

Compressive Sensing Assisted Generalized Quadrature Spatial Modulation for Massive MIMO systems

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Abstract—A novel Multiple-Input and Multiple-Output (MIMO) transmission scheme termed as Generalized Quadrature Spatial Modulation (G-QSM) is proposed. It amalgamates the concept of Quadrature Spatial Modulation (QSM) and spatial multiplexing for the sake of achieving a high throughput, despite relying on low number of Radio Frequency (RF) chains. In the proposed G-QSM scheme, the conventional constellation points of the spatial multiplexing structure are replaced by the QSM symbols, hence the information bits are conveyed both by the antenna indices as well as by the classic Amplitude/Phase Modulated (APM) constellation points. The upper bounds of the Average Bit Error Probability (ABEP) of the proposed G-QSM system in high throughput massive MIMO configurations are derived. Furthermore, an Efficient Multipath Orthogonal Matching Pursuit (EM-OMP) based Compressive Sensing (CS) detector is developed for our proposed G-QSM system. Both our analytical and simulation results demonstrated that the proposed scheme is capable of providing considerable performance gains over the existing schemes in massive MIMO configurations.

Index Terms—Quadrature Spatial Modulation (QSM), Multiple-Input Multiple-Output (MIMO), Vertical Bell Labs Space-Time (VBLAST).

INTRODUCTION

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MASSIVE Multi-input Multi-output (MIMO) [1]–[4] systems are capable of significantly increasing the transmission rate, reliability and system capacity by employing numerous antennas. Hence they may be deemed a key technology in next-generation wireless system. In massive MIMO systems, the classic Vertical Bell Laboratories Space Time (V-BLAST) based MIMO schemes [5]–[6] require numerous Radio Frequency (RF) chains, hence imposing substantial implementation cost and signal processing complexity. As a result, how to design a large-scale MIMO scheme to strike an elegant tradeoff between the throughput and implementation cost has attracted a lot attention.

Recently, Spatial Modulation based variants, which employ the activated antenna indices as an additional means of implicitly conveying information, have been advocated as novel low-cost RF chain based MIMO transmission structures. To be more specific, in SM [7], one out of N_t Transmit Antennas (TAs) is activated to implicitly convey information. Due to the advantages of the single RF chain based MIMO structure, SM scheme has also been studied in massive MIMO communication scenarios [8]–[11]. However, the low number of RF chains limits the transmission rate of the SM scheme. In order to increase its throughput, multiple activated TA based SM variants were proposed in [12]–[19]. Specifically, in the Generalized SM (GSM) scheme [12]–[15], N_u out of N_t TAs are activated to transmit different symbols, so that the transmission rate is significantly increased. In [16]–[17], the number of activated TAs ranged from one to N_t . Additionally, GSM combined with Space Time Block Coding (STBC) structures was designed in [18]–[19]. However, the above SM variants have the common feature that the information bits are conveyed both by the active TA index as well as by the Amplitude Phase Modulation (APM) symbols. In high-throughput massive MIMO scenarios, the number of information bits carried by the index is high, which makes the bit-to-symbol mapping and symbol-to-bit demapping challenging [20].

In order to address the mapping issues in SM based massive MIMO systems, a novel spatial multiplexing based SM scheme is proposed in [20], where the N_t TAs are partitioned into L groups and SM is employed in each group. We refer to this scheme as SM-VBLAST. Compared to the classic VBLAST system, an SM symbol instead

of a single AMP symbol is mapped to each group, so that the throughput is significantly increased at the same number of RF chains. As a result, the SM-VBLAST system circumvents the mapping issues in massive MIMO systems. Moreover, this SM-VBLAST scheme is capable of outperforming the conventional GSM scheme, despite its lower complexity, providing a 6 dB Signal to Noise Ratio (SNR) gain over the V-BLAST system having the same number of RF chains and the same throughput. However, how to fully exploit the advantages of index modulation in massive MIMO systems is still an open question to be investigated.

As an SM variants, Quadrature SM (QSM) [21] also constitutes a high-efficiency MIMO technique, which exploits both the in-phase and quadrature dimensions for enhancing the overall throughput of the conventional SM system. In the QSM scheme, both the indices of the real part as well as of the imaginary part of an M -ary APM symbol and the APM symbol itself convey information, hence resulting in higher throughput than conventional SM. Furthermore, the inter-antenna interference can also be avoided by the QSM scheme, since orthogonal symbols are transmitted by the activated antennas. Recently, the QSM design efforts have been mainly focused on its performance analysis [22]-[24], on its coherent vs non-coherent detection [25]-[27], as well as on practical communication systems [29]-[31]. To the best of our knowledge, QSM based high throughput MIMO schemes have not been investigated. On the other hand, since the SM based MIMO schemes exhibit an inherent transmit symbol sparsity, Compressive Sensing (CS) has been invoked as an efficient technique of detecting the SM based MIMO symbols [32]-[39].

Against the above background, the contributions of this paper are summarized as follows:

- 1) A novel low-complexity high throughput MIMO scheme namely Generalized Quadrature Spatial Modulation (G-QSM) is proposed. The throughput achieved by the G-QSM system is analyzed and the optimal highest-throughput mapping is derived.
- 2) The upper bound of the Average Bit Error Probability (ABEP) of the proposed G-QSM system in high-throughput massive MIMO configurations is derived, where the optimal detector becomes impractical. Hence, the proposed upper bound of ABEP is helpful for performance evaluation in high-throughput communication scenarios.
- 3) Furthermore, an Efficient Multipath Orthogonal Matching Pursuit (EM-OMP) based CS detector is proposed for our G-QSM system. Both our theoretical and simulation results demonstrated that the proposed G-QSM scheme employing a large number of TAs is capable of providing a 20 dB Signal to Noise Ratio (SNR) gain over the V-BLAST system having the same limited number of RF chains and the same throughput.

