

Finite-Cardinality Single-RF Differential Space-Time Modulation for Improving the Diversity-Throughput Tradeoff

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Abstract—The matrix-based differential encoding invoked by Differential Space-Time Modulation (DSTM) typically results in an infinite-cardinality of arbitrary signals, despite the fact that the Transmit Antennas (TAs) can only radiate a limited number of patterns. As a remedy, the recently developed Differential Spatial Modulation (DSM) is capable of avoiding this problem by conceiving a beneficial sparse signal matrix design, which also facilitates low-complexity single-RF signal transmission. Inspired by this development, the Differential Space-Time Block Code using Index Shift Keying (DSTBC-ISK) further introduces a beneficial diversity gain without compromising the DSM's appealingly low transceiver complexity. However, the DSTBC-ISK's performance advantage tends to diminish as the throughput increases, especially when an increased number of Receive Antennas (RAs) is used. By contrast, the classic Differential Group Code (DGC) that actively maximizes its diversity gain for different Multiple-Input Multiple-Output (MIMO) system setups is capable of achieving a superior performance, but its detection complexity grows exponentially with the throughput. Against this background, we propose the Differential Space-Time Shift Keying using Diagonal Algebraic Space-Time (DSTSK-DAST) scheme, which is the first DSTM that is capable of achieving the DGC's superior diversity gain at high throughputs without compromising the DSM's low transceiver complexity. As a further advance, we also conceive a new Differential Space-Time Shift Keying using Threaded Algebraic Space-Time (DSTSK-TAST) arrangement, which is capable of achieving an even further improved diversity gain at a substantially reduced signal detection complexity compared to the best DGCs. Furthermore, in order to strike a practical tradeoff, we develop a generic multi-element and multi-level-ring Amplitude Phase Shift Keying (APSK) design, and we also arrange for multiple reduced-size DSTM sub-blocks to be transmitted in a permuted manner, which exhibits an improved diversity-throughput tradeoff.

Index Terms—Differential space-time modulation, differential

spatial modulation, diagonal algebraic space-time, threaded algebraic space-time, space-time shift keying, group code, single-RF, finite-cardinality, single-stream ML detection, diversity gain, throughput-diversity tradeoff.

I. INTRODUCTION

Differential Space-Time Modulation (DSTM) constitutes a low-complexity design alternative to coherent Multiple-Input Multiple-Output (MIMO) schemes, where the excessive pilot overhead and the high-complexity channel estimation can be eliminated. Based on the classic Space-Time Block Codes (STBCs) [1]–[3], Differential STBCs (DSTBCs) were conceived in [4]–[6]. Moreover, based on the family of Linear Dispersion Codes (LDCs) [7]–[11] that are capable of achieving both the V-BLAST's high throughput and the STBC's full diversity, Differential LDCs (DLDCs) were conceived in [12], [13], where the Cayley transform was invoked in order to employ unitary signal matrices. The recently developed Differential Space-Time Shift Keying (DSTSK) [14]–[16] avoids the non-linear Cayley transform, where only a single unitary dispersion matrix is activated. As a benefit of the unitary and linear design, both DSTBC and DSTSK are capable of decoupling the received signals, where the single-stream ML detection complexity does not grow with the constellation size [6], [17], [18].

In contrast to both the DPSK [19], [20] and star QAM schemes [21]–[25] routinely used in Single-Input Single-Output (SISO) channels, the matrix multiplication invoked by the DSTM's differential encoding results in arbitrary infinite-cardinality of transmit signals for all the aforementioned DSTM schemes. However, in reality, the Transmit Antennas (TAs) can only radiate a limited number of signal transmission patterns [26]–[28]. As a remedy, inspired by the recently proposed Spatial Modulation (SM) [17], [18], [29], [30], the Differential SM (DSM) concept was investigated in [31]–[34], while their applications in the future wireless networks were also received in [35], [36]. More explicitly, the finite-cardinality of the classic PSK/QAM constellations was retained thanks to the DSM's beneficial sparse matrix design, which also facilitates low-complexity single-RF signal transmission. However, the DSM does not perform well, when the number of Receive Antennas (RAs) N is low. In order to introduce a beneficial transmit diversity gain, Field Extension based DSM (FE-DSM) was conceived in [37], and then it was refined in [38] in order to generate a finite-cardinality signal

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This work was supported in part by the EPSRC projects EP/N004558/1 and EP/PO34284/1, the European Research Council's Advanced Fellow Grant under the QuantCom project and the Royal Society's Wolfson Research Merit Award and the GRCF. The work of N. Ishikawa was supported in part by the Japan Society for the Promotion of Science KAKENHI under Grant 17H07036. The work of S. Sugiura was supported in part by the Japan Society for the Promotion of Science KAKENHI under Grants 26709028 and 16KK0120.

set, which became a special case of the DGC $G_{m,r}$ of [39] associated with $r = 1$. Moreover, the recently proposed DSTBC using Index Shift Keying (DSTBC-ISK) [40] transformed the DSTBC signals [4]–[6] to finite-cardinality sets by activating only a single modulated symbol in the STBC's signal structure. Although FE-DSM and DSTBC-ISK significantly outperform DSM at low throughputs, their diversity gains tend to diminish as the throughput increases, especially when a higher number of RAs is used.

As a remedy, the classic Differential Group Code (DGC) [39], [41]–[43] that forms a finite group may be invoked, where the diversity gain is consciously maximized by rotating the PSK signal elements for different Multiple-Input Multiple-Output (MIMO) system setups. However, owing to the associated non-linear phase rotations imposed on the data-carrying PSK signal, the DGC's ML detection complexity grows exponentially with the throughput. In summary, the family of DSTM schemes is characterized in terms of its transceiver complexity versus its performance in Table I. We note that for a generic DSTM, the notations M , N , T and Q represent the numbers of TAs, RAs, transmission time slots and dispersion matrices, respectively. The cardinality of the DSTM's data-carrying matrix set is denoted by I , hence the throughput is $R = \frac{\log_2 I}{T}$. *Against this background, our objective is to mitigate the problem of the diminishing diversity gain of FE-DSM and DSTBC-ISK, by invoking the DGC's diversity gain maximization, without compromising the DSM's low transceiver complexity.* More explicitly, the novel contributions of this paper are:

- 1) We propose a new DSTSK using Diagonal Algebraic Space-Time (DSTSK-DAST) coding. First of all, we portray the DAST of [7] as a LDC scheme using Q dispersion matrices. Then the proposed DSTSK-DAST transforms the LDC's 'matrix-multiplexing' form into the STSK's 'matrix-activation' form, where $\log_2 Q$ bits are assigned to activate a single one out of Q dispersion matrices. In order to retain a **finite-cardinality** signal set, we opt for confining the elements in the DSTSK-DAST's dispersion matrices to an equi-spaced L_{DM} -PSK constellation. As a result, on one hand, the DSTSK's **single-stream ML detection** is invoked without imposing any performance loss. On the other hand, the DGC's phase rotation method is applied to the L_{DM} -PSK elements in the dispersion matrices, so that the DSTSK-DAST's **diversity gain** is actively maximized for different MIMO setups.
- 2) Furthermore, we propose a new DSTSK using Threaded Algebraic Space-Time (DSTSK-TAST) scheme, which transforms the classic TAST of [10] also into a finite-cardinality single-RF and diversity-gain-maximizing DSTSK scheme. More explicitly, in order to efficiently exploit all the available degrees of freedom, the DSTSK-TAST divides the $(T \times T)$ -element space-time signal matrix into T non-overlapping threaded layers, where each threaded layer is a T -element vector. Specifically, the DSTSK-DAST's T -element diagonal signal vector constitutes one of the threaded layers in the DSTSK-

TAST signal structure. As a result, the DSTSK-TAST design may follow the same procedures as the DSTSK-DAST design, but the DSTSK-TAST scheme is capable of conveying $\log_2 T$ extra bits for activating a single one out of T threaded layers for signal transmission.

- 3) In order to improve the throughput, we extend the Single-Element (SE) Amplitude Phase Shift Keying (APSK) design of [40] to a generic Multi-Element (ME) design. Compared to the existing two/four-level-ring ME APSK solutions of [33], [34], we conceive a single-stream ML detector for the generic family of high-level star QAM signals, where the detection complexity does not grow with the constellation size.
- 4) Furthermore, in order to strike a practical tradeoff, we arrange for multiple reduced-size DSTM sub-blocks to be transmitted in a permuted manner, which improves Diversity-Rate (DR) tradeoff. This results in an improved throughput at the cost of a reduced diversity gain. Compared to the conventional DR design of [37], [38] using the FE signal matrix, we opt for invoking the layer-switching matrix of the proposed DSTSK-TAST scheme, so that the resultant bespoke DR design may be applied to the whole family of the diversity-oriented DSTM schemes.

The rest of this paper is organized as follows. The design objectives and preliminaries are presented Sec. II. The new DSTSK-DAST and DSTSK-TAST are proposed in Sec. III and Sec. IV, respectively. The ME APSK design and the DR tradeoff design are detailed in Sec. V. The simulation results are presented in Sec. VI, and our conclusions are offered in Sec. VII

The following notations are used throughout the paper. $\Re(\cdot)$ and $\Im(\cdot)$ represent the real and imaginary parts of complex numbers, respectively. $(\cdot)^*$, $(\cdot)^T$ and $(\cdot)^H$ denote the conjugate of a complex number, the transpose of a matrix and the Hermitian transpose of a complex matrix, respectively. \otimes represents the Kronecker product. $\text{tr}(\cdot)$, $\text{rank}(\cdot)$ and $\det(\cdot)$ take trace, rank and determinant of a matrix, respectively.

II. DESIGN OBJECTIVES AND PRELIMINARIES

In this section, our design objectives are highlighted in the context of the transmitted signal model of Sec. II-A, the received signal model of Sec. II-B as well as the DGC's diversity gain maximization of Sec. II-C. The high-throughput DSM is also briefly reviewed in Sec. II-D.

A. Transmitted Signal Model

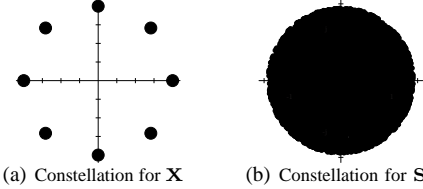
The DSTM's $(T \times M)$ -element signal matrix \mathbf{S}_n is obtained by differential encoding:

$$\mathbf{S}_n = \mathbf{X}_{n-1} \mathbf{S}_{n-1}, \quad (1)$$

where we have the constraint of $\text{tr}(\mathbf{S}_n^H \mathbf{S}_n) = T$ and $(M \leq T)$, while the $(T \times T)$ -element data-signal matrix \mathbf{X}_{n-1} carries source information. The recursion of (1) commences from $\mathbf{S}_1 = \sqrt{\frac{T}{M}} [\mathbf{I}_M, \mathbf{0}]^T$ [43], where the all-zero matrix $\mathbf{0}$ has $[M \times (T - M)]$ elements.

TABLE I: Summary of the DSTM schemes.

	Transmitter complexity		Receiver complexity		Performance	
	Finite-cardinality?	Single-RF transmission?	Single-stream ML detection?	Detection complexity order	Transmit diversity?	Diversity gain maximization?
DSTBC [4]–[6]	×	×	✓	$\mathcal{O}(Q)$	✓	×
DLDC [12], [13]	×	×	×	$\mathcal{O}(I = 2^{RT})$	✓	✓
DSTSK [14]–[16]	×	×	✓	$\mathcal{O}(Q)$	✓	✓
DGC [39], [41]–[43]	✓	✓	×	$\mathcal{O}(I = 2^{RT})$	✓	✓
DSM [31]–[34]	✓	✓	✓	$\mathcal{O}(2^{\lceil \log_2 T \rceil})$	×	×
FE-DSM [37], [38]	✓	✓	✓	$\mathcal{O}(T)$	✓	×
DSTBC-ISK [40]	✓	✓	✓	$\mathcal{O}(T)$	✓	×
DSTSK-DAST	✓	✓	✓	$\mathcal{O}(Q)$	✓	✓
DSTSK-TAST	✓	✓	✓	$\mathcal{O}(QT)$	✓	✓

Fig. 1: Constellations of DSTBC [6] signals in \mathbf{X} and \mathbf{S} of (1), where $M = 2$ TAs and 8PSK signals are used.

The data-signal matrix \mathbf{X}_{n-1} of (1) may invoke any unitary MIMO signal structures including STBC, STSK and Cayley transformed LDC, which form the DSTM schemes of DSTBC [4]–[6], DSTSK [14]–[16] and DLDC [12], [13], respectively. However, when coherent MIMO signals are directly used in DSTM, the matrix multiplication of (1) typically results in arbitrary transmitted signals having a cardinality tending to infinity. This problem is exemplified for the case of DSTBC [6] in Fig. 1. Although the differential design aims for a low transceiver complexity, compared to the coherent STBC, the DSTBC's infinite-cardinality problem actually imposes significantly extra constraints on the speed, precision and power consumption of both the digital circuitry and on the Digital-to-Analog Converter (DAC). Moreover, when the DSTBC signals of Fig. 1(b) are quantized to a finite-cardinality set, the associated quantization error results in a severe performance loss, as demonstrated in [40]. As a remedy, the finite-cardinality design of DSM [31]–[34], FE-DSM [37], [38], DSTBC-ISK [40] and DGC [39], [41]–[43] confines the signal matrix \mathbf{X}_{n-1} of (1) to have only a single non-zero element in each row and column, where the non-zero signal is drawn from the equispaced PSK and star QAM constellations. As a further benefit of using the sparse matrix, only a single RF chain is activated for signal transmission.

We note that a single-carrier system is assumed in this paper, owing to the fact that the single-RF feature cannot be retained by Orthogonal Frequency-Division Multiplexing (OFDM) [44]. Nonetheless, considering that the products of the communication industry are typically standardized, our main objective of preserving the finite-cardinality feature also offers a practical interface to OFDM applications, where the DSTM signals in \mathbf{S}_n may be fed into the Inverse Discrete Fourier Transform (IDFT) block as the frequency-domain input signals. More explicitly, the frequency-domain signals entered into the IDFT block at the OFDM transmitter are generally drawn from the standardized finite-cardinality PSK/QAM constellations. In particular, a variety of practical techniques - including active constellation extension, selected mapping, tone injection and tone reservation [45]–[48] - explicitly aim

for the OFDM time-domain signal's Peak-to-Average Power Ratio (PAPR) reduction based on the frequency-domain input PSK/QAM constellations. Therefore, the finite-cardinality feature that retains the standardized PSK/QAM signal constellation is of practical importance to the applications both in single-carrier systems and in multi-carrier OFDM systems.

