatively low power sharing factor  $\kappa$  based on Theorem 3 and Theorem 4. Otherwise, if the BER performance is the most important metric of the system considered, then more power should be allocated to the information signals. For the GSSK-VLC systems jointly considering the secrecy and BER performance, based on the above results, we can opt for a power sharing factor of  $\kappa = 0.5$  in the relatively high-SNR region and  $\kappa = 1$  in the lower-SNR region.

## VI. CONCLUSIONS

As a recent secrecy enhancement strategy, PLS has been shown to be unbreakable, regardless of the computational capability of E. Given the broadcast nature of the VLC downlink, it is advisable to improve its secrecy. In this paper, we have provided the secrecy performance analysis of a PLS-aided GSSK-VLC system. Accordingly, four major contributions have been proposed. Firstly, by exploiting the input signal characteristics and channels of the proposed GSSK-VLC system, the secrecy performance was analyzed, when the input signals are assumed to have finite discrete distributions subject to specific amplitude and power constraints. From these results, we conclude that without extra secrecy enhancement, E may wiretap the confidential signals at high SNRs, even if its channel conditions are worse than those of D. Moreover, if the S-D channel is degraded, the system fails to support secret communication. Secondly, a friendly optical jamming aided secrecy enhancement scheme was designed for the proposed GSSK-VLC system. Apart from transmitting optical jamming signals by the LEDs, S sends simultaneously its confidential signal using these LEDs under appropriate amplitude and power constraints. We adopted the truncated Gaussian distribution for the optical jamming signals to satisfy these constraints. Furthermore, the optical jamming signals were generated within the nullspace of the S-D channel vector. Thirdly, the secrecy performance of our GSSK-VLC system relying on optical jamming was analyzed. Fourthly, the optimal power sharing strategy of the proposed GSSK-VLC system using optical jamming was considered for maximizing the achievable secrecy rate of the proposed system. Specifically, the closed-form expressions of the optimal power sharing were derived both for the low- and high-SNR regions. Finally, all the analytical results have been verified by computer simulations.

## APPENDIX

## A. Proof of Theorem 1

Based on (19)-(22), the AMI  $\mathbb{I}^J(h_D; Y)$  can be expressed as

$$\begin{aligned} \mathbb{I}^J(h_D; Y) &= \sum_{\omega=1}^M \int_y f_{Y|h_D}(h = h_{D(\omega)}, y) \\ &\quad \times \log_2 \frac{f_{Y|h_D}(y|h_D = h_{D(\omega)})}{p_Y(y)} dy \\ &= \log_2 M - \frac{1}{M} \\ &\quad \times \sum_{\omega=1}^M \mathbb{E}_{w_D} \left[ \log_2 \sum_{\varpi=1}^M \exp \left( \frac{w_D^2 - (w_D + (h_{D(\omega)} - h_{D(\varpi)})s)^2}{2\sigma_D^2} \right) \right]. \end{aligned} \quad (36)$$

Using the notation  $\zeta_{\omega, \varpi} = h_{D(\omega)} - h_{D(\varpi)}$ , (36) can be expressed as

$$\begin{aligned} \mathbb{I}(h_D; Y) &= \log_2 M - \frac{1}{M} \\ &\quad \times \sum_{\omega=1}^M \mathbb{E}_{w_D} \left[ \log_2 \sum_{\varpi=1}^M \exp \left( \frac{1}{2} \varrho_D (w_D^2 - (w_D + \zeta_{\omega, \varpi})^2) \right) \right], \end{aligned} \quad (37)$$

where  $\varrho_D = 1/\sigma_D^2$ . It can be shown that for a particular  $\zeta_{\omega, \varpi}$  and  $s$ ,  $\mathbb{I}(h_D; Y)$  is a monotonically increasing function w.r.t. the SNR  $\varrho_D$ . For  $\varrho_D \rightarrow \infty$ , i.e.,  $\sigma_D^2 = 0$ , we have

$$\lim_{\varrho_D \rightarrow \infty} \mathbb{I}(h_D; Y) = \log_2 M, \quad (38)$$

which is the upper bound of  $\mathbb{I}(h_D; Y)$ .