The remainder of this paper is organized as follows. Section II provides a rudimentary review of the conven-

tional QSM system. In Section III, the system model, the rate achieved and the mapping methods employed in our proposed G-QSM system are introduced. In Section IV, the theoretical analysis of the G-QSM systems having massive MIMO configurations is presented. In Section V, the EM-OMP detector is proposed for our G-QSM system. Section VI portrays our simulation results. Finally, Section VII concludes this paper.

Notation: $\|\cdot\|$ denotes the Frobenious norms of a matrix. $|\cdot|$ represents the magnitude of a complex quantity. $(\cdot)^T$ and $(\cdot)^H$ stand for the transpose and the Hermitian transpose of a vector/matrix, respectively. $A \setminus B$ denotes removing the set B from the set A . C_m^n denotes the binomial coefficient.

QSM SYSTEM

We consider a QSM system having N_t TAs and N_r Receive Antennas (RAs) operating in flat Rayleigh fading channels. In a QSM scheme, the information bits are partitioned in two parts, $2\log_2(N_t)$ bits are used for conveying the active TA indices and $\log_2(M)$ bits mapped to an M -ary APM symbol x . To be specific, $\log_2(N_t)$ bits are assigned to select one TA for transmitting the real part of signal x and another $\log_2(N_t)$ bits are used to select an active TA to transmit the imaginary part of x . Based on the above mapping rule, the transmitted signal $\mathbf{x} \in \mathbb{C}^{N_t \times 1}$ can be represented as

$$\mathbf{x} = \begin{cases} \underbrace{[0, \dots, 0]_{l_{\Re}-1}}_{l_{\Re}-1}, \underbrace{[x_{\Re} + jx_{\Im}, 0, \dots, 0]_{N_t-l_{\Re}}}_{N_t-l_{\Re}}^T, & \text{if } l_{\Re} = l_{\Im}, \\ \underbrace{[0, \dots, 0]_{l_{\Re}-1}}_{l_{\Re}-1}, \underbrace{[x_{\Re}, 0, \dots, 0, jx_{\Im}, 0, \dots, 0]_{N_t-l_{\Im}}}_{N_t-l_{\Im}}^T, & \text{else,} \end{cases} \quad (1)$$

where x_{\Re} and x_{\Im} denote the real and imaginary parts of x , respectively, while $l_{\Re} = (1, \dots, N_t)$ and $l_{\Im} = (1, \dots, N_t)$ represent the corresponding active TA indices.

PROPOSED GENERALIZED QSM SYSTEM

In order to improve the throughput of the QSM systems, a low-complexity yet high-throughput generalized QSM scheme is proposed in this section. Considering a MIMO system having N_t TAs, N_r RAs, the N_t TAs are split into L groups and in each group a single QSM symbol is constructed. Considering $N_t = 4$, $L = 2$ and QPSK for example, the proposed generalized QSM symbol set \mathbb{X}_4^2 is constructed of two QSM symbols $[\mathbf{x}_1, \mathbf{x}_2]$ as seen in (2), where x_1 and x_2 are the constellation symbols.

Generalized System Model of G-QSM

The generalized system model of the proposed G-QSM is shown in Fig. 1. As seen from Fig. 1, the information bits of $R = \sum_{i=1}^L R_i$ are partitioned into L ($1 \leq L \leq N_t/2$) groups and then the QSM mapping principle is employed for each group. Assuming that the i -th ($i = 1, \dots, L$) QSM symbol employs N_Q^i TAs and M_i APM symbols, the i -th block

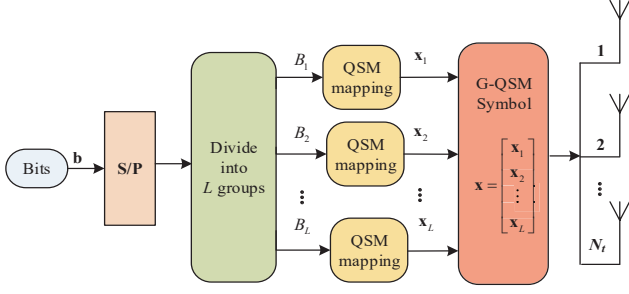


Fig. 1. System model of the proposed G-QSM system

of information bits having a length of $R_i = 2\log_2(N_Q^i) + \log_2(M_i)$ can be mapped to a QSM symbol as

$$\mathbf{x}_i = \begin{cases} [0, \dots, 0, x_{\Re}^i + jx_{\Im}^i, 0, \dots, 0]^T, & \text{if } l_{\Re}^i = l_{\Im}^i, \\ [0, \dots, 0, x_{\Re}^i, 0, \dots, 0, jx_{\Im}^i, 0, \dots, 0]^T, & \text{else,} \end{cases} \quad (3)$$

where x_{\Re}^i and x_{\Im}^i denote the real and imaginary parts of x^i , respectively, while $l_{\Re}^i = (1, \dots, N_Q^i)$ and $l_{\Im}^i = (1, \dots, N_Q^i)$ represent the corresponding active TA indices. Hence, the transmitted signal of the G-QSM system can be expressed as

$$\mathbf{x} = [\mathbf{x}_1^T, \mathbf{x}_2^T, \dots, \mathbf{x}_L^T]^T. \quad (4)$$

Let $\mathbf{H} \in \mathbb{C}^{N_r \times N_t}$ and $\mathbf{n} \in \mathbb{C}^{N_r \times 1}$ be the MIMO channel matrix and noise vector, whose entries are complex-valued Gaussian distributed, yielding $\mathcal{CN}(0, 1)$ and $\mathcal{CN}(0, \sigma^2)$, respectively. The received signal $\mathbf{y} \in \mathbb{C}^{N_r \times 1}$ is written as

$$\mathbf{y} = \mathbf{H}\mathbf{x} + \mathbf{n}. \quad (5)$$

It follows from Eq. (4) that the optimal ML-based demodulator can be formulated as

$$\hat{\mathbf{x}}_{\text{ML}} = \arg \min_{\mathbf{x} \in \mathbb{X}} \|\mathbf{y} - \mathbf{H}\mathbf{x}\|^2, \quad (6)$$

where \mathbb{X} is G-QSM symbol set associated with size of 2^R .