B. Received Signal Model

In order to present a fair comparison against the existing finite-cardinality DSTM schemes of DSM [31]–[34], FE-DSM [37], [38], DSTBC-ISK [40] and DGC [39], [41]–[43], the same narrowband single-carrier model is invoked, where the received signals are given by:

$$\mathbf{Y}_n = \mathbf{S}_n \mathbf{H}_n + \mathbf{V}_n. \quad (2)$$

\mathbf{Y}_n , \mathbf{H}_n and \mathbf{V}_n are the $(T \times N)$ -element received signal matrix, the $(M \times N)$ -element Rayleigh fading channel matrix as well as the $(T \times N)$ -element AWGN matrix, respectively. In quasi-static fading channels, we have $\mathbf{H}_n = \mathbf{H}_{n-1}$, hence (2) becomes $\mathbf{Y}_n = \mathbf{X}_{n-1}(\mathbf{Y}_{n-1} - \mathbf{V}_{n-1}) + \mathbf{V}_n$. This leads to the following differential detection:

$$\hat{\mathbf{X}}_{n-1} = \arg \min_{\mathbf{X}^i} \|\mathbf{Y}_n - \mathbf{X}^i \mathbf{Y}_{n-1}\|^2, \quad (3)$$

which results in the following Pairwise Error Probability (PEP) [49], [50] for the case of $M = T$:

$$p(\mathbf{X}^i \rightarrow \mathbf{X}^{i'}) \leq \left\{ 1 + \left[\frac{1}{4N_0(N_0+2)} \right]^{\text{rank}(\Delta)} \frac{\det(\Delta)}{\det(\Delta) + \frac{1}{4N_0(N_0+2)} \text{tr}(\Delta)} \right\}^{-N}, \quad (4)$$

where the determinant $\det(\Delta) = \prod_{t=1}^T \lambda_t$ and the trace term $\text{tr}(\Delta) = \sum_{t=1}^T \lambda_t$ dominates the PEP in the high- and low-SNR regions, respectively, while $\{\lambda_t\}_{t=1}^T$ are eigenvalues of $\Delta = (\mathbf{X}^i - \mathbf{X}^{i'})^H (\mathbf{X}^i - \mathbf{X}^{i'})$. Accordingly, the trend of (4) may be characterized by the *diversity product* and by the *diversity sum* as [50]:

$$\Lambda_p = \frac{1}{2} \min_{i \neq i'} \det(\Delta)^{\frac{1}{2T}}. \quad (5a)$$

$$\Lambda_s = \frac{1}{2\sqrt{T}} \min_{i \neq i'} \text{tr}(\Delta)^{\frac{1}{2}}. \quad (5b)$$

The tradeoff between Λ_p and Λ_s explicitly characterizes the ubiquitous diversity-throughput tradeoff in MIMO design, which is exemplified in Fig. 2. Detailed introductions to DGC and DSM will be offered in Sec. II-C and Sec. II-D, respectively. It is demonstrated by Fig. 2 that on one hand, the DGC-cyclic associated with the maximized $\Lambda_p = 0.3827$ achieves substantial diversity gains for $N = \{1, 2\}$, but it does not perform well for large $N = \{8, 16\}$. On the other

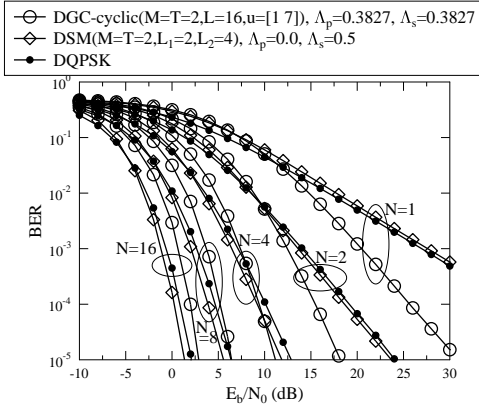


Fig. 2: BER performance comparison between DGC and DSM associated with $M = T = 2$ and $R = 2.0$.

TABLE II: Example of DGC-cyclic ($M = T = 2$, $L = 8$) at $R = 1.5$.

$\mathbf{X}^0 = \begin{bmatrix} 1 & 0 \\ 0 & 1 \end{bmatrix}$	$\mathbf{X}^1 = \begin{bmatrix} w_8 & 0 \\ 0 & w_8^3 \end{bmatrix}$	$\mathbf{X}^2 = \begin{bmatrix} w_8^2 & 0 \\ 0 & w_8^6 \end{bmatrix}$	$\mathbf{X}^3 = \begin{bmatrix} w_8^3 & 0 \\ 0 & w_8^1 \end{bmatrix}$
$\mathbf{X}^4 = \begin{bmatrix} w_8^4 & 0 \\ 0 & w_8^4 \end{bmatrix}$	$\mathbf{X}^5 = \begin{bmatrix} w_8^5 & 0 \\ 0 & w_8^7 \end{bmatrix}$	$\mathbf{X}^6 = \begin{bmatrix} w_8^6 & 0 \\ 0 & w_8^2 \end{bmatrix}$	$\mathbf{X}^7 = \begin{bmatrix} w_8^7 & 0 \\ 0 & w_8^5 \end{bmatrix}$

hand, the DSM associated with $\Lambda_p = 0$ performs even worse than its SISO DPSK counterpart for $N = 1$ in Fig. 2, but the DSM outperforms both DGC-cyclic and DPSK for large $N = \{8, 16\}$. This is because the DSM's throughput of $R = \frac{\lfloor \log_2 T \rfloor + T \log_2 L}{T}$ is higher than the DPSK's $R = \log_2 L$ and the DGC-cyclic's $R = \frac{\log_2 L}{T}$. As a result, upon comparing the performance at the same throughput, the DSM may opt for a lower-level LPSK associated with a higher Euclidean distance, which leads to a higher Λ_s of (5b) and hence a better performance for large N .

Essentially, all operational and future wireless communication systems are adaptive. For instance, the low-complexity transceiver of SM and DSM schemes is favourable for the evolving Internet-of-Things (IoT) in next generation networks. For the billions of connected smart devices, it is especially beneficial to adaptively switch between high-diversity and high-throughput schemes according to the specific service demands and link qualities. Therefore, in this work, we aim for devising new DSTM scheme that is capable of even outperforming the classic DGC, which mitigates the problem of the diminishing diversity gain of the recently developed FE-DSM and DSTBC-ISK, without compromising the DSM's appealingly low transceiver complexity.

C. Differential Group Code (DGC) Diversity Gain Maximization

For cyclic DGCs [39], [41]–[43], the signal-matrix \mathbf{X} of (1) is constructed by $\mathbf{X}^l = \mathbf{G}_c^l$, where $\mathbf{G}_c = \text{diag}([w_L^{u_1}, w_L^{u_2}, \dots, w_L^{u_T}])$ and we have $w_L = \exp(j \frac{2\pi}{L})$. The integer phase-rotation parameters $\mathbf{u} = [u_1, u_2, \dots, u_T]$ are chosen for the sake of maximizing the diversity product of $\Lambda_p = \min_{l \neq 0} [\prod_{t=1}^T |\sin(\frac{\pi u_t l}{L})|]^{\frac{1}{T}}$. Moreover, the dicyclic DGCs [39], [41]–[43] convey an extra bit determining whether the signal matrix is diagonal or anti-diagonal. It is shown in [39], [43] that DGC-cyclic generally outperform DGC-dicyclic.

Given a MIMO system setup, the DGC-cyclic strives for actively maximizing its diversity product of $\Lambda_p =$

TABLE III: Example of DSM associated with $M = T = 2$

Permutation Index	Activation Sequence	Data-Carrying Matrix
$\bar{m} = 1$	$\mathbf{a}_1 = [1, 2]$	$\mathbf{X}^{\{\bar{m}=1, l_1, l_2\}} = \begin{bmatrix} x^{l_1} & 0 \\ 0 & x^{l_2} \end{bmatrix}$
$\bar{m} = 2$	$\mathbf{a}_2 = [2, 1]$	$\mathbf{X}^{\{\bar{m}=2, l_1, l_2\}} = \begin{bmatrix} 0 & x^{l_1} \\ x^{l_2} & 0 \end{bmatrix}$

$\min_{l \neq 0} [\prod_{t=1}^T |\sin(\frac{\pi u_t l}{L})|]^{\frac{1}{T}}$ by searching for the best combination of $\{1 \leq u_t \leq L-1\}_{t=1}^T$. This exhaustive search is simplified in [42] as: (a) $\{u_t\}_{t=1}^T$ are relatively prime to L . Otherwise, we have $[u_t l \bmod L = 0]$ that results in $\Lambda_p = 0$; (b) The ordering of $\{u_t\}_{t=1}^T$ does not change Λ_p , hence we set $u_1 \leq u_2 \leq \dots \leq u_T$; (c) u_t and $(L_{DM} - u_t)$ leads to the same Λ_p , which results in a reduced range of $\{1 \leq u_t \leq L/2 - 1\}_{t=1}^T$; (d) $u_1 = 1$ is set, because of rule (c) and also because multiplying $\{u_t\}_{t=2}^T$ by a common integer does not change Λ_p .

Example 1: For a DGC-cyclic scheme associated with $M = T = 2$ and $L = 8$, the legitimate parameters according to the above search rules are given by $\mathbf{u} = [1, 1]$ and $\mathbf{u} = [1, 3]$, which lead to $\Lambda_p = 0.3827$ and $\Lambda_p = 0.5946$, respectively. Therefore, the latter is chosen, and the resultant signal matrices are given by $\{\mathbf{X}^l = \text{diag}([w_8^l, w_8^{3l}])\}_{l=0}^7$, which are detailed in Table II.

Example 2: For DGC-cyclic associated with $M = T = 4$ and $L = 32$, there are a total of 120 candidates for \mathbf{u} ranging from $\mathbf{u} = [1, 1, 1, 1]$ to $\mathbf{u} = [1, 15, 15, 15]$ according to the above search rules. The maximized $\Lambda_p = 0.3827$ is given by $\mathbf{u} = [1, 7, 9, 15]$, which leads to the DGC signal matrices of $\{\mathbf{X}^l = \text{diag}([w_{32}^l, w_{32}^{7l}, w_{32}^{9l}, w_{32}^{15l}])\}_{l=0}^{31}$.

Owing to the non-linear phase rotations on the data-carrying LPSK signal, the DGC detection invoking (3) has a complexity order of $\mathcal{O}(I)$, which grows exponentially with the throughput.

D. Differential Spatial Modulation (DSM)

The DSM [31]–[34] modulates a total of $(M=TQ)$ $\{L_t\}_{t=1}^T$ -PSK symbols $\{x^{lt} = \exp(j \frac{2\pi}{L_t} \tilde{l}_t)\}_{t=1}^T$ from a total of $\sum_{t=1}^T \log_2 L_t$ bits. Moreover, a total of $\lfloor \log_2 M! \rfloor = \log_2 \bar{M}$ bits are assigned for determining the activation sequence $\mathbf{a}_{\bar{m}} = [a_{\bar{m},1}, a_{\bar{m},2}, \dots, a_{\bar{m},T}]$, which obeys $1 \leq \{a_{\bar{m},t}\}_{t=1}^T \leq M$ and $a_{\bar{m},1} \neq a_{\bar{m},2} \neq \dots \neq a_{\bar{m},T}$. The permutation index \bar{m} ranges $1 \leq \bar{m} \leq \bar{M}$. This presentation indicates that the $a_{\bar{m},t}$ -th element on the t -th row of the $(T \times T)$ -element data-carrying matrix \mathbf{X}_{n-1} in (1) is activated to transmit x^{lt} :

$$\mathbf{X}^i(r, c) = \begin{cases} x^{lt}, & \text{if } r = t \text{ and } c = a_{\bar{m},t} \\ 0, & \text{all the other elements} \end{cases}, \quad (6)$$

where $\mathbf{X}^i(r, c)$ denotes the element on the r -th row and c -th column in \mathbf{X}^i . The example of DSM associated with $(M = 2)$ is characterized by Table III.

Therefore, a pair of DSM codewords associated with $\bar{m} = \bar{m}'$, $l_1 = l'_1$ and $\{l_t \neq l'_t\}_{t=2}^T$ leads to $\Lambda_p = 0$ in (5a). Nonetheless, in the family of the existing finite-cardinality single-RF DSTM schemes, DSM offers the highest throughput of $R = \frac{\lfloor \log_2 T \rfloor + T \log_2 L}{T}$, which is beneficial for using large N , as discussed in Sec. II-B. Moreover, we note that DSM may invoke the single-stream ML detector designed in Sec. IV of [40], which has a reduced complexity order of $\mathcal{O}(\bar{M})$.

III. DIFFERENTIAL SPACE-TIME SHIFT KEYING USING DIAGONAL ALGEBRAIC SPACE-TIME (DSTSK-DAST)

In this section, we propose the new DSTSK-DAST scheme, which simultaneously achieves our design objectives of (1) forming a finite-cardinality set of transmitted signals, (2) retaining a single-stream ML detection complexity and (3) maximizing the diversity gain.

A. The Relationship Between DAST and LDC

Before the development of space-time diversity for MIMO systems, the temporal diversity was imposed onto a sequence of Q modulated symbols $\mathbf{x} = [x_1, \dots, x_Q]^T$ in SISO systems by a series of Q dispersion elements $\{\{\bar{a}_{t,q}\}_{t=1}^Q\}_{q=1}^Q$ as $\bar{x}_t = \sum_{q=1}^Q \bar{a}_{t,q} x_q$. As a result, the receiver is capable of recovering the dispersed sequence of T correlated symbols $\bar{\mathbf{x}} = [\bar{x}_1, \dots, \bar{x}_T]^T$ unless the whole sequence experiences a deep fade. This SISO diversity design was termed as lattice-based multidimensional constellation [51]–[53], where the best rotations $\{\{\bar{a}_{t,q}\}_{t=1}^Q\}_{q=1}^Q$ are obtained according to algebraic number theory. More explicitly, according to the construction based on PSK signalling in [51], the dispersed sequence may be represented by:

$$\bar{\mathbf{x}} = \mathbf{G}_A \mathbf{x}, \quad (7)$$

where $\mathbf{G}_A = \text{VDM}_{T \times Q}(g_1, \dots, g_T)$ is a $(T \times Q)$ -element Vandermonde matrix:

$$\mathbf{G}_A = \begin{bmatrix} \bar{a}_{1,1} & \bar{a}_{1,2} & \bar{a}_{1,3} & \dots & \bar{a}_{1,Q} \\ \bar{a}_{2,1} & \bar{a}_{2,2} & \bar{a}_{2,3} & \dots & \bar{a}_{2,Q} \\ \vdots & \vdots & \vdots & \ddots & \vdots \\ \bar{a}_{T,1} & \bar{a}_{T,2} & \bar{a}_{T,3} & \dots & \bar{a}_{T,Q} \end{bmatrix} = \begin{bmatrix} 1 & g_1 & g_1^2 & \dots & g_1^{Q-1} \\ 1 & g_2 & g_2^2 & \dots & g_2^{Q-1} \\ \vdots & \vdots & \vdots & \ddots & \vdots \\ 1 & g_T & g_T^2 & \dots & g_T^{Q-1} \end{bmatrix}. \quad (8)$$

Therefore, the diversity is imposed by the spreading operation of $\bar{x}_t = \sum_{q=1}^Q \bar{a}_{t,q} x_q = \sum_{q=1}^Q g_t^{q-1} x_q$, where $\{g_t\}_{t=1}^T$ are taken from a $4T$ -PSK constellation according to [51].