In a similar way, when using the notation of  $\xi_{\omega, \varpi} = h_{E(\omega)} - h_{E(\varpi)}$ , the AMI of the S-E channel can be expressed as

$$\begin{aligned} \mathbb{I}^J(h_E; Z) &= \log_2 M - \frac{1}{M} \\ &\quad \times \sum_{\omega=1}^M \mathbb{E}_{\tilde{w}_E} \left[ \log_2 \sum_{\varpi=1}^M \exp \left( \frac{1}{2} \left( \tilde{w}_E^2 - (\tilde{w}_E + \Omega_E^{-1/2} \xi_{\omega, \varpi} s)^2 \right) \right) \right]. \end{aligned} \quad (39)$$

## B. Proof of Theorem 2

The proof of (25) and that of (26) are the same. Therefore, in the following, we only detail the proof of (26).

From (24), we have

$$\begin{aligned} \mathbb{I}^J(h_E; Z) &= \log_2 M - \frac{1}{M} \sum_{\omega=1}^M \mathbb{E}_{\tilde{w}_E} \left[ \log_2 \exp \left( \frac{\tilde{w}_E^2}{2} \right) \right] - \frac{1}{M} \\ &\quad \times \sum_{\omega=1}^M \mathbb{E}_{\tilde{w}_E} \left[ \log_2 \sum_{\varpi=1}^M \exp \left( -\frac{(\tilde{w}_E + \Omega_E^{-1/2} \xi_{\omega, \varpi} s)^2}{2} \right) \right] \\ &= \log_2 M - I_1^J - I_2^J. \end{aligned} \quad (40)$$

The second term at the RHS of (40), i.e.,  $I_1^J$ , can be simplified as

$$\begin{aligned} I_1^J &= \frac{1}{M} \sum_{\omega=1}^M \mathbb{E}_{\tilde{w}_E} \left[ \log_2 \exp \left( \frac{\tilde{w}_E^2}{2} \right) \right] \\ &= \log_2 e \mathbb{E}_{\tilde{w}_E} \left[ \frac{\tilde{w}_E^2}{2} \right] = \frac{1}{2} \log_2 e. \end{aligned} \quad (41)$$

Due to the concavity of  $\log_2(\cdot)$ , the third term at the RHS of (40), i.e.,  $I_2^J$ , can be upper bounded by applying Jensen's inequality as

$$\begin{aligned} I_2^J &\leq \frac{1}{M} \sum_{\omega=1}^M \log_2 \sum_{\varpi=1}^M \mathbb{E}_{\tilde{w}_E} \left[ \exp \left( -\frac{(\tilde{w}_E + \Omega_E^{-1/2} \xi_{\omega, \varpi} s)^2}{2} \right) \right] \\ &= -\frac{1}{2} + \frac{1}{M} \sum_{\omega=1}^M \log_2 \sum_{\varpi=1}^M \exp \left( -\frac{(\xi_{\omega, \varpi} s)^2}{4\Omega_E} \right). \end{aligned} \quad (42)$$

Finally, upon substituting (41) and (42) into (40), we can arrive at:

$$\begin{aligned} \mathbb{I}^J(h_E; Z) &\geq \log_2 M - \frac{1}{2} (\log_2 e - 1) \\ &\quad - \frac{1}{M} \sum_{\omega=1}^M \log_2 \sum_{\varpi=1}^M \exp \left( -\frac{(\xi_{\omega, \varpi} s)^2}{4\Omega_E} \right), \end{aligned} \quad (43)$$

which completes the proof of Theorem 2.