Achievable Rate of the G-QSM system

In this section, the bits/channel use (bpcu) throughput of the G-QSM system is analyzed. According to the system model of the G-QSM system, the transmit rate of the G-QSM system is formulated as

$$R = \sum_{i=1}^L [2\log_2(N_Q^i) + \log_2(M_i)]. \quad (7)$$

where N_Q^i is a power of 2^1 . The transmission rates of the SM-VBLAST and VBLAST are given as follows

$$R_{\text{SM-VBLAST}} = \sum_{i=1}^L [\log_2(N_Q^i) + \log_2(M_i)], \quad (8)$$

$$R_{\text{VBLAST}} = \sum_{i=1}^L \log_2(M_i).$$

For a fixed value of L and $M_i = M$, the maximum transmission rate can be optimized as follows

$$R_{\max} = \max_{N_Q^i} \left[\sum_{i=1}^L 2\log_2(N_Q^i) + \log_2(M)L \right] \quad (9)$$

$$s.t. \sum_{i=1}^L N_Q^i = N_t, N_Q^i = 2^{(R_i - \log_2(M))/2}.$$

Then the optimization problem can be analyzed for the following two cases

- 1) If N_t is a power of 2 with $N_t = 2^c$, the optimization problem is represented as

$$R_{\max} = \max_{B_i} 2 \sum_{i=1}^L B_i + L \log_2(M), \quad (10)$$

$$s.t. \sum_{i=1}^L 2^{B_i} = 2^c, B_i = (1, 2, \dots, c-1),$$

where $B_i = \log_2(N_Q^i)$.

- 2) If N_t is not a power of 2 with $2^c < N_t < 2^{c+1}$, the optimization problem is represented as

$$R_{\max} = \max_{B_i} 2 \sum_{i=1}^L (B_i) + L \log_2(M), \quad (11)$$

$$s.t. \sum_{i=1}^L 2^{B_i} = N_t, B_i = (1, 2, \dots, c).$$

To acquire the values N_Q^1, \dots, N_Q^L , $L-1$ decompositions are needed for N_t TAs. For the l -th $1 \leq l \leq L-1$ decomposition, we have $N_Q^1 = 2^{c_1}$, $N_Q^2 = 2^{c_2}, \dots, N_Q^L = 2^{c_L}$ $c_1 \leq c_2 \leq \dots \leq c_L$ and the transmit rate becomes $R_l = \sum_{i=1}^l 2c_i + \log_2(M)l$. In the $(l+1)$ -th decomposition, we choose any N_Q^i to be divided into the values of $N_Q^i/2, N_Q^i/2$, so the transmit rate becomes $R_{l+1} = 2(c_1 + \dots + c_{l-1} + c_l - 1 + c_l - 1 + c_{l+1} + \dots + c_L) + \log_2(M)l + \log_2(M)$. Accordingly, we have

$$R_{l+1} = R_l + 2(c_l - 2) + \log_2(M). \quad (12)$$

Hence, the different selection of c_i results in different transmission rates.

In [20], we introduce the so-called Large Number First (LNF) principle to obtain N_Q^1, \dots, N_Q^L . For the LNF principle, in the l -th $l \in (2, \dots, L-1)$ step, the largest number N_{l-1}^{lar} in the $(l-1)$ -st set \mathcal{N}_{l-1} is decomposed into two identical numbers as $(N_{l-1}^{lar}/2, N_{l-1}^{lar}/2)$, where $\mathcal{N}_1 = [N_t]$.

¹For simplicity, in this paper, we have focused our attention on the achievable transmit rate of the proposed G-QSM system associated with N_Q^i being a power of 2. When N_Q^i is not a power of 2, the mapping scheme of [19] designed for an arbitrary number of antennas can be employed and the achievable transmit rate analysis following a similar process.

As a result, the decomposed set in the l -th step can be expressed as $N_l = (N_{l-1} \setminus N_{l-1}^{lar}, N_{l-1}^{lar}/2, N_{l-1}^{lar}/2)$. These processes continue, until we get L values as N_Q^1, \dots, N_Q^L . Next, we will demonstrate by the mathematical techniques of induction that the LNF principle invoked for both cases of $N_t = 2^c$ and $N_t \neq 2^c$ can always achieve the maximum transmission rate of (9).

In the $k = 1$ -st decomposition associated with N_t TAs, there are two cases, namely

- 1) If $N_t = 2^c$, we have $L = 2$ and $N_Q^1 = N_Q^2 = 2^{c-1}$, so that we have $B_1 = B_2 = c - 1$ in order to satisfy (10);
- 2) If $2^c < N_t < 2^{c+1}$ with $N_t = \sum_{i=1}^m 2^{c_i}$, we have $L = m$ and $N_Q^1 = 2^{c_1}, N_Q^2 = 2^{c_2}, \dots, N_Q^m = 2^{c_m}$, so that we have $B_1 = c_1, B_2 = c_2, \dots, B_m = c_m$ in order to satisfy (11).

Hence, the maximum transmission rate $R_{k=1}$ of the ($k = 1$)-st decomposition is expressed as

$$R_{k=1} = \begin{cases} 2(2c - 2) + 2\log_2(M), & \text{if } N_t = 2^c, \\ \sum_{i=1}^m 2c_i + m\log_2(M), & \text{else.} \end{cases} \quad (13)$$

- 1) If $k = 2$, we have $L = 3$ or $L = m + 1$. According to (12) and the LNF principle, we have
 - a) If $N_t = 2^c$, we have $L = 3$ and $N_Q^1 = N_Q^2 = 2^{c-2}, N_Q^3 = 2^{c-1}$ with $B_1 = B_2 = c - 2, B_3 = c - 1$;
 - b) If $2^c < N_t < 2^{c+1}$ with $N_t = \sum_{i=1}^m 2^{c_i}$, $c_1 \leq c_2 \leq \dots \leq c_m$, we have $L = m + 1$ and $N_Q^1 = 2^{c_1}, N_Q^2 = 2^{c_2}, \dots, N_Q^m = N_Q^{m+1} = 2^{c_m-1}$ with $B_1 = c_1, B_2 = c_2, \dots, B_m = B_{m+1} = c_m - 1$.