In order to achieve the space-time diversity in MIMO systems, the DAST scheme in [7] proposed to transmit the dispersed sequence of $\bar{\mathbf{x}}$ on diagonal of the MIMO's space-time signal matrix. More explicitly, DAST using PSK may be constructed by:

$$\mathbf{X} = \text{diag}(\mathbf{G}_A \mathbf{x}) = \text{diag}\left(\left[\sum_{q=1}^Q g_1^{q-1} x_q, \sum_{q=1}^Q g_2^{q-1} x_q, \dots, \sum_{q=1}^Q g_T^{q-1} x_q\right]\right). \quad (9)$$

As a result, the DAST signal matrix of (9) may be further expressed in the LDC form [9] as:

$$\mathbf{X} = \sum_{q=1}^Q x_q \mathbf{A}_q, \quad (10)$$

where we have the equivalent LDC dispersion matrices of $\{\mathbf{A}_q = \text{diag}([g_1^{q-1}, \dots, g_T^{q-1}])\}_{q=1}^Q$, while $Q = T$ is assumed for the DAST scheme in [7].

Although both $\{g_t\}_{t=1}^T$ and $\{x_q\}_{q=1}^Q$ are PSK signals, the dispersed symbols $\{\bar{x}_t = \sum_{q=1}^Q g_t^{q-1} x_q\}_{t=1}^T$ do not remain

TABLE IV: Examples of DGC-cyclic and DSTSK-DAST associated with $M = T = 4$.

DGC-cyclic	DSTSK-DAST
$\mathbf{X}^l = \mathbf{G}_c^l = \begin{bmatrix} w_L^{u_1} & 0 & 0 & 0 \\ 0 & w_L^{u_2} & 0 & 0 \\ 0 & 0 & w_L^{u_3} & 0 \\ 0 & 0 & 0 & w_L^{u_4} \end{bmatrix}^l$	$\mathbf{X}^i = x^i \mathbf{A}_q = w_L^i \begin{bmatrix} w_{L_{DM}}^{u_1} & 0 & 0 & 0 \\ 0 & w_{L_{DM}}^{u_2} & 0 & 0 \\ 0 & 0 & w_{L_{DM}}^{u_3} & 0 \\ 0 & 0 & 0 & w_{L_{DM}}^{u_4} \end{bmatrix}^q$

limited to the PSK constellation owing to the additions of " $\sum_{q=1}^Q$ ". Consequently, when the signal matrix of (10) is invoked in the differential encoding of (1), the infinite-cardinality problem discussed in Sec. II-A arises. Moreover, we note that the DAST's signal matrix of (10) is not unitary, which prevents the employment of a single-stream ML detector.

B. Finite-Cardinality Design

In order to conceive a finite-cardinality DSTSK-DAST scheme, we revise the LDC's matrix multiplexing form of (10) to the STSK's matrix activation form [14] as:

$$\mathbf{X} = x^l \mathbf{A}_q. \quad (11)$$

More explicitly, a total of $\log_2 L$ bits are assigned to modulate a L -PSK symbol $\{x^l = w_L^l\}_{l=0}^{L-1}$, and a total of $\log_2 Q$ bits are assigned to activate a single out of Q dispersion matrix \mathbf{A}_q . The dispersion matrices are given by $\{\mathbf{A}_q = \text{diag}([w_{L_{DM}}^{(q-1)u_1}, w_{L_{DM}}^{(q-1)u_2}, \dots, w_{L_{DM}}^{(q-1)u_T}])\}_{q=1}^Q$, where the dispersion elements are taken from an L_{DM} -PSK constellation. The sequence of integers $\mathbf{u} = [u_1, u_2, \dots, u_T]$ are invoked for the sake of diversity gain maximization in the same way as the DGC, which will be further detailed in Sec. III-D.

The revised signal matrix of (11) is equivalent to the DAST of (9), where the $(T \times Q)$ -element Vandermonde matrix \mathbf{G}_A is now given by:

$$\mathbf{G}_A = \text{VDM}_{T \times Q}(w_{L_{DM}}^{u_1}, \dots, w_{L_{DM}}^{u_T}) = \begin{bmatrix} 1 & w_{L_{DM}}^{u_1} & w_{L_{DM}}^{2u_1} & \dots & w_{L_{DM}}^{(Q-1)u_1} \\ 1 & w_{L_{DM}}^{u_2} & w_{L_{DM}}^{2u_2} & \dots & w_{L_{DM}}^{(Q-1)u_2} \\ \vdots & \vdots & \vdots & \ddots & \vdots \\ 1 & w_{L_{DM}}^{u_T} & w_{L_{DM}}^{2u_T} & \dots & w_{L_{DM}}^{(Q-1)u_T} \end{bmatrix}, \quad (12)$$

and the Q -element signal vector \mathbf{x} is now given by:

$$\mathbf{x} = [\underbrace{0 \dots 0}_{q-1}, x^l, \underbrace{0 \dots 0}_{Q-q}]. \quad (13)$$

Therefore, in the absence of signal additions, the DSTSK-DAST signal matrix of (11) complies with the finite-cardinality design of Sec. II-A. We note that (11) requires $L_{DM} \geq LQ$ for creating distinct codewords $\{\mathbf{X}^i\}_{i=0}^I$. We opt for choosing $L_{DM} = LQ$ for the sake of simplicity. Moreover, for the proposed DSTSK-DAST, Q and T do not have to be equal, hence the throughput is now given by $R = \frac{\log_2 L + \log_2 Q}{T}$.

C. Single-Stream ML Detection

The relationship between DGC-cyclic and DSTSK-DAST is further exemplified in Table IV, which demonstrates that instead of directly rotating the data-carrying L PSK signal,

the DSTSK-DAST applies the phase-rotation parameters of $\mathbf{u} = [u_1, u_2, \dots, u_T]$ to the L_{DM} PSK signals in the dispersion matrix. As a result, the detection of (3) may be decoupled for the DSTSK-DAST as:

$$\hat{\mathbf{X}}_{n-1} = \arg \max_{\forall \mathbf{X}_{n-1}} \Re [\text{tr}(\mathbf{X}_{n-1}^H \mathbf{Z})] = \arg \max_{\forall q, l} \Re [(x^l)^* \tilde{z}_q]. \quad (14)$$

We have $\mathbf{Z} = \mathbf{Y}_n \mathbf{Y}_{n-1}^H$ and $\tilde{z}_q = \text{tr}(\bar{\mathbf{A}}_q^H \mathbf{Z}) = \sum_{t=1}^T w_{L_{DM}}^{-(q-1)u_t} z_{t,t}$, where $z_{t,t}$ denotes the element on the t -th row and t -th column in \mathbf{Z} . As a result, the single-stream L -PSK demodulator $\mathbb{M}_{\text{PSK}}^{-1}(\cdot)$ may be invoked a total of Q times for all indices q as:

$$\hat{l}_q = \mathbb{M}_{\text{PSK}}^{-1}(\tilde{z}_q), \quad (15)$$

where the phase of \tilde{z}_q may be directly rounded¹ as $\hat{l}_q = \lfloor \frac{L}{2\pi} \angle \tilde{z}_q \rfloor$. Following this, the optimum activation index may be detected by invoking (14) for a reduced number of Q times as:

$$\hat{q} = \arg \max_{\forall q} \Re [(x^{\hat{l}_q})^* \tilde{z}_q]. \quad (16)$$

Finally, the optimum L -PSK index is $\hat{l} = \hat{l}_{\hat{q}}$, which is the one obtained in (15) associated with \hat{q} of (16). In summary, the DSTSK-DAST's single-stream-based ML detection complexity order is given by $\mathcal{O}(Q)$, which does not grow with the L -PSK constellation size.

D. Diversity Gain Maximization

In order to simplify the parameter search, let us firstly introduce the following proposition for the simplified evaluation of the DSTSK-DAST's diversity product:

Proposition 1: The DSTSK-DAST's diversity product of (5a) is given by

$$\Lambda_p = \min_{\forall (l \neq 0 || q \neq 1)} \prod_{t=1}^T \left| \sin \left[\frac{\pi(qu_t - u_t + lL_{DM}/L)}{L_{DM}} \right] \right|^{1/T}. \quad (17)$$

Proof. Owing to the fact that DSTSK-DAST's signal matrices of (11) form a finite group under multiplication [39], [41]–[43], the determinant term in (5a) may be further expressed as:

$$\begin{aligned} \det(\Delta) &= \det[(\mathbf{I}_T - x^l \mathbf{A}_q)^H (\mathbf{I}_T - x^l \mathbf{A}_q)] \\ &= \prod_{t=1}^T \left| 1 - w_L^l w_{L_{DM}}^{(q-1)u_t} \right|^2 \\ &= \prod_{t=1}^T \left| 1 - \exp \left\{ \frac{2\pi j[(q-1)u_t + lL_{DM}/L]}{L_{DM}} \right\} \right|^2, \end{aligned} \quad (18)$$

which results in (17) according to the relationship of $|1 - \exp(j\theta)|^2 = 4 \sin^2 \frac{\theta}{2}$. \square

As a result, with given a MIMO system setup, the DSTSK-DAST may also actively maximize Λ_p of (17) by searching over all the possible combinations of $\{1 \leq u_t \leq L_{DM}-1\}_{t=1}^T$. The search space may be further reduced according to:

¹We note that the detected index \hat{l} recovers the L PSK index l that can be translated back to source bits as $\text{dec2bin}(l) = (b_1 \dots b_{\log_2 L})$. Moreover, l is the Gray coded version of the natural index \hat{l} as defined in the L PSK modulation of $x^l = \exp(\frac{2\pi j}{L} \hat{l})$.

TABLE V: Example of DSTSK-DAST($M = T = 2$) using $Q = 2$, $L = 4$, $L_{DM} = 8$ and $\mathbf{u} = [1, 5]$ at $R = 1.5$.

Input bits	PSK symbol	Dispersion matrix	Data-carrying matrix
000	$x^0 = 1$	$\mathbf{A}_1 = \begin{bmatrix} 1 & 0 \\ 0 & 1 \end{bmatrix}$	$\mathbf{X} = x^0 \mathbf{A}_1 = \begin{bmatrix} 1 & 0 \\ 0 & 1 \end{bmatrix}$
001	$x^0 = 1$	$\mathbf{A}_2 = \begin{bmatrix} w_8 & 0 \\ 0 & w_8^5 \end{bmatrix}$	$\mathbf{X} = x^0 \mathbf{A}_2 = \begin{bmatrix} w_8 & 0 \\ 0 & w_8^5 \end{bmatrix}$
010	$x^1 = j$	$\mathbf{A}_1 = \begin{bmatrix} 1 & 0 \\ 0 & 1 \end{bmatrix}$	$\mathbf{X} = x^1 \mathbf{A}_1 = \begin{bmatrix} w_8^2 & 0 \\ 0 & w_8^2 \end{bmatrix}$
011	$x^1 = j$	$\mathbf{A}_2 = \begin{bmatrix} w_8 & 0 \\ 0 & w_8^5 \end{bmatrix}$	$\mathbf{X} = x^1 \mathbf{A}_2 = \begin{bmatrix} w_8^3 & 0 \\ 0 & w_8^7 \end{bmatrix}$
100	$x^2 = -j$	$\mathbf{A}_1 = \begin{bmatrix} 1 & 0 \\ 0 & 1 \end{bmatrix}$	$\mathbf{X} = x^2 \mathbf{A}_1 = \begin{bmatrix} w_8^6 & 0 \\ 0 & w_8^6 \end{bmatrix}$
101	$x^2 = -j$	$\mathbf{A}_2 = \begin{bmatrix} w_8 & 0 \\ 0 & w_8^5 \end{bmatrix}$	$\mathbf{X} = x^2 \mathbf{A}_2 = \begin{bmatrix} w_8^7 & 0 \\ 0 & w_8^3 \end{bmatrix}$
110	$x^3 = -1$	$\mathbf{A}_1 = \begin{bmatrix} 1 & 0 \\ 0 & 1 \end{bmatrix}$	$\mathbf{X} = x^3 \mathbf{A}_1 = \begin{bmatrix} w_8^4 & 0 \\ 0 & w_8^4 \end{bmatrix}$
111	$x^3 = -1$	$\mathbf{A}_2 = \begin{bmatrix} w_8 & 0 \\ 0 & w_8^5 \end{bmatrix}$	$\mathbf{X} = x^3 \mathbf{A}_2 = \begin{bmatrix} w_8^5 & 0 \\ 0 & w_8 \end{bmatrix}$

- (a) $\{u_t\}_{t=1}^T$ are relatively prime to L_{DM} . Otherwise, there exists $[u_t(q-1) \bmod L_{DM} = 0]$ for some integer q , which results in $\Lambda_p = 0$ in (17).
- (b) The ordering of $\{u_t\}_{t=1}^T$ does not change Λ_p , hence we set $u_1 \leq u_2 \leq \dots \leq u_T$.

We note that unlike DGC of Sec. II-C, u_t and $(L_{DM} - u_t)$ are no longer equivalent in (17). In summary, the resultant search space is bounded by a total of $[(\frac{L_{DM}}{2})^T / 2 + (\frac{L_{DM}}{2})^{T/2} / 2]$ combinations for $\{u_t\}_{t=1}^T$.

The brute-force exhaustive search constitutes a better choice than random search [8], [9], [14] or the classic gradient-ascent based search [12], [54], [55], because the cost function of (17) is not concave. Nonetheless, when the search space becomes excessive as Q and L increase, we may still resort to random search, since the optimization of (17) is no longer sensitive to its values at a high throughput. As a further benefit, instead of storing the dispersion matrices at the conventional DSTSK transmitters [14]–[16], the DSTSK-DAST transmitter only has to store the T -element integer phase-rotation parameters $\mathbf{u} = [u_1, \dots, u_T]$.

Example 3: Let us consider the DSTSK-DAST using $M = T = 2$ at the throughput of $R = \frac{\log_2 LQ}{T} = 1.5$, where we have $L_{DM} = LQ = 8$. Considering that the DSTSK-DAST detection complexity order in Sec. III-C is given by $\mathcal{O}(Q)$, we commence from the parameter combination of $Q = 1$ and $L = 8$. The exhaustive search gives $\mathbf{u} = [1, 1]$ associated with the diversity product of $\Lambda_p = 0.3827$, which is still smaller than the $\Lambda_p = 0.5946$ value of its DGC counterpart in Example 1. Therefore, we further increase Q and use the updated combination of $Q = 2$ and $L = 4$. The exhaustive search gives $\mathbf{u} = [1, 5]$ achieving $\Lambda_p = 0.5946$. The corresponding signal matrices are summarized in Table V. The resultant DSTSK-DAST detection complexity order is $\mathcal{O}(Q = 2)$, which is substantially lower than the full-search based DGC detection complexity order of $\mathcal{O}(I = 8)$ in Example 1, despite the fact that both DSTSK-DAST and DGC achieve the same maximized diversity product of $\Lambda_p = 0.5946$.

Example 4: For a DSTSK-DAST using $M = T = 4$ at $R = 1.25$, following the same procedures as Example 3, we arrive at the parameters of $Q = 4$, $L = 8$, $L_{DM} = 32$ and $\mathbf{u} = [7, 15, 23, 31]$, which achieves the same $\Lambda_p = 0.3827$ as the DGC in Example 2. The resultant DSTSK-DAST signal matrices are given by $\{\{\mathbf{X}^i =$

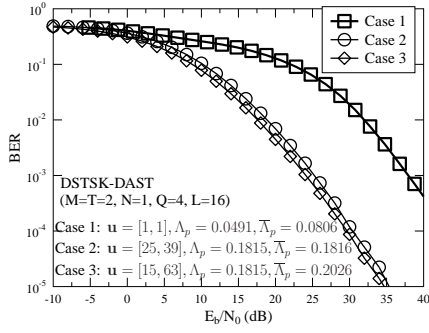


Fig. 3: Performance of DSTSK-DAST ($M=T=2$, $N=1$, $Q=4$, $L=16$) associated with different \mathbf{u} .