## C. Proof of Theorem 3

Let us first consider the situation, where the SNR is very low, that is  $\varrho \ll 1$ . Then, from (29), we have

$$\begin{aligned} \Upsilon_1 &= \sum_{\omega=1}^M \log_2 \sum_{\varpi=1}^M \exp \left( -\frac{(\zeta_{\omega, \varpi} s)^2}{4\sigma^2} \right) \\ &\stackrel{(a)}{\approx} \sum_{\omega=1}^M \log_2 \sum_{\varpi=1}^M \left( 1 - \frac{\varrho \kappa \mathcal{P} \zeta_{\omega, \varpi}^2}{4} \right) \\ &\stackrel{(b)}{\approx} M \log_2 M + \log_2 \left( 1 - \frac{\varrho \kappa \mathcal{P}}{4} \sum_{\omega=1}^M d_{\omega} \right) \\ &= M \log_2 M + \log_2 \left( 1 - \frac{1}{4} \varrho \kappa \mathcal{P} M \bar{d} \right), \end{aligned} \quad (44)$$

where we have (a) by applying the approximation by the Taylor series. For small  $\varrho$ , we have (b) by letting  $\varrho \rightarrow 0$ .

Similarly, from (30), we have

$$\begin{aligned} \Upsilon_2 &= \sum_{\omega=1}^M \log_2 \sum_{\varpi=1}^M \exp \left( -\frac{\varrho \kappa \mathcal{P} \xi_{\omega, \varpi}^2}{4(\varrho(1-\kappa)\mathcal{P}\chi + 1)} \right) \\ &\stackrel{(a)}{\approx} \sum_{\omega=1}^M \log_2 \sum_{\varpi=1}^M \exp \left( -\frac{\varrho \kappa \mathcal{P} \xi_{\omega, \varpi}^2}{4} \right) \\ &\stackrel{(b)}{\approx} \sum_{\omega=1}^M \log_2 \sum_{\varpi=1}^M \left( 1 - \frac{\varrho \kappa \mathcal{P} \xi_{\omega, \varpi}^2}{4} \right) \\ &= \sum_{\omega=1}^M \log_2 \left( M \left( 1 - \frac{\varrho \kappa \mathcal{P} g_{\omega}}{4} \right) \right) \\ &\stackrel{(c)}{\approx} M \log_2 M + \log_2 \left( 1 - \frac{\varrho \kappa \mathcal{P}}{4} \sum_{\omega=1}^M g_{\omega} \right) \\ &= M \log_2 M + \log_2 \left( 1 - \frac{1}{4} \varrho \kappa \mathcal{P} M \bar{g} \right), \end{aligned} \quad (45)$$

where again, we have (a) and (c) due to  $\varrho \ll 1$  and we have (b) by the Taylor expansion and approximation. By substituting (44) and (45) into (31), we obtain

$$I_{S, \text{sec}}^J(\kappa) = \mathbb{I}_S^J(h_D; Y) - \mathbb{I}_S^J(h_E; Z) = \frac{1}{M} (\Upsilon_2 - \Upsilon_1) \quad (46)$$

$$= \frac{1}{M} \log_2 \frac{1 - \frac{1}{4} \varrho \mathcal{P} M \bar{g} \kappa}{1 - \frac{1}{4} \varrho \mathcal{P} M \bar{d} \kappa}. \quad (47)$$

We can readily verify that when  $\bar{d} > \bar{g}$ ,  $R_{S,\text{sec}}^J(\kappa)$  is a monotonically increasing function of  $\kappa \in [0, 1]$ . By contrast, when  $\bar{d} < \bar{g}$ ,  $R_{S,\text{sec}}^J(\kappa)$  is a monotonically decreasing function of  $\kappa \in [0, 1]$ . Therefore, when operated in the low-SNR region, and when  $\bar{d} > \bar{g}$ ,  $R_{S,\text{sec}}^J(\kappa)$  attains its maximum at  $\kappa = 1$ . When  $\bar{d} < \bar{g}$ ,  $R_{S,\text{sec}}^J(\kappa)$  achieves its maximum at  $\kappa = 0$ , implying that we should allocate as much power as possible to the jamming signals. It should be noted that the achievable secrecy rate  $R_{S,\text{sec}}^J(\kappa)$  is zero, when  $\bar{d} < \bar{g}$  and  $\rho \ll 1^1$ .