Hence, according to (12)-(13), $R_{k=2}$ can be simplified to

$$R_{k=2}^{LNF} = \begin{cases} R_{k=1} + 2(c - 1 - 2) + \log_2(M), & \text{if } N_t = 2^c, \\ R_{k=1} + 2(c_m - 2) + \log_2(M), & \text{else.} \end{cases} \quad (14)$$

It is evident that the LNF principle is capable of achieving the maximum transmission rate R_{\max} .

- 2) Assuming that the ($k = K$)-th decomposition associated with c_1, c_2, \dots, c_K has the maximum transmit rate R_{\max}^K , according to (12), the maximum rate of the ($K + 1$)-st decomposition is expressed as

$$R_{\max}^{K+1} = R_{\max}^K + \max_i 2(c_i - 2) + \log_2(M). \quad (15)$$

It is evident that the LNF principle allows (15) to achieve the maximum transmission rate of (9).

- 3) Based on the analysis of 1) and 2), it is concluded that the LNF principle allows (9) to remain true for any value of N_t in conjunction with L ($1 \leq L \leq N_t/2$).

In order to provide further insights, Figs.2-3 present the achievable rate of the QPSK aided SM-VBLAST and of

²In this case, $L = m$ is the smallest number that the decomposition can have. Taking $N_t = 12$ and $N_t = 14$ for example, we have $L = 2$ and $L = 3$, respectively.

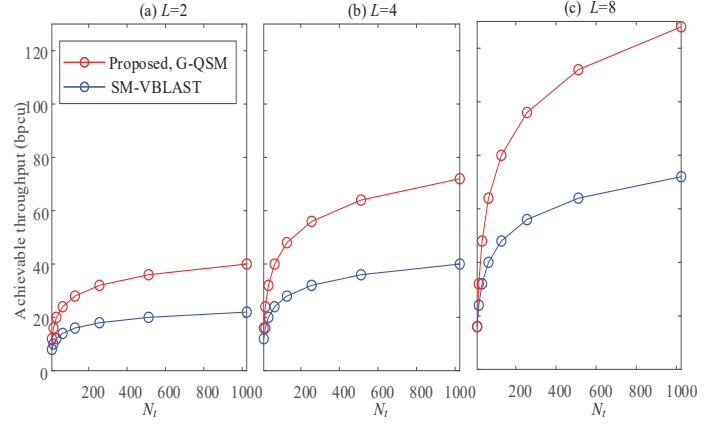


Fig. 2. The achievable rate of the QPSK aided SM-VBLAST and of the proposed G-QSM system having diverse values of N_t and a fixed value of L . a) $L = 2$; b) $L = 4$; c) $L = 8$.

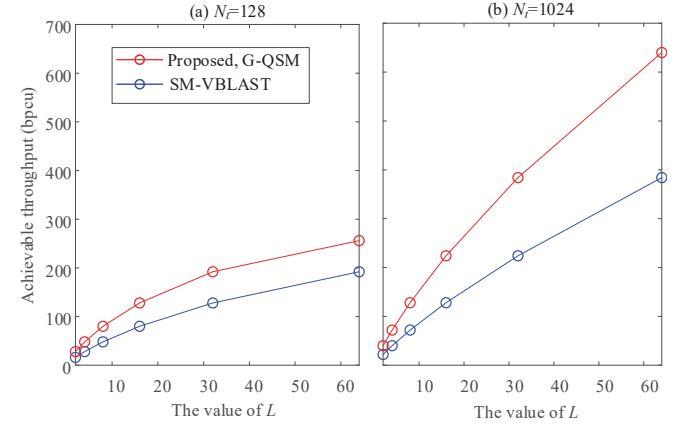


Fig. 3. The achievable rate of the QPSK aided SM-VBLAST and of the proposed G-QSM system having a fixed number of TAs and different values of L . a) $N_t = 128$; b) $N_t = 1024$.

the proposed G-QSM system for different values of L and N_t . Observe from Fig. 2 that the rate achieved by the G-QSM system is significantly higher than that of the SM-VBLAST system for a fixed value of L . The higher the value of L is, the higher the rate advantage becomes. More specifically, for a small value of $L = 8$, the transmission rate of the proposed G-QSM system may reach 128 bpcu, while that of the SM-VBLAST system is 72 bpcu. Observe from Fig. 3 that the rate achieved by the G-QSM system is significantly increased as L increases. The transmission rate of the G-QSM system for $N_t = 1024, L = 64$ becomes 640 bpcu, while that of the SM-VBLAST system is 384 bpcu. Hence, the proposed G-QSM system achieves a high throughput, despite using a low number of RF chains.

Bit-to-Symbol Mapping and Symbol-to-Bit Demapping

In this section, the bit-to-symbol mapping and symbol-to-bit demapping are introduced for our G-QSM system.

Symbol-to-Bit Mapping

The symbol-to-bit mapping of the proposed G-QSM system is introduced in detail as follows.

Step 1: Obtain the numbers $(N_Q^1, N_Q^2, \dots, N_Q^L)$ for constructing the QSM symbols based on the LNF principle [26].

Step 2: Construct the i -th QSM symbol \mathbf{x}_i using N_Q^i TAs and obtain the indices $(q_{\mathcal{R}}^i, q_{\mathcal{I}}^i)$ of the real and imaginary parts of x_i .