$w_8^l \mathbf{A}_q \}_{l=0}^7 \}_{q=1}^4$, where the dispersion matrices are given by $\{\mathbf{A}_q = \text{diag}([w_{32}^{7(q-1)}, w_{32}^{15(q-1)}, w_{32}^{23(q-1)}, w_{32}^{31(q-1)}])\}_{q=1}^Q$. Similarly, the resultant DSTSK-DAST detection complexity order of $\mathcal{O}(Q=4)$ is substantially lower than the DGC complexity order of $\mathcal{O}(I=32)$ in Example 2.

E. Further Considerations Concerning Average BER

It was demonstrated in [54] that the diversity product Λ_p of (5a) that characterizes the worst PEP $p(\mathbf{X}^i \rightarrow \mathbf{X}^{i'})$ of (4) does not always reflect the trend of the averaged BER:

$$\bar{P}_{e, \text{bit}} \leq \mathbb{E} \left\{ \sum_{i=0}^{I-1} \sum_{i'=0}^{I-1} \frac{d_H(i, i')}{I \log_2 I} p(\mathbf{X}^i \rightarrow \mathbf{X}^{i'}) \right\}, \quad (19)$$

where $d_H(i, i')$ refers to the Hamming distance between the bit-mappings of \mathbf{X}^i and $\mathbf{X}^{i'}$. Against this background, the *average diversity product* was conceived based on (19) in [40] as:

$$\bar{\Lambda}_p = \frac{1}{2} \left(\frac{2}{\log_2 I + 1} \sum_{\forall i} \sum_{\forall i' \neq i} \frac{d_H(i, i')}{I \log_2 I \det(\Delta)} \right)^{-\frac{1}{2T}}. \quad (20)$$

The performance results of DSTSK-DAST ($M=T=2$, $Q=4$, $L=16$) associated with different Λ_p and $\bar{\Lambda}_p$ are exemplified in Fig. 3, which demonstrates that when both $\mathbf{u} = [25, 39]$ and $\mathbf{u} = [15, 63]$ achieve the maximum of $\Lambda_p = 0.1815$, the latter choice associated with a higher $\bar{\Lambda}_p = 0.2026$ exhibits a better performance. *Therefore, we propose to use $\bar{\Lambda}_p$ of (20) as the secondary objective function for all diversity-gain-maximizing DSTM schemes.* Explicitly, when several candidates \mathbf{u} achieve the same maximum Λ_p , the one associated with the highest $\bar{\Lambda}_p$ is chosen.

IV. DIFFERENTIAL SPACE-TIME SHIFT KEYING USING THREADED ALGEBRAIC SPACE-TIME (DSTSK-TAST)

In order to achieve a further improved performance, we propose a new DSTSK-TAST. More explicitly, we firstly represent the TAST [10] in the LDC form in Sec. IV-A, and then transform it into the STSK form in Sec. IV-B. The resultant finite-cardinality DSTSK-TAST invokes the single-stream DSTSK detectors in Sec. IV-C, and its diversity gain is maximized in Sec. IV-D. Finally, the diversity product of the proposed DSTSK-TAST is compared to the existing diversity-oriented finite-cardinality DSTM schemes in Sec. IV-E.

TAST:			
1	4	3	2
2	1	4	3
3	2	1	4
4	3	2	1

DSTSK-TAST:			
$\tau=1$	$\tau=2$	$\tau=3$	$\tau=4$
1			
2			
3			
4			

Fig. 4: Examples of TAST signal matrix [10] and DSTSK-TAST signal matrix for $M=T=4$.

A. The Relationship Between TAST and LDC

Instead of only using the diagonal matrix, the TAST in [10] divides a $(T \times T)$ -element space-time matrix into T non-overlapping threaded layers, where each threaded layer transmits a dispersed sequence $\bar{\mathbf{x}}_\tau = \{\mathbf{G}_A \mathbf{x}_\tau\}_{\tau=1}^T$. An example of the TAST matrix associated with $(T=4)$ is portrayed in Fig. 4. More explicitly, for the PSK signalling, the TAST is constructed by [10]:

$$\mathbf{X} = \sum_{\tau=1}^T \phi_r^{\tau-1} \text{diag}(\mathbf{G}_A \mathbf{x}_\tau) \mathbf{G}_r^{\tau-1}, \quad (21)$$

where the phase rotations $\{\phi_r^{\tau-1} = w_{L_r}^{\tau-1}\}_{\tau=1}^T$ are taken from a L_r -PSK constellation, which are invoked for the sake of retaining the full diversity order. The dispersion-element generator matrix \mathbf{G}_A is given by (8). Moreover, the layer-swithing generator matrix \mathbf{G}_r is given by:

$$\mathbf{G}_r = \begin{bmatrix} 0 & \cdots & 0 & 1 \\ 1 & \cdots & 0 & 0 \\ \vdots & \ddots & \vdots & \vdots \\ 0 & \cdots & 1 & 0 \end{bmatrix}. \quad (22)$$

Similar to the DAST of (10), the TAST of (21) may also be expressed in the form of LDC:

$$\mathbf{X} = \sum_{\tau=1}^T \sum_{q=1}^Q x_{\tau,q} \mathbf{A}_{\tau,q}, \quad (23)$$

where the dispersion matrices are $\{\{\mathbf{A}_{\tau,q} = \phi_r^{\tau-1} \text{diag}([g_1^{q-1}, \dots, g_T^{q-1}]) \mathbf{G}_r^{-1}\}_{\tau=1}^T\}_{q=1}^Q$, while the TAST in [10] assumes $Q=T$.

B. Finite-Cardinality Design

For the sake of forming a finite-cardinality signals set, we now propose the DSTSK-TAST, which revises the LDC's multiplexing form of (23) to the STSK's activation form [14]:

$$\mathbf{X} = x^l \mathbf{A}_{\tau,q}, \quad (24)$$

where $\log_2 L$, $\log_2 T$ and $\log_2 Q$ number of bits are separately assigned to the L -PSK index l , the layer activation index τ and the dispersion activation index q , respectively. Based on the DSTSK-DAST design of (11), the dispersion matrices in (24) are now given by $\{\{\mathbf{A}_{\tau,q} = \phi_r^{\tau-1} \text{diag}([w_{L_{DM}}^{(q-1)u_1}, \dots, w_{L_{DM}}^{(q-1)u_T}]) \mathbf{G}_r^{-1}\}_{\tau=1}^T\}_{q=1}^Q$. Therefore, the throughput of DSTSK-TAST is $R = \frac{\log_2 L + \log_2 Q + \log_2 T}{T}$, where $\log_2 T$ extra bits are conveyed compared to DSTSK-DAST.

Moreover, (24) requires that $L_r \geq \max\{LT, LQ, L_{DM}\}$ for creating distinct codewords. Once again, for the sake of simplicity, we generally choose $L_{DM} = LQ$ and $L_r = \max\{LT, L_{DM}\}$.

C. Single-Stream ML Detection

Similar to DSTSK-DAST, the ML detection of (3) may be decoupled for DSTSK-TAST as $\hat{\mathbf{X}}_{n-1} = \arg \max_{\forall \tau, q, l} \Re[(x^l)^* \tilde{z}_{\tau, q}]$, where we have $\tilde{z}_{\tau, q} = \text{tr}(\bar{\mathbf{A}}_{\tau, q}^H \mathbf{Z})$, which may be further simplified as $\tilde{z}_{\tau, q} = (\phi_r^{\tau-1})^* \sum_{t=1}^T w_{L_{DM}}^{-(q-1)u_t} z_{\bar{t}, t}$ associated with $\bar{t} = \{[T+(t-1) - (\tau-1)] \bmod T\} + 1$ and $\bar{\tau} = \{[T-(\tau-1)] \bmod T\} + 1$. As a result, the L -PSK demodulator is invoked TQ times as $\{\{\hat{l}_{\tau, q} = \mathbb{M}_{\text{PSK}}^{-1}(\tilde{z}_{\tau, q})\}_{\tau=1}^T\}_{q=1}^Q$. Then the optimum activation indices may be detected as $(\hat{\tau}, \hat{q}) = \arg \max_{\forall \tau, q} \Re[(x^{\hat{l}_{\tau, q}})^* \tilde{z}_{\tau, q}]$. Finally, the optimum L -PSK index is directly given by $\hat{l} = \hat{l}_{\hat{\tau}, \hat{q}}$. In summary, the DSTSK-TAST's single-stream ML detection complexity order is given by $\mathcal{O}(TQ)$.

D. Diversity Gain Maximization

Since DSTSK-TAST forms a finite group, its diversity product of (5a) may be expressed as:

$$\Lambda_p = \frac{1}{2} \min_{\forall i \neq 0} \det[(\mathbf{I}_T - x^i \mathbf{A}_{\tau, q})^H (\mathbf{I}_T - x^i \mathbf{A}_{\tau, q})]^{\frac{1}{2T}}. \quad (25)$$

It becomes more difficult to further simplify (25) because of the associated layer-swithing. Nonetheless, we offer the following two propositions for the cases of $T = 2$ and $T = 4$:

Proposition 2: For $T = 2$, the DSTSK-TAST's diversity product of (25) is given by:

$$\begin{aligned} \Lambda_p &= \min_{\tau=\{1,2\}} \Lambda_p^\tau, \text{ where} \\ \Lambda_p^{\tau=1} &= \min_{\forall (l \neq 0 || q \neq 1)} \prod_{t=1}^T \left| \sin \left[\frac{\pi(qu_t - u_t + lL_{DM}/L)}{L_{DM}} \right] \right|^{\frac{1}{T}} \\ \Lambda_p^{\tau=2} &= \min_{\forall l \forall q \forall (\tau \neq 0)} \frac{\sqrt{2}}{2} \left| \sin \left[\frac{(q-1)(\sum_{t=1}^T u_t)\pi}{L_{DM}} + \frac{T\pi l}{L} + \frac{T\pi(\tau-1)}{L_r} \right] \right|^{\frac{1}{T}}. \end{aligned} \quad (26)$$

Proof. For the case of $\tau = 1$, DSTSK-TAST's signal matrix of (24) is the same as DSTSK-DAST of (11), hence $\Lambda_p^{\tau=1}$ of (26) is given by (17) in Proposition 1. For $\tau = 2$, the determinant term in (25) is further extended as:

$$\begin{aligned} \det(\Delta) &= 4 - |\phi_r^* w_{L_{DM}}^{-(q-1)u_1} w_L^{-l} + \phi_r w_{L_{DM}}^{(q-1)u_2} w_L^l|^2 \\ &= 4 \sin^2 \left[\frac{(q-1)(u_1+u_2)\pi}{L_{DM}} + \frac{2\pi l}{L} + \frac{2\pi(\tau-1)}{L_r} \right], \end{aligned} \quad (27)$$

which results in $\Lambda_p^{\tau=2}$ of (26) associated with $T = 2$. \square

Proposition 3: For $T = 4$, the DSTSK-TAST's diversity product of (25) is given by $\Lambda_p = \min_{\tau=\{1,2,3\}} \Lambda_p^\tau$, where $\Lambda_p^{\tau=1}$ and $\Lambda_p^{\tau=2}$ are given by (26) associated with $T = 4$. Moreover, the third term $\Lambda_p^{\tau=3}$ is given by:

$$\begin{aligned} \Lambda_p^{\tau=3} &= \frac{1}{\sqrt{2}} \left| \sin \left[\frac{(q-1)(u_1+u_3)\pi}{L_{DM}} + \frac{2\pi l}{L} + \frac{2\pi}{L_r} \right] \right|^{1/4} \\ &\quad \left| \sin \left[\frac{(q-1)(u_2+u_4)\pi}{L_{DM}} + \frac{2\pi l}{L} + \frac{2\pi}{L_r} \right] \right|^{1/4}. \end{aligned} \quad (28)$$

Proof. For the cases of $\tau = 1$ and $\tau = 2$, $\Lambda_p^{\tau=1}$ and $\Lambda_p^{\tau=2}$ may be obtained in the same way following the proof of Proposition 2. For the case of $\tau = 3$, we have \mathbf{G}_r^3 in (24), which implies that the activated layer is shifted twice away

TABLE VI: Example of DSTSK-TAST ($M = T = 2$) using $Q = 2$, $L = 2$, $L_{DM} = 4$, $L_r = 4$ and $\mathbf{u} = [1, 3]$ at $R = 1.5$.

Input bits	PSK symbol	Index τ	Index q	Dispersion matrix	Data-carrying matrix
000	$x^0=1$	$\tau=1$	$q=1$	$\mathbf{A}_{1,1} = \begin{bmatrix} 1 & 0 \\ 0 & 1 \end{bmatrix}$	$\mathbf{X} = x^0 \mathbf{A}_{1,1} = \begin{bmatrix} 1 & 0 \\ 0 & 1 \end{bmatrix}$
001	$x^0=1$	$\tau=1$	$q=2$	$\mathbf{A}_{1,2} = \begin{bmatrix} j & 0 \\ 0 & -j \end{bmatrix}$	$\mathbf{X} = x^0 \mathbf{A}_{1,2} = \begin{bmatrix} j & 0 \\ 0 & -j \end{bmatrix}$
010	$x^0=1$	$\tau=2$	$q=1$	$\mathbf{A}_{2,1} = \begin{bmatrix} 0 & j \\ j & 0 \end{bmatrix}$	$\mathbf{X} = x^0 \mathbf{A}_{2,1} = \begin{bmatrix} 0 & j \\ j & 0 \end{bmatrix}$
011	$x^0=1$	$\tau=2$	$q=2$	$\mathbf{A}_{2,2} = \begin{bmatrix} 0 & -1 \\ 1 & 0 \end{bmatrix}$	$\mathbf{X} = x^0 \mathbf{A}_{2,2} = \begin{bmatrix} 0 & -1 \\ 1 & 0 \end{bmatrix}$
100	$x^1=-1$	$\tau=1$	$q=1$	$\mathbf{A}_{1,1} = \begin{bmatrix} 1 & 0 \\ 0 & 1 \end{bmatrix}$	$\mathbf{X} = x^1 \mathbf{A}_{1,1} = \begin{bmatrix} -1 & 0 \\ 0 & -1 \end{bmatrix}$
101	$x^1=-1$	$\tau=1$	$q=2$	$\mathbf{A}_{1,2} = \begin{bmatrix} j & 0 \\ 0 & -j \end{bmatrix}$	$\mathbf{X} = x^1 \mathbf{A}_{1,2} = \begin{bmatrix} -j & 0 \\ 0 & j \end{bmatrix}$
110	$x^1=-1$	$\tau=2$	$q=1$	$\mathbf{A}_{2,1} = \begin{bmatrix} 0 & j \\ j & 0 \end{bmatrix}$	$\mathbf{X} = x^1 \mathbf{A}_{2,1} = \begin{bmatrix} 0 & -j \\ -j & 0 \end{bmatrix}$
111	$x^1=-1$	$\tau=2$	$q=2$	$\mathbf{A}_{2,2} = \begin{bmatrix} 0 & -1 \\ 1 & 0 \end{bmatrix}$	$\mathbf{X} = x^1 \mathbf{A}_{2,2} = \begin{bmatrix} 0 & 1 \\ -1 & 0 \end{bmatrix}$

from the diagonal line. The corresponding determinant term in (25) may be evaluated by:

$$\det(\Delta) = [4 - |1 + \phi_r^2 w_{L_{DM}}^{(q-1)(u_1+u_3)} w_L^{2l}|^2] [4 - |1 + \phi_r^2 w_{L_{DM}}^{(q-1)(u_2+u_4)} w_L^{2l}|^2], \quad (29)$$

which results in $\Lambda_p^{\tau=3}$ of (28). Moreover, for the case of $\tau = 4$, we have $\mathbf{G}_r^3 = (\mathbf{G}_r)^T$ in (24), which is equivalent to the case of using $\tau = 2$ associated with \mathbf{G}_r . \square

Consequently, similar to DSTSK-DAST, the DSTSK-TAST's integer phase-rotation parameters $\{1 \leq u_t \leq L_{DM} - 1\}_{t=1}^T$ may also be obtained by brute-force exhaustive search for maximizing the diversity product of (25). However, unlike DSTSK-DAST, the ordering of $\{u_t\}_{t=1}^T$ may change Λ_p for DSTSK-TAST due to the phase rotations of $\phi_r^{\tau-1}$. Nonetheless, the DSTSK-TAST's search space may still be reduced to $(\frac{L_{DM}}{2})^T$ number of combinations for candidate integers $\{u_t\}_{t=1}^T$, which are still supposed to be relatively prime to L_{DM} .