#### D. Proof of Theorem 4

Let us now consider the situation of the high-SNR region with  $\rho \gg 1$ . Due to value ranges of  $\rho, \kappa$ , one of the following three cases should be satisfied by the product of  $\rho\kappa$ , which are 1)  $\rho\kappa \rightarrow 0$ ; 2)  $\rho\kappa = c_1$ , where  $c_1$  is a finite real-valued constant; 3)  $\rho\kappa \rightarrow \infty$ .

In the first case, (31) is used for calculating  $R_{S,\text{sec}}^J(\kappa)$ , which implies that  $\kappa \rightarrow 0$  and gives  $R_{S,\text{sec}}^J(\kappa) = 0$ . Hence, there exist no values of  $\kappa \in (0, 1)$  that are optimal. In the second case,  $\rho \rightarrow \infty$  also results in  $\kappa \rightarrow 0$ , hence we have  $R_{S,\text{sec}}^J(\kappa) = 0$ . In the third case, when  $\rho\kappa \rightarrow \infty$ , we can infer from (31) that  $R_{S,\text{sec}}^J(\kappa) \approx \frac{1}{M} \sum_{\omega=1}^M \log_2 \sum_{\varpi=1}^M \exp\left(-\frac{\kappa \mathcal{P} \xi_{\omega,\varpi}^2}{4(1-\kappa)\mathcal{P}\chi}\right) + c$ , where  $c$  is a constant. Explicitly, the highest  $R_{S,\text{sec}}^J(\kappa)$  is obtained with an optimal  $\kappa$  approaching 0. From the analysis of the above three cases, we can conclude that the optimal value of  $\kappa$  is close to zero, when operating in the high-SNR region. In other words, in the high-SNR region, the optimal  $\kappa$  should be a relative small value. Based on these observations, we define the optimal  $\kappa$  in the high-SNR region as follows.

We commence by considering  $\Upsilon_1$  in the high-SNR region, which can be approximated as

$$\begin{aligned} \Upsilon_1 &= \sum_{\omega=1}^M \log_2 \sum_{\varpi=1}^M \exp\left(-\frac{(\zeta_{\omega,\varpi}s)^2}{4\sigma^2}\right) \\ &= \sum_{\omega=1}^M \log_2 \left[1 + \sum_{\varpi=1, \varpi \neq \omega}^M \exp\left(-\frac{\rho\kappa \mathcal{P} \zeta_{\omega,\varpi}^2}{4}\right)\right]. \end{aligned} \quad (48)$$

Since  $\exp(-\delta)$  is a rapidly decaying function w.r.t. a positive  $\delta$ , we can introduce the approximation of

$$\sum_{\varpi=1, \varpi \neq \omega}^{N_i} \exp\left(-\frac{\rho\kappa \mathcal{P} \zeta_{\omega,\varpi}^2}{4}\right) \approx \ell_{\omega} \exp\left(-\frac{\rho\kappa \mathcal{P} \zeta_{\omega}^{\min}}{4}\right). \quad (49)$$

Then,  $\Upsilon_1$  can be approximated as

$$\begin{aligned} \Upsilon_1 &\approx \sum_{\omega=1}^M \log_2 \left[1 + \ell_{\omega} \exp\left(-\frac{\rho\kappa \mathcal{P} \zeta_{\omega}^{\min}}{4}\right)\right] \\ &= \log_2 \prod_{\omega=1}^M \left[1 + \ell_{\omega} \exp\left(-\frac{\rho\kappa \mathcal{P} \zeta_{\omega}^{\min}}{4}\right)\right] \\ &\stackrel{(a)}{\approx} \log_2 \left[1 + \sum_{\omega=1}^M \ell_{\omega} \exp\left(-\frac{\rho\kappa \mathcal{P} \zeta_{\omega}^{\min}}{4}\right)\right] \\ &\stackrel{(b)}{\approx} \log_2 \left[1 + \ell \exp\left(-\frac{\rho\kappa \mathcal{P} \zeta^{\min}}{4}\right)\right], \end{aligned} \quad (50)$$

<sup>1</sup>In practical systems, we should use the resource-allocation to ensure that  $\bar{d} > \bar{g}$ .

where both (a) and (b) hold due to the fact that  $\kappa$  is a small real number approaching 0.