Step 3: Obtain the whole G-QSM symbol \mathbf{x} by (4) based on the L obtained QSM symbols of Step 2. Assuming that $(l_{\mathcal{R}}^i, l_{\mathcal{I}}^i)$ are the indices of the real and imaginary parts of x_i in the \mathbf{x} , the relationship between $(l_{\mathcal{R}}^i, l_{\mathcal{I}}^i)$ and $(q_{\mathcal{R}}^i, q_{\mathcal{I}}^i)$ is expressed as

$$l_{\mathcal{R}}^i = \sum_{t=1}^{i-1} N_Q^t + q_{\mathcal{R}}^i; l_{\mathcal{I}}^i = \sum_{t=1}^{i-1} N_Q^t + q_{\mathcal{I}}^i. \quad (16)$$

Symbol-to-Bit Demapping

At the receiver, the symbol-to-bit de-mapping processes rely on the values of (N_Q^1, \dots, N_Q^L) and the values of $(l_{\mathcal{R}}^i, l_{\mathcal{I}}^i, x_i)$ ($i = 1, \dots, L$), which operate in detail as follows.

Step 1: Obtain the values of $(q_{\mathcal{R}}^i, q_{\mathcal{I}}^i, x_i)$ ($i = 1, \dots, L$) based on the detector and (15).

Step 2: Demodulate $(q_{\mathcal{R}}^i, q_{\mathcal{I}}^i)$ into the information bits according to $\mathbf{b}_{\mathcal{R}}^i = \text{de2bi}[q_{\mathcal{R}}^i, \log_2(N_Q^i)]$ and $\mathbf{b}_{\mathcal{I}}^i = \text{de2bi}[q_{\mathcal{I}}^i, \log_2(N_Q^i)]$, where $\text{de2bi}(m, n)$ is a function that transforms a decimal number m to n bits. Demodulation x_i into information bits generates \mathbf{b}_x^i . Hence the i -th QSM symbol is demodulated into bits as $\mathbf{b}_i = [\mathbf{b}_{\mathcal{R}}^i, \mathbf{b}_{\mathcal{I}}^i, \mathbf{b}_x^i]$.

Step 3: Finally, we obtain the demodulated bits as $\mathbf{b} = [\mathbf{b}_1, \dots, \mathbf{b}_i, \dots, \mathbf{b}_L]$.

PERFORMANCE ANALYSIS FOR THE PROPOSED GENERALIZED QSM SYSTEM

The ABEP of the Proposed G-QSM System

Let us denote the transmit signal and receive signal of our G-QSM system by $\mathbf{x}^i = [\mathbf{x}_1^i, \dots, \mathbf{x}_L^i]$ and $\mathbf{x}^j = [\mathbf{x}_1^j, \dots, \mathbf{x}_L^j]$, respectively. The ABEP upper bound is given by

$$P_b = \frac{1}{R2^R} \sum_{i=1}^{2^R} \sum_{j=1, j \neq i}^{2^R} d(\mathbf{x}^i, \mathbf{x}^j) P(\mathbf{x}^i \rightarrow \mathbf{x}^j), \quad (17)$$

where $P(\mathbf{x}^i \rightarrow \mathbf{x}^j)$ denotes the Pairwise Error Probability (PEP), $d(\mathbf{x}^i, \mathbf{x}^j)$ is the Hamming Distance (HD) associated with the corresponding PEP event. According to [20], $P(\mathbf{x}^i \rightarrow \mathbf{x}^j)$ is expressed as

$$P(\mathbf{x}^i \rightarrow \mathbf{x}^j) = F(\bar{\zeta}) = \gamma(\bar{\zeta})^{N_r} \sum_{k=0}^{N_r-1} \binom{N_r-1+k}{k} [1-\gamma(\bar{\zeta})]^k, \quad (18)$$

where $\gamma(\bar{\zeta}) = \frac{1}{2} \left(1 - \sqrt{\frac{\bar{\zeta}/2}{1+\bar{\zeta}/2}} \right)$ and $\bar{\zeta}$ is the mean value of $\zeta = \|\mathbf{H}(\mathbf{x}^i - \mathbf{x}^j)\|^2 / 2\sigma^2$ with $N_r = 1$ as

$$\bar{\zeta} = E\{\|\mathbf{H}(\mathbf{x}^i - \mathbf{x}^j)\|^2 / 2\sigma^2 | N_r = 1\} = \|\mathbf{x}^i - \mathbf{x}^j\|^2 / 2\sigma^2. \quad (19)$$

According to (17), the ABEP is obtained by calculating 2^{2R} PEP events and HDs, which becomes impractical in scenarios having a high throughput. In order to evaluate the large-scale G-QSM's performance, simplified ABEP

upper bound expressions are derived. Specifically, for QPSK modulation, Eq. (17) can be further simplified to

$$P_b = \frac{1}{R} \sum_{j=2}^{2^R} d(\mathbf{x}^1, \mathbf{x}^j) P(\mathbf{x}^1 \rightarrow \mathbf{x}^j). \quad (20)$$

For 16-QAM modulation, (17) cannot be represented by (20), hence the simplified ABEP upper bound calculation becomes again more complicated. Hence, the simplified ABEP upper bound expressions analysis are only valid for our proposed G-QSM system using QPSK modulation, whilst other high-order schemes will be consider in our future work. Specifically, according to [20], there are only n different values $\bar{\zeta}$ as $\bar{\zeta}_1, \dots, \bar{\zeta}_n$ in (18), Eq. (20) can be finally simplified to

$$P_b = \frac{\sum_{t=1}^{\varphi_1} d_{\bar{\zeta}_1}^t F(\bar{\zeta}_1) + \sum_{t=1}^{\varphi_2} d_{\bar{\zeta}_2}^t F(\bar{\zeta}_2) + \dots + \sum_{t=1}^{\varphi_n} d_{\bar{\zeta}_n}^t F(\bar{\zeta}_n)}{R}, \quad (21)$$

$$= D(\bar{\zeta}_1)F(\bar{\zeta}_1) + D(\bar{\zeta}_2)F(\bar{\zeta}_2) + \dots + D(\bar{\zeta}_n)F(\bar{\zeta}_n)$$

with

$$D(\bar{\zeta}_p) = \frac{\sum_{t=1}^{\varphi_p} d_{\bar{\zeta}_p}^t}{R}, p = 1, \dots, n, \quad (22)$$