Example 5: Let us now consider the DSTSK-TAST using $M = T = 2$ at the throughput of $R = \frac{\log_2 LTQ}{T} = 1.5$, where we have $L_{DM} = LQ = 4$. Since the DSTSK-TAST detection complexity order in Sec. IV-C is given by $\mathcal{O}(TQ)$, we start from the parameter combination of $Q = 1$, $L = 4$ and $L_r = \max\{LT, L_{DM}\} = 8$, and the exhaustive search gives $\mathbf{u} = [1, 1]$ associated with $\Lambda_p = 0.5946$, which is the same as the diversity product of its DGC counterpart of Example 1 and that of its DSTSK-DAST counterpart of Example 3. Furthermore, once we further increase Q and use the updated combination of $Q = 2$, $L = 2$ and $L_r = \max\{LT, L_{DM}\} = 4$, the exhaustive search gives $\mathbf{u} = [1, 3]$ achieving $\Lambda_p = 0.7071$, which is higher than that of both DGC and DSTSK-DAST. The corresponding DSTSK-TAST signal matrices are summarized in Table VI. Specifically, compared to the DGC of Example 1, DSTSK-TAST achieves a higher $\Lambda_p = 0.7071$ at a lower detection complexity order of $\mathcal{O}(TQ = 4)$.

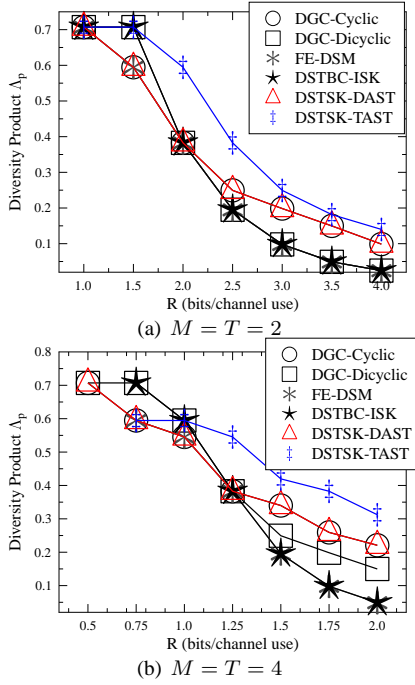


Fig. 5: Comparison of diversity products of finite-cardinality DSTM schemes.

Example 6: For a DSTSK-TAST using $M = T = 4$ at $R = 1.25$, we obtain the parameters of $Q = 2$, $L = 4$, $L_{DM} = 8$, $L_r = 16$ and $\mathbf{u} = [1, 3, 7, 5]$, which achieves $\Lambda_p = 0.5453$ that is higher than $\Lambda_p = 0.3827$ of the DGC and DSTSK-DAST of Examples 2 and 4. The resultant DSTSK-TAST signal matrices are given by $\{\{\{\mathbf{X}^i = w_4^l \mathbf{A}_{\tau,q} \}_{l=0}^3 \}_{\tau=1}^4 \}_{q=1}^2$, where the dispersion matrix is given by $\mathbf{A}_{\tau,q} = w_{16}^{\tau-1} \text{diag}([w_8^{(q-1)}, w_8^{3(q-1)}, w_8^{7(q-1)}, w_8^{5(q-1)}]) \mathbf{G}_r^{\tau-1}$. The associated detection complexity order is $\mathcal{O}(TQ = 8)$, which is lower than the DGC's $\mathcal{O}(I = 32)$ in Example 2.

E. Comparison of Diversity Products

All of the diversity-oriented finite-cardinality DSTM schemes are summarized in Table VII, which may be categorized into three types. The first type is the classic cyclic and dicyclic DGCs, which actively maximize their diversity products by using the parameters $\mathbf{u} = [u_1, \dots, u_T]$, but the DGC detection complexity grows exponentially with the throughput, as demonstrated by Table VII. The second type is the full-diversity FE-DSM and DSTBC-ISK schemes that achieve a non-zero Λ_p in Table VII. However, in the absence of diversity gain maximization, both FE-DSM and DSTBC-ISK generally suffer from a performance loss compared to the DGCs as the throughput increases. Nonetheless, as a benefit, both FE-DSM and DSTBC-ISK are capable of invoking the single-stream ML detection. Finally, the third type is constituted by our proposed DSTSK-DAST and DSTSK-TAST featured in Table VII, which are capable of achieving an actively maximized diversity gain at the low-cost of a single-stream ML detection complexity.

The diversity products of (5a) achieved by the above DSTM schemes are compared in Fig. 5, where the parameters are summarized in Table VIII. Firstly, Fig. 5 evidences that DSTSK-ISK achieves higher diversity products than FE-DSM

at $R = 1.5$ for $M = T = 2$ and $R = \{0.75, 1.0\}$ for $M = T = 4$, but the diversity products of both DSTSK-ISK and FE-DSM decrease rapidly and become substantially lower than that of DGC-cyclic as the throughput increases. Secondly, it is evidenced by Fig. 5 that the proposed DSTSK-DAST is capable of achieving exactly the same diversity gains as DGC-cyclic, while the proposed DSTSK-TAST is capable of achieving even further improved diversity gains, apart from the low-throughput exceptions of $R = 1$ for $M = T = 2$ and $R = 0.75$ for $M = T = 4$, where the diversity products of DSTSK-TAST, DSTSK-DAST and DGC-cyclic are the same. Once again, the diversity advantages of DSTSK-DAST and DSTSK-TAST are achieved at a substantially lower detection complexity than that of DGC. Moreover, we also note that the better-performing DSTSK-TAST has a higher detection complexity order of $\mathcal{O}(QT)$ than $\mathcal{O}(Q)$ of DSTSK-DAST, despite the fact that neither of the two single-stream ML detection complexities increase with the constellation size. Our more detailed performance versus complexity comparisons will be presented in Sec. VI.

V. AMPLITUDE-PHASE SHIFT KEYING (APSK) AND DIVERSITY-RATE (DR) TRADEOFF

In order to improve the throughput, in Sec. V-A, we further extend the SE-APSK design of [40] to the ME-APSK, while a generic DR arrangement is conceived for a practical DSTM diversity-throughput tradeoff design in Sec. V-B.

A. Multi-Element (ME) Amplitude-Phase Shift Keying (APSK)

According to [40], the generic L_A -level ring-amplitude is applied to the PSK signals in (1) by the Differential Amplitude Shift Keying (DASK) as $\Gamma_n = \gamma_{n-1} \Gamma_{n-1}$, where $\{\Gamma_{n-1} = \frac{\alpha^{\mu_{n-1}}}{\sqrt{\beta}}, \Gamma_n = \frac{\alpha^{\mu_n}}{\sqrt{\beta}}\}$ and $\{\gamma_{n-1} = \alpha^a\}_{a=0}^{L_A-1}$ represent the transmitted ring-amplitudes and data-carrying ring-amplitude, respectively, while α and $\beta = \frac{\sum_{\mu=0}^{L_A-1} \alpha^{2\mu}}{L_A}$ respectively represent the ring ratio and the associated normalization factor. More explicitly, the indices of the transmitted ring-amplitudes $\{\Gamma_{n-1}, \Gamma_n\}$ are given by the integers $\{\mu_{n-1}, \mu_n\} \in [0, L_A - 1]$, while the data-carrying ring-amplitude modulates the difference between μ_{n-1} and μ_n as $a = [(\mu_{n-1} + \tilde{a}) \bmod L_A] - \mu_{n-1}$, where a is the Gray coded version of the natural index \tilde{a} .

In order to further extend the SE design of [40] and the two/four-ring ME design of [33], [34], the differential encoding of (1) may be revised for the generic ME-APSK as [33], [34]:

$$\tilde{\mathbf{S}}_n = \mathbf{X}_{n-1} \tilde{\mathbf{S}}_{n-1} \mathbf{\Lambda}_{n-1}, \quad (30)$$

where the $(T \times T)$ -element data-carrying ring-amplitude matrix is $\mathbf{\Lambda}_{n-1} = \text{diag}([\gamma_{n-1,1}, \dots, \gamma_{n-1,T}])$, which forms a star-QAM constellation for $\tilde{\mathbf{S}}_{n-1} = \mathbf{S}_{n-1} \text{diag}([\Gamma_{n-1,1}, \dots, \Gamma_{n-1,T}])$ and $\tilde{\mathbf{S}}_n = \mathbf{S}_n \text{diag}([\Gamma_{n,1}, \dots, \Gamma_{n,T}])$ owing to the DASK of $\{\Gamma_{n,t} = \gamma_{n-1,t} \Gamma_{n-1,t}\}_{t=1}^T$.

It was proven in [34] that for a single-RF sparse signal matrix $\tilde{\mathbf{S}}_{n-1}$, we have the relationship of $\tilde{\mathbf{S}}_{n-1} \mathbf{\Lambda}_{n-1} = \mathbf{\Lambda}_{n-1} \tilde{\mathbf{S}}_{n-1}$, where the permuted matrix is given by $\mathbf{\Lambda}_{n-1} =$

TABLE VII: Summary of diversity-oriented finite-cardinality DSTM schemes.

DSTM scheme	Diversity product	Throughput	Detection complexity
DGC-cyclic	$\Lambda_p = \min_{\forall l \neq 0} \prod_{t=1}^T \sin(\frac{\pi}{L} u_t l) ^{1/T}$	$R = \frac{\log_2 L}{T}$	$\mathcal{O}(L = 2^{RT})$
DGC-dicyclic	$\Lambda_p = \min\{\frac{1}{\sqrt{2}}, \min_{\forall l \neq 0} \prod_{t=1}^{T/2} \sin(\frac{\pi}{L} u_t l) ^{2/T}\}$	$R = \frac{\log_2 2L}{T}$	$\mathcal{O}(2L = 2^{RT})$
FE-DSM	$\Lambda_p = \min_{\forall (l \neq 0 q \neq 1)} \prod_{t=1}^T \sin(\frac{\pi}{L} l + \frac{\pi(L-L+1)}{LT}(q-1) $	$R = \frac{\log_2 L + \log_2 T}{T}$	$\mathcal{O}(T)$
DSTBC-ISK	$T = 2: \Lambda_p = \min\{\sin(\pi/L), \frac{1}{\sqrt{2}}\}$	$R = \frac{\log_2 L + \log_2 T}{T}$	$\mathcal{O}(T)$
	$T \geq 4: \Lambda_p = \min\{\sin(\pi/L), \frac{1}{\sqrt{2}} \sqrt{\sin(\frac{4\pi}{QL})}\}$		
DSTSK-DAST	$\Lambda_p = \min_{\forall (l \neq 0 q \neq 1)} \prod_{t=1}^T \sin(\frac{\pi(qu_t - u_t + lL_{DM}/L)}{L_{DM}}) ^{1/T}$	$R = \frac{\log_2 L + \log_2 Q}{T}$	$\mathcal{O}(Q)$
DSTSK-TAST	$\Lambda_p = \frac{1}{2} \min_{\forall i \neq 0} \det[(\mathbf{I}_T - \mathbf{x}^i \mathbf{A}_{\tau,q})^H (\mathbf{I}_T - \mathbf{x}^i \mathbf{A}_{\tau,q})]^{1/2T}$ where $\mathbf{A}_{\tau,q} = \phi_r^{\tau-1} \text{diag}([w_{L_{DM}}^{(q-1)u_1}, \dots, w_{L_{DM}}^{(q-1)u_T}]) \mathbf{G}_r^{\tau-1}$	$R = \frac{\log_2 L + \log_2 Q + \log_2 T}{T}$	$\mathcal{O}(QT)$

TABLE VIII: Parameters of DSTSK-TAST, DSTSK-DAST and DGC-cyclic used for Fig. 5.

(a) DSTSK-TAST, $M = T = 2$			
$R = 1.0: Q = 1, L = 2, L_{DM} = 2, L_r = 4, u = [1, 1]$	$R = 1.5: Q = 2, L = 2, L_{DM} = 4, L_r = 4, u = [1, 3]$	$R = 2.0: Q = 4, L = 2, L_{DM} = 8, L_r = 8, u = [1, 3]$	$R = 2.5: Q = 2, L = 8, L_{DM} = 16, L_r = 16, u = [3, 13]$
$R = 3.0: Q = 4, L = 8, L_{DM} = 32, L_r = 32, u = [29, 3]$	$R = 3.5: Q = 4, L = 16, L_{DM} = 64, L_r = 64, u = [7, 57]$		
$R = 4.0: Q = 8, L = 16, L_{DM} = 128, L_r = 128, u = [3, 117]$			
(b) DSTSK-DAST, $M = T = 2$			
$R = 1.0: Q = 1, L = 4, L_{DM} = 4, u = [1, 1]$	$R = 1.5: Q = 2, L = 4, L_{DM} = 8, u = [1, 5]$		
$R = 2.0: Q = 2, L = 8, L_{DM} = 16, u = [1, 9]$	$R = 2.5: Q = 4, L = 8, L_{DM} = 32, u = [1, 25]$		
$R = 3.0: Q = 16, L = 4, L_{DM} = 64, u = [1, 37]$	$R = 3.5: Q = 8, L = 16, L_{DM} = 128, u = [1, 81]$		
$R = 4.0: Q = 64, L = 4, L_{DM} = 256, u = [75, 255]$			
(c) DGC-cyclic, $M = T = 2$			
$R = 1.0: L = 4, u = [1, 1]$	$R = 1.5: L = 8, u = [1, 3]$	$R = 2.0: L = 16, u = [1, 7]$	$R = 2.5: L = 32, u = [1, 7]$
$R = 3.0: L = 64, u = [1, 19]$	$R = 3.5: L = 128, u = [1, 47]$	$R = 4.0: L = 256, u = [1, 75]$	
(d) DSTSK-TAST, $M = T = 4$			
$R = 0.75: Q = 1, L = 2, L_{DM} = 2, L_r = 8, u = [1, 1, 1, 1]$	$R = 1.0: Q = 2, L = 2, L_{DM} = 4, L_r = 8, u = [1, 1, 3, 3]$		
$R = 1.25: Q = 2, L = 4, L_{DM} = 8, L_r = 16, u = [1, 3, 7, 5]$	$R = 1.5: Q = 4, L = 4, L_{DM} = 16, L_r = 32, u = [1, 5, 15, 7]$		
$R = 1.75: Q = 4, L = 8, L_{DM} = 32, L_r = 32, u = [5, 25, 11, 31]$	$R = 2.0: Q = 8, L = 8, L_{DM} = 64, L_r = 64, u = [15, 41, 57, 63]$		
(e) DSTSK-DAST, $M = T = 4$			
$R = 0.5: Q = 1, L = 4, L_{DM} = 4, u = [1, 1, 1, 1]$	$R = 0.75: Q = 2, L = 4, L_{DM} = 8, u = [1, 1, 5, 5]$		
$R = 1.0: Q = 4, L = 4, L_{DM} = 16, u = [1, 5, 9, 13]$	$R = 1.25: Q = 4, L = 8, L_{DM} = 32, u = [7, 15, 23, 31]$		
$R = 1.5: Q = 16, L = 4, L_{DM} = 64, u = [1, 29, 37, 49]$	$R = 1.75: Q = 32, L = 4, L_{DM} = 128, u = [83, 91, 103, 127]$		
$R = 2.0: Q = 64, L = 4, L_{DM} = 256, u = [1, 41, 137, 221]$			
(f) DGC-cyclic, $M = T = 4$			
$R = 0.5: L = 4, u = [1, 1, 1, 1]$	$R = 0.75: L = 8, u = [1, 1, 3, 3]$	$R = 1.0: L = 16, u = [1, 3, 5, 7]$	
$R = 1.25: L = 32, u = [1, 7, 9, 15]$	$R = 1.5: L = 64, u = [1, 15, 27, 29]$	$R = 1.75: L = 128, u = [1, 25, 37, 45]$	
$R = 2.0: L = 256, u = [1, 35, 41, 119]$			