The term  $\Upsilon_2$  defined under (31) in the high-SNR region can be approximated as

$$\begin{aligned} \Upsilon_2 &\stackrel{(a)}{\approx} \sum_{\omega=1}^M \log_2 \sum_{\varpi=1}^M \exp\left(-\frac{\kappa \xi_{\omega,\varpi}^2}{4(1-\kappa)\chi}\right) \\ &\stackrel{(b)}{\approx} \sum_{\omega=1}^M \log_2 \sum_{\varpi=1}^M \exp\left(-\frac{\kappa \xi_{\omega,\varpi}^2}{4\chi}\right) \\ &\stackrel{(c)}{\approx} \sum_{\omega=1}^M \log_2 \sum_{\varpi=1}^M \left(1 - \frac{\kappa \xi_{\omega,\varpi}^2}{4\chi}\right) \\ &= \sum_{\omega=1}^M \log_2 \left(M \left(1 - \frac{\kappa g_{\omega}}{4\chi}\right)\right) \\ &\approx M \log_2 M + \log_2 \left(1 - \frac{\kappa}{4\chi} \sum_{\omega=1}^M g_{\omega}\right) \\ &= M \log_2 M + \log_2 \left(1 - \frac{1}{4\chi} \kappa M \bar{g}\right), \end{aligned} \quad (51)$$

where we have (a) due to  $\rho \rightarrow \infty$ ; we have (b) owing to  $\kappa \rightarrow 0$ ; (c) follows according to the Taylor approximation.

Upon substituting (50) and (51) into (31), we obtain

$$\begin{aligned} R_{S,\text{sec}}^J(\kappa) &= \log_2 M + \frac{1}{M} \left[ \log_2 \left(1 - \frac{1}{4\chi} \kappa M \bar{g}\right) \right. \\ &\quad \left. - \log_2 \left[1 + \ell \exp\left(-\frac{\rho\kappa \mathcal{P} \zeta^{\min}}{4}\right)\right] \right]. \end{aligned} \quad (52)$$

Since  $\log_2 \left(1 - \frac{1}{4\chi} \kappa M \bar{g}\right)$  is a log-concave function and  $1 + \ell \exp\left(-\frac{\rho\kappa \mathcal{P} \zeta^{\min}}{4}\right)$  is a log-convex function both w.r.t.  $\kappa$ ,  $R_{S,\text{sec}}^J(\kappa)$  is a concave function of  $\kappa$ . Therefore, the maximum of  $R_{S,\text{sec}}^J(\kappa)$  is achieved, when  $\kappa$  satisfies

$$\begin{aligned} \frac{\partial R_{S,\text{sec}}^J(\kappa)}{\partial \kappa} &= \frac{1}{M} \left[ \frac{-\frac{1}{4\chi} M \bar{g}}{1 - \frac{\kappa M \bar{g}}{4\chi}} - \frac{\ell \exp\left(-\frac{\rho\kappa \mathcal{P} \zeta^{\min}}{4}\right) \left(-\frac{\rho \mathcal{P} \zeta^{\min}}{4}\right)}{1 + \ell \exp\left(-\frac{\rho\kappa \mathcal{P} \zeta^{\min}}{4}\right)} \right] \\ &= 0. \end{aligned} \quad (53)$$

Explicitly, we have  $1 - \frac{\kappa M \bar{g}}{4\chi} \xrightarrow{\kappa \rightarrow 0} 1$ . With the aid of this approximation, the desired result in (34) can be obtained by the solution of (53).

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