$$\sum_{t=1}^{\varphi_1} d_{\bar{\zeta}_1}^t + \dots + \sum_{t=1}^{\varphi_n} d_{\bar{\zeta}_n}^t = \sum_{j=1}^{2^R} d(\mathbf{x}^1, \mathbf{x}^j) = \sum_{r=1}^R C_R^r r = R2^{R-1},$$

where φ_p is the number of $\bar{\zeta}_p$ for $(\mathbf{x}^1, \mathbf{x}^j)$ and $\sum_{r=1}^R C_R^r r = R2^{R-1}$, which can be seen in the Appendix. As a result, the calculation of ABEP only has to find the n different values of $\bar{\zeta}_1, \dots, \bar{\zeta}_n$ and to compute their corresponding values of $F(\bar{\zeta}_1), \dots, F(\bar{\zeta}_n)$ and $D(\bar{\zeta}_1), \dots, D(\bar{\zeta}_n)$. Furthermore, $F(\bar{\zeta}_1), \dots, F(\bar{\zeta}_n)$ can be obtained by (17), so that the calculation of $D(\bar{\zeta}_1), \dots, D(\bar{\zeta}_n)$ is the key issue, which is introduced in detail as follows.

ABEP of G-QSM with $L = 1$

The system model of the G-QSM system with $L = 1$ is the same as that of the QSM system. In this section, a simplified expression of ABEP upper bound is derived for large-scale QSM system. Assuming that $l_{\mathcal{R}}^i, l_{\mathcal{I}}^i, x_{\mathcal{R}}^i, x_{\mathcal{I}}^i$ are the antenna indices and symbols of the QSM symbol \mathbf{x}^i , the Euclidean Distance (ED) between \mathbf{x}^i and \mathbf{x}^j is given by:

$$\|\mathbf{x}^i - \mathbf{x}^j\|^2 = \begin{cases} |x_{\mathcal{R}}^i - x_{\mathcal{R}}^j|^2 + |x_{\mathcal{I}}^i - x_{\mathcal{I}}^j|^2, & \text{if } l_{\mathcal{R}}^i = l_{\mathcal{R}}^j, l_{\mathcal{I}}^i = l_{\mathcal{I}}^j; \\ |x_{\mathcal{R}}^i - x_{\mathcal{R}}^j|^2 + |x_{\mathcal{I}}^i|^2 + |x_{\mathcal{I}}^j|^2, & \text{if } l_{\mathcal{R}}^i = l_{\mathcal{R}}^j, l_{\mathcal{I}}^i \neq l_{\mathcal{I}}^j; \\ |x_{\mathcal{R}}^i|^2 + |x_{\mathcal{R}}^j|^2 + |x_{\mathcal{I}}^i - x_{\mathcal{I}}^j|^2, & \text{if } l_{\mathcal{R}}^i \neq l_{\mathcal{R}}^j, l_{\mathcal{I}}^i = l_{\mathcal{I}}^j; \\ |x_{\mathcal{R}}^i|^2 + |x_{\mathcal{I}}^i|^2 + |x_{\mathcal{R}}^j|^2 + |x_{\mathcal{I}}^j|^2, & \text{if } l_{\mathcal{R}}^i \neq l_{\mathcal{R}}^j, l_{\mathcal{I}}^i \neq l_{\mathcal{I}}^j. \end{cases} \quad (23)$$

For the case of QPSK modulation, there are only 4 different values of $\bar{\zeta}_1 = 1/2\sigma^2$, $\bar{\zeta}_2 = 2/2\sigma^2$, $\bar{\zeta}_3 = 3/2\sigma^2$ and $\bar{\zeta}_4 = 4/2\sigma^2$ for each different i in 2^R EDs $\|\mathbf{x}^i - \mathbf{x}^j\|^2$ $j = (1, \dots, 2^R)$, so that (17) can be represented to (20). For the case of 16-QAM,

2^R EDs $\|\mathbf{x}^i - \mathbf{x}^j\|^2$ $j = (1, \dots, 2^R)$ are different for different values of i . Specifically, we have $\|\mathbf{x}_i - \mathbf{x}_j\|^2 \in \{0.2, 0.4, 0.6, 0.8, 1, 1.2, 1.4, 1.6, 1.8, 2, 2.2, 2.6, 2.8, 3.2, 3.4, 3.6, 3.8, 4, 4.6, 5.2, 5.4, 7.2\}$ for 16-QAM. For the transmit signal \mathbf{x}_1 , we have $\|\mathbf{x}_1 - \mathbf{x}_j\|^2 \in \{0.4, 0.8, 1, 1.4, 1.6, 1.8, 2, 2.2, 2.6, 2.8, 3.2, 3.4, 3.6, 4, 4.6, 5.2, 5.4, 7.2\}$, while $\|\mathbf{x}_2 - \mathbf{x}_j\|^2 \in \{0.2, 0.4, 0.6, 0.8, 1, 1.2, 1.4, 1.6, 1.8, 2, 2.2, 2.6, 2.8, 3.2, 3.4, 3.6, 3.8, 4, 4.6, 5.2\}$ for $(j = 1, \dots, 2^R)$ holds for \mathbf{x}_2 , hence (17) cannot be represented by (20). For the large-scale QPSK aided QSM ABEP calculation, we only have to compute four corresponding values of $F(\bar{\zeta}_1), \dots, F(\bar{\zeta}_4)$ and $D(\bar{\zeta}_1), \dots, D(\bar{\zeta}_4)$.