$\text{diag}([\gamma_{n-1,t_1}, \dots, \gamma_{n-1,t_T}])$, as $\{t_\ell\}_{\ell=1}^T$ denotes the activation index of ℓ -th row in \mathbf{S}_{n-1} . As a result, the received signal model of (2) may be further extended as $\mathbf{Y}_n = \mathbf{X}_{n-1} \bar{\mathbf{A}}_{n-1} (\mathbf{Y}_{n-1} - \mathbf{V}_{n-1}) + \mathbf{V}_n$ under the assumption of $\mathbf{H}_{n-1} = \mathbf{H}_n$, which leads to the following ME-APSK detection:

$$\{\hat{\mathbf{X}}_{n-1}, \hat{\bar{\mathbf{A}}}_{n-1}\} = \arg \min_{\substack{\forall \mathbf{X}_{n-1} \in \{\mathbf{X}^i\}_{i=0}^{I-1} \\ \forall \bar{\mathbf{A}}_{n-1} \in \{\bar{\mathbf{A}}^{a_1, \dots, a_T}\}_{a_1, \dots, a_T \in [0, L_A-1]}} \|\mathbf{Y}_n - \mathbf{X}_{n-1} \bar{\mathbf{A}}_{n-1} \mathbf{Y}_{n-1}\|^2. \quad (31)$$

The ME-DASK design of (30) improves the DSTM throughput from $R = \frac{\log_2 I}{T}$ to $R = \frac{\log_2 I + T \log_2 L_A}{T}$, but the detection complexity of (31) also grows exponentially with the throughput as $\mathcal{O}(IL_A^T)$.

Against this background, we further propose to invoke the single-stream ML detection for our ME-APSK applications. More explicitly, the decision metric of (31) may be further extended as:

$$\begin{aligned} \|\mathbf{Y}_n - \mathbf{X}_{n-1} \bar{\mathbf{A}}_{n-1} \mathbf{Y}_{n-1}\|^2 &= \kappa_n^2 + \sum_{\ell=1}^T \gamma_{n-1,t_\ell}^2 \|\mathbf{Y}_{n-1}^\ell\|^2 \\ &\quad - 2\Re[\text{tr}(\bar{\mathbf{A}}_{n-1} \mathbf{X}_{n-1}^H \mathbf{Z})] \\ &= \kappa_n^2 + \sum_{\ell=1}^T \kappa_{n-1,\ell}^2 (\gamma_{n-1,t_\ell} - \bar{z}_{\ell,\ell} / \kappa_{n-1,\ell})^2 - \sum_{\ell=1}^T \bar{z}_{\ell,\ell}^2 / \kappa_{n-1,\ell}^2, \end{aligned} \quad (32)$$

where we define $\kappa_n^2 = \|\mathbf{Y}_n\|^2$ and $\mathbf{Z} = \mathbf{Y}_n \mathbf{Y}_{n-1}^H$, while \mathbf{Y}_{n-1}^ℓ denotes the ℓ -th row in \mathbf{Y}_{n-1} . We also define $\kappa_{n-1,\ell}^2 =$

TABLE IX: Summary of detection complexity of DSM using ME-APSK.

Scheme	Complexity order	Real-valued multiplications
Two-ring-amplitude design [33]	$\mathcal{O}(IL_A^T)$	$(6T^2N + 8T^2N)IL_A^T$
Four-ring-amplitude design [34]	$\mathcal{O}(IL_A^T)$	$(6T^2N + 2T^2N)IL_A^T$
New generic design	$\mathcal{O}(\bar{M})$	$4T^2N + (1 + 6T^2N + 6T^3)\bar{M}$

$\|\mathbf{Y}_{n-1}^\ell\|^2$, and $\bar{z}_{\ell,\ell}$ denotes the element on the ℓ -th row and ℓ -th column of $\bar{\mathbf{Z}} = \Re(\mathbf{X}^H \mathbf{Z})$. Owing to the fact that $(\kappa_n^2 - \sum_{\ell=1}^T \bar{z}_{\ell,\ell}^2 / \kappa_{n-1,\ell}^2)$ in (32) is a constant for the ring-amplitudes, the ME-APSK detection may be decoupled as $\{\gamma_{n-1,t_\ell} = \mathbb{M}_{\text{DASK}}^{-1}(\bar{z}_{\ell,\ell} / \kappa_{n-1,\ell}^2)\}_{\ell=1}^T$ without any performance loss, where the linearized DASK detector is given by (48) of [40]. As a result, when the ME-APSK design is applied to the finite-cardinality schemes of DSM, DSTBC-ISK, FE-DSM, DSTSK-DAST and DSTSK-TAST, both the linearized PSK detector $\mathbb{M}_{\text{PSK}}^{-1}(\cdot)$ exemplified in (15) and the DASK detector $\mathbb{M}_{\text{DASK}}^{-1}(\cdot)$ are invoked, so that the resultant single-stream ML detection complexity does not grow with the star QAM constellation size.

Let us consider the DSM example of Sec. (II-D) using ME-ASK. Firstly, in order to detect the unitary matrix \mathbf{X} in (32), according to Sec. IV of [40], the PSK detector is invoked $\bar{M}T$ times as $\{\{\hat{l}_{\bar{m},t} = \mathbb{M}_{\text{PSK}}^{-1}(\mathbf{Z}(t, \mathbf{a}_{\bar{m},t}))\}_{\bar{m}=1}^{\bar{M}}\}_{t=1}^T$, where $\mathbf{Z}(t, \mathbf{a}_{\bar{m},t})$ denotes the element on the t -th row and

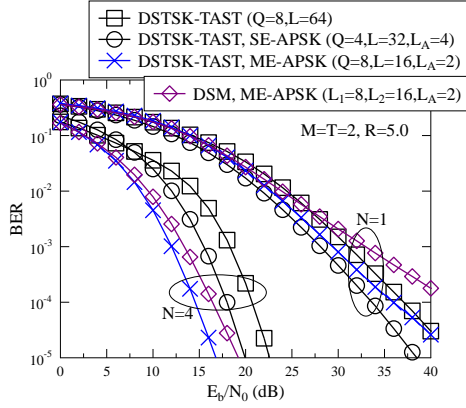


Fig. 6: Performance of DSTSK-TAST employing ME-APSK, where we have $M = T = 2$ and $R = 5.0$.

$\mathbf{a}_{\bar{m},t}$ -column in $\mathbf{Z} = \mathbf{Y}_n \mathbf{Y}_{n-1}^H$. Following this, a total of \bar{M} legitimate candidates can be formed for \mathbf{X} in (32) according to $\{\{\hat{l}_{\bar{m},t}\}_{\bar{m}=1}^{\bar{M}}\}_{t=1}^T$. Secondly, the DASK detector is invoked $\bar{M}T$ times as $\{\{\hat{a}_{\bar{m},t}\}_{\bar{m}=1}^{\bar{M}}\}_{t=1}^T$, which further forms \bar{M} candidates for $\bar{\mathbf{A}}_{n-1}$ in (32). As a result, the permutation index can be obtained by evaluating (32) for a reduced number of \bar{M} times as $\hat{m} = \arg \min_{\bar{m} \in [1, \bar{M}]} \|\mathbf{Y}_n - \mathbf{X}_{n-1} \bar{\mathbf{A}}_{n-1} \mathbf{Y}_{n-1}\|^2$. Finally, the T data-carrying PSK and ring-amplitude indices are given by $\{\hat{l}_t = \hat{l}_{\hat{m},t}\}_{t=1}^T$ and $\{\hat{a}_t = \hat{a}_{\hat{m},t}\}_{t=1}^T$, respectively. In summary, the complexity comparison between the proposed generic design and the existing ME-APSK solutions [33], [34] is shown in Table IX.

It is straightforward to see that $\bar{\mathbf{A}}_{n-1}$ of (32) does not comply with the full-rank requirement of (4), as a pair of $\bar{\mathbf{A}}_{n-1}$ and $\bar{\mathbf{A}}'_{n-1}$ associated with $\gamma_{n-1,t_1} = \gamma'_{n-1,t_1}$ and $\{\gamma_{n-1,t_i} \neq \gamma'_{n-1,t_i}\}_{i=2}^T$ has rank-1. It is exemplified by Fig. 6 that for the case of DSTSK-TAST using $N = 1$, the full-diversity SE-APSK arrangement performs better. However, Fig. 6 also demonstrates that for the case of using $N = 4$, the ME-APSK arrangement helps the DSTSK-TAST scheme once again to achieve a performance advantage over DSM.

Moreover, as a DASK counterpart, the ASK using the absolute-amplitude of $\Gamma_n = \gamma_{n-1} = \frac{\alpha}{\sqrt{\beta}}$ is capable of achieving a better performance in channel coded scenarios [40], but this is beyond the scope of this treatise. Nonetheless, for the sake of technical completeness, we note that for the applications of ME-APSK design using ASK, the DASK's $\{\gamma_{n-1,t}\}_{t=1}^T$ in $\bar{\mathbf{A}}_{n-1}$ values of (30) are replaced by $\{\gamma_{n-1,t}/\Gamma_{n-1,t}\}_{t=1}^T$, while the single-stream ASK detector is invoked as $\{\gamma_{n-1,t_i}\}_{i=1}^T = \mathbb{M}_{\text{ASK}}^{-1}(\Gamma_{n-1,t_i} \bar{z}_{i,t_i}/\kappa_{n-1,t_i}^2)\}_{i=1}^T$ for (32) according to (49) in [40].

B. Diversity-Rate Tradeoff

The DR arrangement in [37], [38] partitions the $(T \times T)$ DSTM signal space into $(T_D \times T_D)$ sub-blocks, where $V = T/T_D$ number of $(T_D \times T_D)$ DSTM sub-blocks are transmitted in a permuted manner. This is exemplified in Fig. 7 for the case of using $M = T = 4$ and $V = 2$, where two 2×2 DSTM sub-blocks are permuted in the 4×4 DSTM signal space. As a result, the DR design retains a reduced diversity order of T_D in exchange for an improved throughput of $R = \frac{\sum_{u=1}^V \log_2 I_u + \log_2 V}{T}$, where $\{I_u\}_{u=1}^V$ denotes the

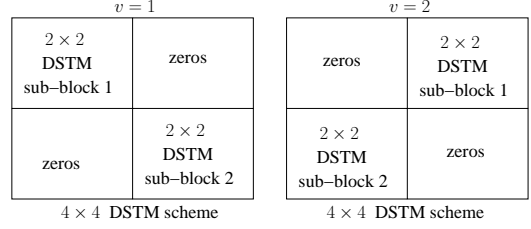


Fig. 7: Examples of a (4×4) DSTM scheme using DR associated with $M = T = 4$ and $V = 2$.