Specifically, the information bits of the G-QSM system with $L = 1$ conveyed consists of four parts: $\mathbf{b}_\mathbb{R}$, $\mathbf{b}_\mathbb{I}$, $\mathbf{c}_\mathbb{R}$ and $\mathbf{c}_\mathbb{I}$, where $\mathbf{b}_\mathbb{R}$ is the bit vector that is conveyed by the real symbol index, $\mathbf{b}_\mathbb{I}$ is the bit vector conveyed by the imaginary symbol index, $\mathbf{c}_\mathbb{R}$ is the bit vector conveyed by the real symbol, and $\mathbf{c}_\mathbb{I}$ is the bit vector conveyed by the imaginary symbol. Due to the symmetry, we set $i = 1$ for simplicity. The information bit vectors that the \mathbf{x}^1 and \mathbf{x}^j symbols conveyed can be expressed as

$$\begin{aligned} \mathbf{x}^1 &\rightarrow [\underbrace{0, \dots, 0}_{\mathbf{b}_\mathbb{R}^1}, \underbrace{0, \dots, 0}_{\mathbf{b}_\mathbb{I}^1}, \underbrace{0}_{\mathbf{c}_\mathbb{R}^1}, \underbrace{0}_{\mathbf{c}_\mathbb{I}^1}], \\ \mathbf{x}^j &\rightarrow [\underbrace{b_{r,1}, \dots, b_{r,B}}_{\mathbf{b}_\mathbb{R}^j}, \underbrace{b_{j,1}, \dots, b_{j,B}}_{\mathbf{b}_\mathbb{I}^j}, \underbrace{c_1}_{\mathbf{c}_\mathbb{R}^j}, \underbrace{c_2}_{\mathbf{c}_\mathbb{I}^j}]. \end{aligned} \quad (24)$$

Then the HDs between $(\mathbf{b}_\mathbb{R}^1, \mathbf{b}_\mathbb{R}^j)$, $(\mathbf{b}_\mathbb{I}^1, \mathbf{b}_\mathbb{I}^j)$, $(\mathbf{c}_\mathbb{R}^1, \mathbf{c}_\mathbb{R}^j)$ and $(\mathbf{c}_\mathbb{I}^1, \mathbf{c}_\mathbb{I}^j)$ are defined as

$$\begin{aligned} d_\zeta^r &= \text{Diff}(\mathbf{b}_\mathbb{R}^1, \mathbf{b}_\mathbb{R}^j) \in \mathbb{B} = [0, \underbrace{1, \dots, 1}_{C_B^1}, \dots, \underbrace{b, \dots, b}_{C_B^b}, \dots, \underbrace{B}_{C_B^B}], \\ d_\zeta^j &= \text{Diff}(\mathbf{b}_\mathbb{I}^1, \mathbf{b}_\mathbb{I}^j) \in \mathbb{B} = [0, \underbrace{1, \dots, 1}_{C_B^1}, \dots, \underbrace{b, \dots, b}_{C_B^b}, \dots, \underbrace{B}_{C_B^B}], \\ d_\zeta^i &= \text{Diff}([c_\mathbb{R}^1, c_\mathbb{R}^j] \in \mathbb{P}, d_\zeta^q = \text{Diff}([c_\mathbb{R}^1, c_\mathbb{R}^j] \in \mathbb{P}), \end{aligned} \quad (25)$$

where $B = \log_2(N_t)$, $b \in (0, B)$, $\text{Diff}(x, y)$ is a function returning the difference between x and y , \mathbb{B} is the HD set of the bits that antenna indices conveyed, and \mathbb{P} is the HD set of the bits that the AMP symbols carried and we have $\mathbb{P} = [0, 1]$ for QPSK. Then the Average HD (AHD) between \mathbf{x}^1 and \mathbf{x}^j can be represented as

$$D(\bar{\zeta}) = \frac{\sum d_{\bar{\zeta}}}{R} = \frac{\sum_{i=1}^2 \sum_{q=1}^2 \sum_{j=1}^{2^B} \sum_{r=1}^{2^B} (d_{\bar{\zeta}}^i + d_{\bar{\zeta}}^q + d_{\bar{\zeta}}^j + d_{\bar{\zeta}}^r)}{2B + 2}. \quad (26)$$

According to (23)-(24), the values of $D(\bar{\zeta}_1), \dots, D(\bar{\zeta}_4)$ are counted as follows.

Case 1 $\bar{\zeta}_1 = 1/2\sigma^2$: According to (21), there are two scenarios $l_\mathbb{R}^1 \neq l_\mathbb{R}^j, l_\mathbb{I}^1 = l_\mathbb{I}^j, x_\mathbb{R}^1 = x_\mathbb{R}^j$ and $l_\mathbb{I}^1 \neq l_\mathbb{I}^j, l_\mathbb{R}^1 = l_\mathbb{R}^j, x_\mathbb{R}^1 = x_\mathbb{R}^j$ having the result of $\bar{\zeta}_1 = 1/2\sigma^2$. The HDs between \mathbf{x}^1 and \mathbf{x}^j for these two cases are computed as follows.

- 1) If $l_\mathbb{R}^1 \neq l_\mathbb{R}^j, l_\mathbb{I}^1 = l_\mathbb{I}^j, x_\mathbb{R}^1 = x_\mathbb{R}^j$ is satisfied, we have the following HDs as $d_\zeta^i = d_\zeta^j = 0$, $d_\zeta^r \neq 0$, and $d_\zeta^q \in [0, 1]$.

There are a total of $\varphi_{1,1} = 2 \times (2^B - 1)$ scenarios having this result, so that the total HDs between \mathbf{x}^1 and \mathbf{x}^j in these conditions can be expressed as

$$\begin{aligned} H_{1,1} &= \sum_{i=1}^2 \sum_{r=2}^{2^B} (d_\zeta^i + d_\zeta^r) = (2^B - 1) \sum_{i=1}^2 d_\zeta^i + 2 \times \sum_{r=2}^{2^B} d_\zeta^r \\ &= (2^B - 1) + 2 \sum_{b=1}^B C_B^b b. \end{aligned} \quad (27)$$

- 2) If $l_\mathbb{I}^1 \neq l_\mathbb{I}^j, l_\mathbb{R}^1 = l_\mathbb{R}^j, x_\mathbb{R}^1 = x_\mathbb{R}^j$ is satisfied, we have $d_\zeta^r = d_\zeta^i = 0$, $d_\zeta^j \neq 0$ and $d_\zeta^q \in [0, 1]$. There are also a total of $\varphi_{1,2} = 2 \times (2^B - 1)$ scenarios having this result, so that the total HDs between \mathbf{x}^1 and \mathbf{x}^j in these conditions can be expressed as

$$\begin{aligned} H_{1,2} &= \sum_{q=1}^2 \sum_{j=2}^{2^B} (d_\zeta^q + d_\zeta^j) \\ &= (2^B - 1) \sum_{q=1}^2 d_\zeta^q + 2 \times \sum_{j=2}^{2^B} d_\zeta^j = (2^B - 1) + 2 \sum_{b=1}^B C_B^b b. \end{aligned} \quad (28)$$