TABLE X: Pseudocode for the single-stream ML detection of DSTM-DR.

```

1: for  $v = 1$  to  $V$ 
2:    $\bar{\mathbf{Z}}^v = (\mathbf{G}_r^{v-1} \otimes \mathbf{I}_{T_D})^H \mathbf{Y}_n \mathbf{Y}_{n-1}^H$ 
3:   for  $u = 1$  to  $V$ 
4:      $\hat{i}_{v,u} = \arg \max_{i_{v,u} \in [0, I_u-1]} \Re \left[ \text{tr} \left( (\bar{\mathbf{X}}_u^{i_{v,u}})^H \bar{\mathbf{Z}}_{(u)}^v \right) \right]$ 
5:      $d_{v,u} = \Re \left[ \text{tr} \left( (\bar{\mathbf{X}}_u^{i_{v,u}})^H \bar{\mathbf{Z}}_{(u)}^v \right) \right]$ 
6:   end for
7: end for
8:  $\hat{v} = \max_{v \in \{0, V-1\}} \sum_{u=1}^V d_{v,u}$ 
9:  $\{\hat{i}_u = \hat{i}_{\hat{v},u}\}_{u=1}^V$ 

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cardinality of the u -th constituent DSTM scheme, while the permutation index $\{v\}_{v=1}^V$ carries an extra number of $\log_2 V$ bits. Furthermore, since there is no interference between the DSTM sub-blocks, the single-stream ML detection may be also invoked. Explicitly, the DR design forms the $(T \times T)$ -element signal matrix as:

$$\mathbf{X}^i = (\mathbf{G}_r^{v-1} \otimes \mathbf{I}_{T_D}) \text{diag}\{\bar{\mathbf{X}}_1^{i_1}, \dots, \bar{\mathbf{X}}_V^{i_V}\}, \quad (33)$$

where instead of using a FE matrix in [37], [38], the $(V \times V)$ -element switching matrix \mathbf{G}_r defined in (22) is invoked for the generic DSTM design. Moreover, the $(T_D \times T_D)$ -element signal matrix $\{\bar{\mathbf{X}}_u^{i_u}\}_{u=1}^V$ of (33) denotes the constituent DSTM sub-block.

As a result, upon obtaining the received signals of (2), the decision metric of (3) may be further extended for the DR detection as:

$$\|\mathbf{Y}_n - \mathbf{X}^i \mathbf{Y}_{n-1}\|^2 = \kappa_n^2 + \sum_{u=1}^V \left\{ \|\mathbf{Y}_{n-1,(u)}\|^2 - 2 \Re \left[\text{tr} \left((\bar{\mathbf{X}}_u^{i_u})^H \bar{\mathbf{Z}}_{(u)}^v \right) \right] \right\}, \quad (34)$$

where we define the $(T_D \times N)$ -element matrix $\mathbf{Y}_{n-1,(u)} = \mathbf{Y}_{n-1}([(u-1)T_D + 1 : uT_D, :])$ as the submatrix taken from the $[(u-1)T_D + 1]$ -th to the (uT_D) -th rows of \mathbf{Y}_{n-1} , and we also define the $(T_D \times T_D)$ -element matrix $\bar{\mathbf{Z}}_{(u)}^v = \bar{\mathbf{Z}}^v([(u-1)T_D + 1 : uT_D, (u-1)T_D + 1 : uT_D])$ as the submatrix taken from the $[(u-1)T_D + 1]$ -th to the (uT_D) -th rows and columns of $\bar{\mathbf{Z}}^v = (\mathbf{G}_r^{v-1} \otimes \mathbf{I}_{T_D})^H \mathbf{Y}_n \mathbf{Y}_{n-1}^H$. Moreover, we note that $\sum_{u=1}^V \|\mathbf{Y}_{n-1,(u)}\|^2 = \kappa_{n-1}^2$ in (34) is a constant. As a result, the detection of $\{\bar{\mathbf{X}}_u^{i_u}\}_{u=1}^V$ is decoupled in Eq. (34), which is detailed in the pseudocode of Table X. Moreover, when the constituent DSTM sub-blocks use FE-DSM, DSTBC-ISK as well as the proposed DSTSK-DAST and DSTSK-TAST schemes, their respective single-stream ML detector may be directly invoked for line 4 in Table X.

In order to once again introduce star-QAM signalling without compromising the diversity, we further propose the DR-APSK design, where a single ring-amplitude is assigned to each $(T_D \times T_D)$ -element DSTM sub-

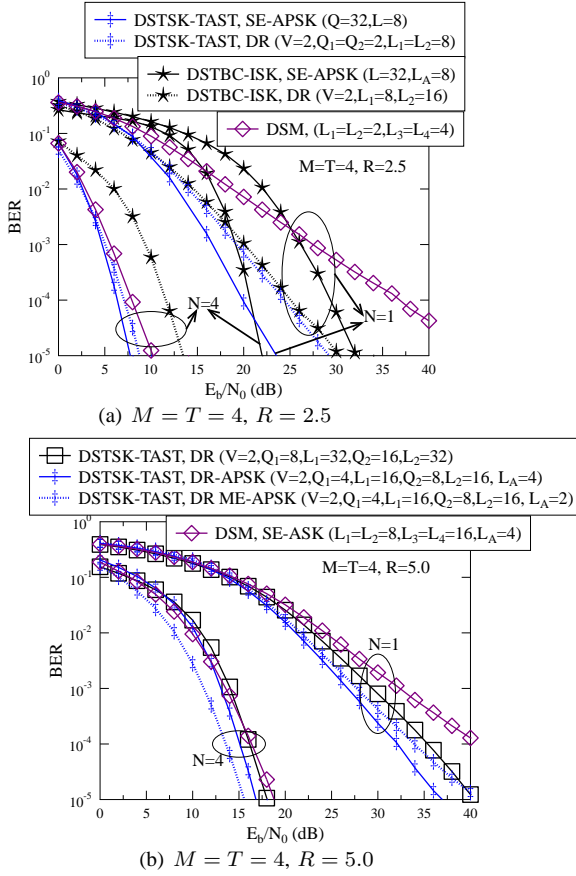


Fig. 8: Performance of DSTSK-TAST and DSTBC-ISK employing the DR and APSK arrangements, where we have $M = T = 4$.

block. More explicitly, for the DR-APSK design, the $(T \times T)$ -element ring-amplitude matrix in (30) is given by $\mathbf{\Lambda}_{n-1} = \text{diag}([\gamma_{n-1,1} \mathbf{I}_{T_D}, \dots, \gamma_{n-1,V} \mathbf{I}_{T_D}])$. The resultant star QAM signal matrices are $\tilde{\mathbf{S}}_{n-1} = \mathbf{S}_{n-1} \text{diag}([\Gamma_{n-1,1} \mathbf{I}_{T_D}, \dots, \Gamma_{n-1,V} \mathbf{I}_{T_D}])$ and $\tilde{\mathbf{S}}_n = \mathbf{S}_n \text{diag}([\Gamma_{n,1} \mathbf{I}_{T_D}, \dots, \Gamma_{n,V} \mathbf{I}_{T_D}])$, where we have $\{\Gamma_{n,u} = \gamma_{n-1,u} \Gamma_{n-1,u}\}_{u=1}^V$ for using DASK. The sparse matrix feature leads to the relationship of $\tilde{\mathbf{S}}_{n-1} \mathbf{\Lambda}_{n-1} = \tilde{\mathbf{\Lambda}}_{n-1} \tilde{\mathbf{S}}_{n-1}$, where the permuted matrix is $\tilde{\mathbf{\Lambda}}_{n-1} = \text{diag}([\gamma_{n-1,t_1} \mathbf{I}_{T_D}, \dots, \gamma_{n-1,t_V} \mathbf{I}_{T_D}])$, as $\{t_u\}_{u=1}^V$ denotes the sub-block activation index. As a result, the decision metric of (34) is revised for DR-APSK as:

$$\left\| \mathbf{Y}_n - \mathbf{X}^i \tilde{\mathbf{\Lambda}}^{a_1, \dots, a_V} \mathbf{Y}_{n-1} \right\|^2 = \kappa_n^2 + \sum_{u=1}^V \left\{ \gamma_{n-1,t_u}^2 \left\| \mathbf{Y}_{n-1,(u)} \right\|^2 - 2\gamma_{n-1,t_u} \Re \left[\text{tr} \left((\bar{\mathbf{X}}_u^{i_u})^H \bar{\mathbf{Z}}_{(u)}^v \right) \right] \right\}. \quad (35)$$

Hence, following the unitary matrix detection of line 4 in Table X, the ring-amplitude detection is also decoupled as $\{\gamma_{n-1,t_u} = \mathbb{M}_{\text{DASK}}^{-1}(\Re[\text{tr}((\bar{\mathbf{X}}_u^{i_u})^H \bar{\mathbf{Z}}_{(u)}^v)]) / \|\mathbf{Y}_{n-1,(u)}\|^2\}_{u=1}^V$, which leads to the ring-amplitude indices $\hat{a}_{v,u}$. Accordingly, the decision metric of line 5 in Table X may be revised for DR-APSK as $d_{v,u} = (\gamma_{\hat{a}_{v,u}})^2 \|\mathbf{Y}_{n-1,(u)}\|^2 - 2\gamma_{\hat{a}_{v,u}} \Re[\text{tr}((\bar{\mathbf{X}}_u^{i_{v,u}})^H \bar{\mathbf{Z}}_{(u)}^v)]$. Finally, the detected ring-amplitude indices are given by $\{\hat{a}_u = \hat{a}_{v,u}\}_{u=1}^V$ after line 9 in Table X.

Moreover, the DR design may also invoke the SE-APSK

of [40] and the ME-APSK of Sec. V-A. More explicitly, Fig. 8(a) exemplifies that for the case of $M = T = 4$ and $R = 2.5$, DSTBC-ISK using DR substantially improves the performance of DSTBC-ISK using SE-APSK, but the DR design does not appear to be beneficial for DSTSK-TAST. This is because DSTSK-TAST actively maximizes its rank-4 diversity gain for $M = T = 4$, which is more advantageous than its DR applications. Nonetheless, for the example of using $M = T = 4$ at a very high rate of $R = 5.0$ that conveys a total of 20 bits, optimizing the diversity gain for the DSTSK-TAST cardinality of 2^{20} becomes unrealistic. Against this background, the DSTSK-TAST using DR may invoke a pair of constituent DSTSK-TAST schemes of $(Q_1, L_1) = (8, 32)$ and $(Q_1, L_1) = (16, 32)$ that are optimized for $M = T = 2$ at $R = 4.5$ and $R = 5.0$, respectively. As a result, it is confirmed in Fig. 8(b) that for the case of $M = T = 4$ and $R = 5.0$, the DSTSK-TAST using DR-ASK achieves the best performance for $N = 1$, while the DSTSK-TAST using DR and ME-ASK ensures a further diversity gain over DSM for $N = 4$.

VI. PERFORMANCE RESULTS

First of all, in Sec. VI-A, the performance results of the full-diversity design are presented, where $M = \{2, 4\}$ TAs and a small number of $N = \{1, 2\}$ RAs are used. Secondly, the diversity-throughput tradeoff is further discussed in Sec. VI-B using both small and large number of $N = \{1, 2, 4, 8, 16\}$ RAs.

A. Performance Results of Full-Diversity Design

The parameters of DSTSK-TAST, DSTSK-DAST and DGC-cyclic at high throughputs used in this section are summarized in Table XI, where the parameters of the low-throughput arrangements are the ones given in Table VIII. The parameters for DSTBC-ISK, FE-DSM and DSM can be retrieved from [31], [32], [38], [40]. Moreover, the star-QAM ring-amplitude ratios of $\alpha = \{2.0, 1.4, 1.2, 1.1\}$ are used for $L_A = \{2, 4, 8, 16\}$, respectively.

The performance comparisons are portrayed both in terms of their DCMC capacity and their BER performance in Fig. 9 for the case of $M = T = 2$ and $R = 2.0$. The evaluation of DSTM's DCMC capacity can be found in [40]. Fig. 9 demonstrates that the proposed DSTSK-DAST exhibits a similar performance to the classic DGC-cyclic, where both DSTSK-DAST and DGC-cyclic achieve substantial diversity gains over DSM. Furthermore, it is also evidenced by Fig. 9 that the new DSTSK-TAST achieves an even further improved performance over DGC-cyclic. We once again note that the diversity advantages of DSTSK-DAST and DSTSK-TAST are achieved at a substantially lower detection complexity than DGC-cyclic.

The performance results of the finite-cardinality DSTM schemes and their SISO counterparts of DPSK/ADPSK are summarized in terms of the E_b/N_0 required for achieving $\text{BER}=10^{-4}$ in Fig. 10 for the case of $M = T = \{2, 4\}$ and $N = 1$, while the corresponding complexity comparisons are portrayed in Fig. 11. The complexities are evaluated in terms of the total number of real-valued multiplications, which are

TABLE XI: Parameters of DSTSK-TAST, DSTSK-DAST and DGC-cyclic used in Sec. VI-A.

(a) DSTSK-TAST, SE-APSK, $M = T = 2$

$R=3.0$: $Q=2$, $L=8$, $L_{DM}=16$, $L_r=16$, $u=[3, 13]$, $L_A=2$	$R=3.5$: $Q=4$, $L=8$, $L_{DM}=32$, $L_r=32$, $u=[29, 3]$, $L_A=2$
$R=4.0$: $Q=4$, $L=8$, $L_{DM}=32$, $L_r=32$, $u=[29, 3]$, $L_A=4$	$R=4.5$: $Q=4$, $L=16$, $L_{DM}=64$, $L_r=64$, $u=[7, 57]$, $L_A=4$
$R=5.0$: $Q=4$, $L=32$, $L_{DM}=128$, $L_r=128$, $u=[127, 97]$, $L_A=4$	$R=5.5$: $Q=4$, $L=32$, $L_{DM}=128$, $L_r=128$, $u=[127, 97]$, $L_A=8$
$R=6.0$: $Q=8$, $L=32$, $L_{DM}=256$, $L_r=256$, $u=[11, 237]$, $L_A=8$	

(b) DSTSK-DAST, SE-APSK, $M = T = 2$

$R=2.5$: $Q=2$, $L=8$, $L_{DM}=16$, $u=[1, 9]$, $L_A=2$	$R=3.0$: $Q=4$, $L=8$, $L_{DM}=32$, $u=[1, 25]$, $L_A=2$
$R=3.5$: $Q=4$, $L=16$, $L_{DM}=64$, $u=[15, 63]$, $L_A=2$	$R=4.0$: $Q=4$, $L=16$, $L_{DM}=64$, $u=[15, 63]$, $L_A=4$
$R=4.5$: $Q=8$, $L=16$, $L_{DM}=128$, $u=[1, 81]$, $L_A=4$	$R=5.0$: $Q=8$, $L=32$, $L_{DM}=256$, $u=[1, 161]$, $L_A=4$
$R=5.5$: $Q=8$, $L=32$, $L_{DM}=256$, $u=[1, 161]$, $L_A=8$	$R=6.0$: $Q=8$, $L=64$, $L_{DM}=512$, $u=[319, 511]$, $L_A=8$

(c) DGC-cyclic, SE-APSK, $M = T = 2$

$R=2.5$: $L=16$, $u=[1, 7]$, $L_A=2$	$R=3.0$: $L=32$, $u=[1, 7]$, $L_A=2$	$R=3.5$: $L=64$, $u=[1, 19]$, $L_A=2$
$R=4.0$: $L=64$, $u=[1, 19]$, $L_A=4$	$R=4.5$: $L=128$, $u=[1, 47]$, $L_A=4$	$R=5.0$: $L=256$, $u=[1, 75]$, $L_A=4$
$R=5.5$: $L=256$, $u=[1, 75]$, $L_A=8$	$R=6.0$: $L=512$, $u=[1, 149]$, $L_A=8$	

(d) DSTSK-TAST, $M = T = 4$

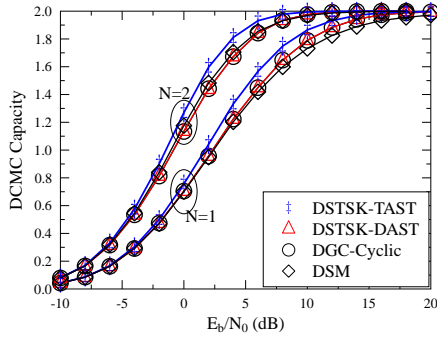
$R=1.5$: $Q=2$, $L=8$, $L_{DM}=16$, $L_r=32$, $u=[1, 5, 11, 15]$	$R=1.75$: $Q=4$, $L=8$, $L_{DM}=32$, $L_r=32$, $u=[5, 25, 11, 31]$
$R=2.0$: $Q=8$, $L=8$, $L_{DM}=64$, $L_r=64$, $u=[15, 41, 57, 63]$	$R=2.25$: $Q=8$, $L=16$, $L_{DM}=128$, $L_r=128$, $u=[31, 89, 113, 127]$
$R=2.5$: $Q=16$, $L=16$, $L_{DM}=256$, $L_r=256$, $u=[1, 41, 55, 135]$	$R=2.75$: $Q=32$, $L=16$, $L_{DM}=512$, $L_r=512$, $u=[71, 41, 489, 503]$
$R=3.0$: $Q=64$, $L=16$, $L_{DM}=1024$, $L_r=1024$, $u=[633, 603, 559, 797]$	

(e) DSTSK-DAST, $M = T = 4$

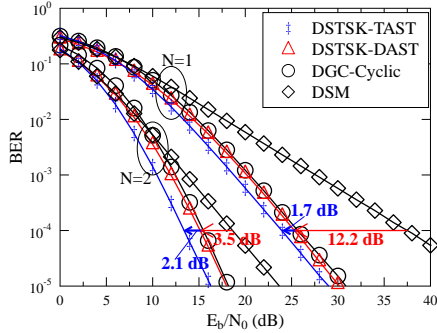
$R=1.5$: $Q=8$, $L=8$, $L_{DM}=64$, $u=[7, 23, 47, 63]$	$R=1.75$: $Q=16$, $L=8$, $L_{DM}=128$, $u=[71, 79, 119, 127]$
$R=2.0$: $Q=32$, $L=8$, $L_{DM}=256$, $u=[1, 41, 49, 137]$	$R=2.25$: $Q=32$, $L=16$, $L_{DM}=512$, $u=[49, 193, 401, 513]$
$R=2.5$: $Q=64$, $L=16$, $L_{DM}=1024$, $u=[5, 999, 1303, 2587]$	$R=2.75$: $Q=64$, $L=32$, $L_{DM}=2048$, $u=[1119, 2047, 2271, 2687]$
$R=3.0$: $Q=128$, $L=32$, $L_{DM}=4096$, $u=[404576931204511501]$	

(f) DGC-cyclic, $M = T = 4$

$R=2.25$: $L=512$, $u=[1, 67, 123, 231]$	$R=2.5$: $L=1024$, $u=[1, 187, 221, 351]$	$R=2.75$: $L=2048$, $u=[1, 277, 325, 919]$
$R=3$: $L=4096$, $u=[1, 493, 695, 1851]$		



(a) DCMC capacity



(b) BER performance

Fig. 9: DCMC capacity and BER comparison of DSTM schemes associated with $M = T = 2$ and $R = 2.0$. The DSTSK-DAST's diversity gains over DSM at $\text{BER}=10^{-4}$ are marked by red arrows, while the DSTSK-TAST's performance improvements over DGC at $\text{BER}=10^{-4}$ are marked by blue arrows.

TABLE XII: Complexity (number of real-valued multiplications) of DSTM schemes.

	Using PSK	Using APSK
DSTSK-TAST	$4NT^2 + 4T^2Q + 5TQ$	$4NT^2 + 4T^2Q + 5TQ + 2N + 3$
DSTSK-DAST	$4NT^2 + 4TQ + 5Q$	$4NT^2 + 4TQ + 5Q + 2N + 3$
DGC-cyclic	$4NT^2 + 5TL$	$4NT^2 + 5TL + 2N + 3$
FE-DSM	$4NT^2 + 4T^2 + 5T$	$4NT^2 + 4T^2 + 5T + 2N + 3$
DSTBC-ISK	$4NT^2 + 2NT + 4T$	$4NT^2 + 2NT + 4T + 2N + 3$
DSM	$4NT^2 + 5Q2^{\lceil \log_2 T \rceil}$	$4NT^2 + 5Q2^{\lceil \log_2 T \rceil} + 2N + 3$
DPSK/APSK	$4N + 1$	$6N + 9$

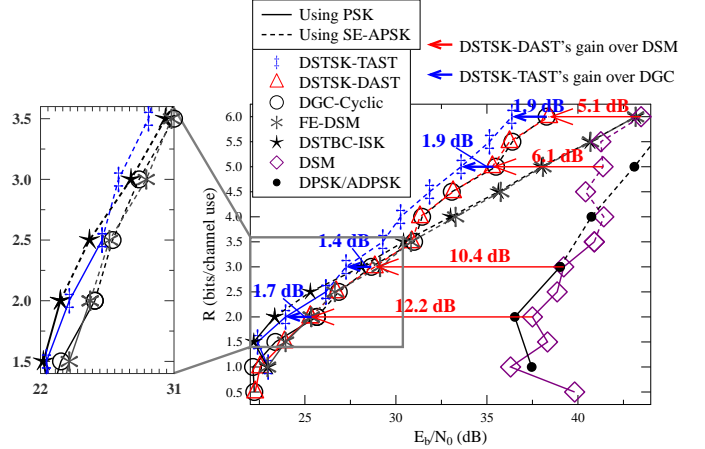
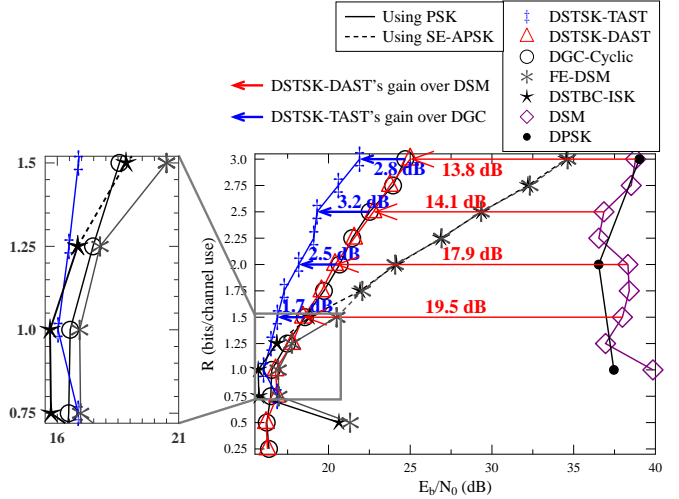
(a) $M = T = 2$ (b) $M = T = 4$

Fig. 10: Comparison of E_b/N_0 required for DSTM schemes associated with $M = T = \{2, 4\}$ and $N = 1$ to achieve $\text{BER}=10^{-4}$. The DSTSK-DAST's diversity gains over DSM at $\text{BER}=10^{-4}$ are marked by red arrows, while the DSTSK-TAST's performance improvements over DGC at $\text{BER}=10^{-4}$ are marked by blue arrows.

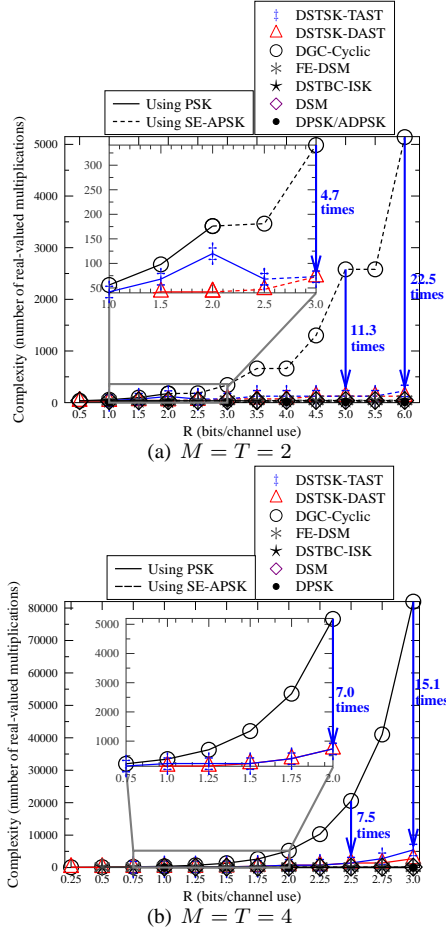


Fig. 11: Complexity comparison for DSTM schemes associated with $M = T = \{2, 4\}$ and $N = 1$, where the DSTSK-TAST's complexity reduction compared to DGC-cyclic are marked by blue arrows.

summarized in Table XII. We also note that the evaluation of the DGC [39], [41]–[43] detection complexity and that of the DSM [31], [32] detection complexity are based on their simplified detectors summarized in [40].

First of all, it is evidenced by Fig. 10 that the DSM does not perform well for a small $N = 1$. Secondly, Fig. 10 demonstrates that DSTBC-ISK outperforms FE-DSM at $1.5 \leq R \leq 4.0$ for $M = T = 2$ and $0.5 \leq R \leq 1.75$ for $M = T = 4$. Moreover, DSTBC-ISK even outperforms DGC-cyclic at the low throughputs of $1.5 \leq R \leq 3.5$ for $M = T = 2$ and $0.75 \leq R \leq 1.25$ for $M = T = 4$, as shown in Fig. 10. However, Fig. 10 also demonstrates that the diversity advantages of both DSTBC-ISK and FE-DSM over DSM diminish, as the throughput increases. By contrast, both the proposed DSTSK-DAST and the classic DGC-cyclic achieve substantial non-diminishing diversity gains over the DSM in Fig. 10, but the DSTSK-DAST's single-stream ML detection complexity is significantly lower than the DGC-cyclic's exponentially increasing complexity portrayed in Fig. 11. Furthermore, the proposed DSTSK-TAST is seen in Fig. 10 to achieve an even better performance than DGC-cyclic, despite its substantially reduced detection complexity seen in Fig. 11.

Specifically, for the case of $M = 2$ and $N = 1$, DSTSK-TAST associated with $R = 3.0$ and $R = 6.0$ achieves **1.4 dB** and **1.9 dB** performance gain over DGC-cyclic in Fig. 10(a) at a **4.7 times** and **22.5 times** lower detection complexity in Fig. 11(a), respectively. Moreover, for the case of $M = 4$ and

$N = 1$, DSTSK-TAST associated with $R = 2.5$ and $R = 3.0$ achieves even more significant gains of **3.2 dB** and **2.8 dB** over DGC-cyclic in Fig. 10(b) at a **7.5 times** and **15.1 times** lower complexity in Fig. 11(b), respectively.

We note that the full-transmit-diversity arrangements are particularly beneficial for the popular Unmanned Aerial Vehicles (UAVs) applications, which are envisioned to be an important part of the IoT. More explicitly, grave airframe-induced shadowing [56] is encountered, when the critical air-ground link becomes blocked by the chassis/fuselage of the UAV. It is reported in [56] that using multiple RAs at the Ground Station (GS) is unable to mitigate this problem, but the transmit diversity design [57] using multiple TAs is capable of significantly improving the connectivity. Therefore, we suggest that the proposed finite-cardinality single-RF DSTSK-DAST and DSTSK-TAST may become suitable candidates for the UAV applications as a benefit of their low complexity and maximized transmit diversity gain, as discussed in [56]–[58] and the references within.

B. Diversity-Throughput Tradeoff

In this section, the ME-APSK design and the DR design are applied to the diversity-oriented DSTM schemes of DSTSK-TAST and DSTBC-ISK for a variety of different MIMO system setups. More explicitly, DSTSK-TAST and DSTBC-ISK are compared to DSM for $M = T = \{2, 4\}$ and $N = \{1, 2, 4, 8, 16\}$ both at low and high throughputs of $R = \{2.0, 5.0\}$ in Figs. 12(a)–(d). First of all, although DSTBC-ISK achieves a better performance for $M = T = 2$ and $N = 1$ at a low $R = 2.0$ in Fig. 12(a), its diversity gain diminishes, as the number of RAs N and/or the throughput R increase, as seen in Figs. 12(a)–(d). By contrast, although the DSTSK-TAST's diversity gain over DSM also decreases upon increasing N , Figs. 12(a)–(d) manifestly demonstrate that the proposed DSTSK-TAST is capable of achieving a persistent performance advantage over DSM for the cases of $M = T = \{2, 4\}$ using both small and large number of RAs $N = \{1, 2, 4, 8, 16\}$ at both low and high throughputs $R = \{2.0, 5.0\}$. Once again, DSTSK-TAST shares the same low transceiver complexity as the DSM, including the finite-cardinality single-RF transmitter and the single-stream ML receiver architectures.

In conclusion, based on our extensive simulations, we suggest that the DSTBC-ISK using the simple STBC signal structure is beneficial for the cases of using a small number of RAs $N = \{1, 2\}$ at low throughputs, such as $R \leq 3.0$ for $M = 2$ and $R \leq 1.5$ for $M = 4$. Furthermore, with the help of the ME-APSK design and the DR design, the DSTSK-TAST is capable of achieving a persistent performance advantage for a variety of different MIMO system setups, such as $M = T = \{2, 4\}$ and $N \in [1, 16]$ at $R \leq 6.0$. Nonetheless, we also note that the construction of the DSTSK-TAST signal matrix becomes extremely challenging at very high throughputs of $R > 6.0$, where the DSM is preferred especially for using $N > 2$.

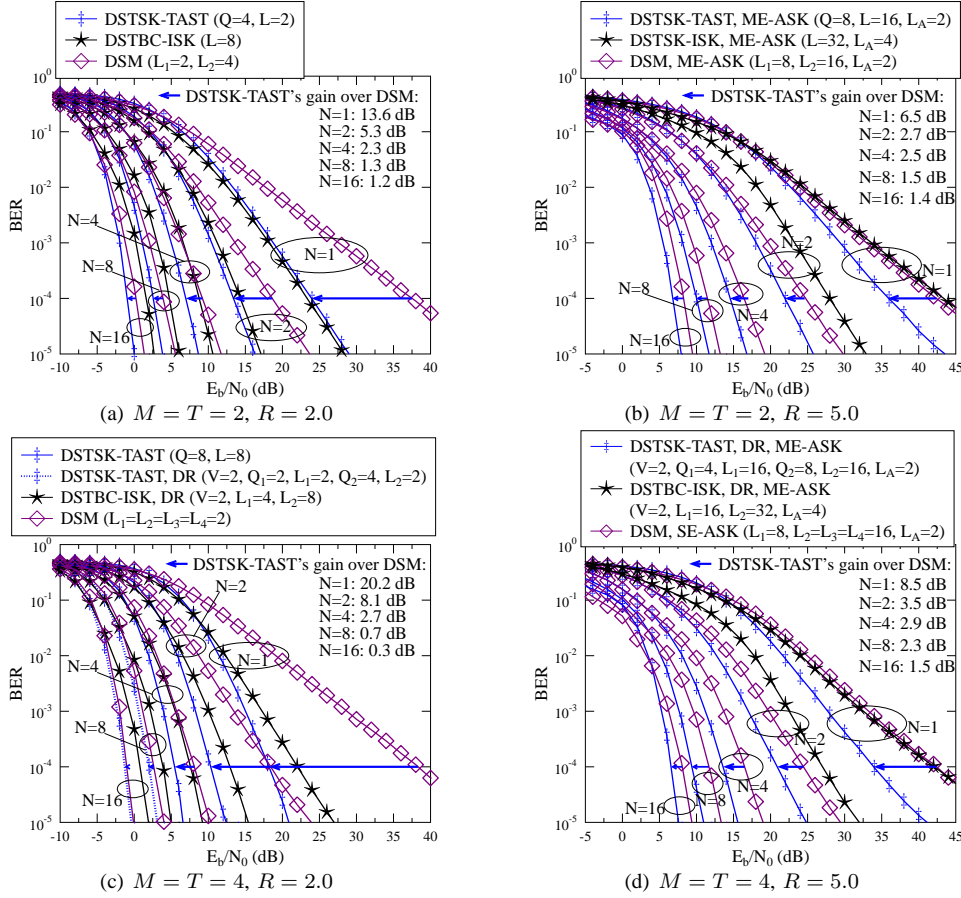


Fig. 12: Performance comparison between the DSTSK-TAST, DSTBC-ISK and DSM arrangements for $M = T = \{2, 4\}$ and $N = \{1, 2, 4, 8, 16\}$ at $R = \{2.0, 5.0\}$. The diversity gains of DSTSK-TAST over DSM at $\text{BER}=10^{-4}$ are marked by blue arrows and also listed in figures.

VII. CONCLUSIONS

In summary, our proposed finite-cardinality single-RF DSTM schemes of DSTSK-DAST and DSTSK-TAST mitigate the problem of the diminishing diversity gain of FE-DSM and DSTBC-ISK over DSM, without compromising its appealingly low transceiver complexity. In order to improve the throughput of the diversity-oriented DSTM schemes, we also conceived the ME-APSK and the DR designs, which assists the proposed DSTSK-TAST in achieving a persistent performance advantage over DSM for a variety of MIMO system setups.

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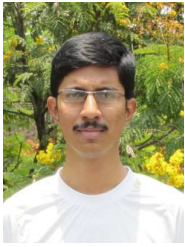
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