Hence, the AHD $D(\bar{\zeta}_1)$ is computed by

$$D(\bar{\zeta}_1 = \frac{1}{2\sigma^2}) = \frac{H_1}{R} = \frac{H_{1,1} + H_{1,2}}{2B + 2} = \frac{2(2^B - 1) + 4 \sum_{b=1}^B C_B^b b}{2B + 2}, \quad (29)$$

and the total number of conditions having $\bar{\zeta}_1 = 1/2\sigma^2$ is expressed as $\varphi_1 = \varphi_{1,1} + \varphi_{1,2} = 4(2^B - 1)$.

Case 2 $\bar{\zeta}_2 = 2/2\sigma^2$: According to (21), there are three conditions $(l_\mathbb{R}^1 \neq l_\mathbb{R}^j, l_\mathbb{I}^1 \neq l_\mathbb{I}^j)$, $(l_\mathbb{I}^1 = l_\mathbb{I}^j, l_\mathbb{R}^1 = l_\mathbb{R}^j, x_\mathbb{R}^1 \neq x_\mathbb{R}^j)$, $(l_\mathbb{I}^1 = l_\mathbb{I}^j, l_\mathbb{R}^1 = l_\mathbb{R}^j, x_\mathbb{I}^1 \neq x_\mathbb{I}^j)$ having the result of $\bar{\zeta}_2 = 2/2\sigma^2$. The total HDs for these three cases are computed as follows.

- 1) If $l_\mathbb{R}^1 \neq l_\mathbb{R}^j, l_\mathbb{I}^1 \neq l_\mathbb{I}^j$ is satisfied, we have $d_\zeta^i \neq 0$, $d_\zeta^r \neq 0$, $d_\zeta^j \in [0, 1]$, $d_\zeta^q \in [0, 1]$. There are a total of $\varphi_{2,1} = 4(2^B - 1)^2$ scenarios having this result, so that the total HDs between \mathbf{x}^1 and \mathbf{x}^j in these conditions can be expressed as

$$\begin{aligned} H_{2,1} &= \sum_{i=1}^2 \sum_{q=1}^2 \sum_{j=2}^{2^B} \sum_{r=2}^{2^B} (d_\zeta^i + d_\zeta^q + d_\zeta^j + d_\zeta^r) \\ &= 2w^2 \sum_{i=1}^2 d_\zeta^i + 2w^2 \sum_{i=1}^2 d_\zeta^q + 4w \sum_{i=1}^2 d_\zeta^j + 4w \sum_{i=1}^2 d_\zeta^r \\ &= 2w^2 + 2w^2 + 8w \sum_{b=1}^B C_B^b b, \end{aligned} \quad (30)$$

where $w = 2^B - 1$.

- 2) If $l_\mathbb{I}^1 = l_\mathbb{I}^j, l_\mathbb{R}^1 = l_\mathbb{R}^j, x_\mathbb{R}^1 = x_\mathbb{R}^j, x_\mathbb{I}^1 \neq x_\mathbb{I}^j$ is satisfied, we have $d_\zeta^i = d_\zeta^j = d_\zeta^r = 0$, $d_\zeta^q = 1$. There is only one $\varphi_{2,2} = 1$ scenario having this result, hence the HD for this case can be expressed as

$$H_{2,2} = d_\zeta^i + d_\zeta^j + d_\zeta^r + d_\zeta^q = 1. \quad (31)$$

- 3) If $l_\mathbb{I}^1 = l_\mathbb{I}^j, l_\mathbb{R}^1 = l_\mathbb{R}^j, x_\mathbb{I}^1 = x_\mathbb{I}^j, x_\mathbb{R}^1 \neq x_\mathbb{R}^j$ is satisfied, we have $d_\zeta^i = d_\zeta^j = d_\zeta^r = 0$, $d_\zeta^q = 1$. There is only one $\varphi_{2,3} = 1$ scenario having this result, hence the HD

for this case can be formulated as

$$H_{2,3} = d_{\bar{\zeta}}^i + d_{\bar{\zeta}}^j + d_{\bar{\zeta}}^r + d_{\bar{\zeta}}^q = 1. \quad (32)$$

Hence, the AHD $D(\bar{\zeta}_2)$ between \mathbf{x}^1 and \mathbf{x}^j having $\bar{\zeta}_2 = 2/2\sigma^2$ is expressed by

$$D(\bar{\zeta}_2 = \frac{2}{2\sigma^2}) = \frac{H_2}{R} = \frac{H_{2,1}+H_{2,2}+H_{2,3}}{2B+2} = \frac{4w^2+8w \sum_{b=1}^B C_B^b b+2}{2B+2}, \quad (33)$$

and the total number of conditions having $\bar{\zeta}_2 = 2/2\sigma^2$ is expressed as $\varphi_2 = 4(2^B - 1)^2 + 2$.

Case 3 $\bar{\zeta}_1 = 3/2\sigma^2$: According to (21), there are two conditions as $l_{\mathbb{R}}^1 \neq l_{\mathbb{R}}^2, l_{\mathbb{S}}^1 = l_{\mathbb{S}}^2, x_{\mathbb{S}}^1 \neq x_{\mathbb{S}}^2$ and $l_{\mathbb{R}}^1 = l_{\mathbb{R}}^2, l_{\mathbb{S}}^1 \neq l_{\mathbb{S}}^2, x_{\mathbb{R}}^1 \neq x_{\mathbb{R}}^2$ having the result of $\bar{\zeta}_3 = 3/2\sigma^2$. The HDs for these two cases are computed as follows.

- 1) If $l_{\mathbb{R}}^1 \neq l_{\mathbb{R}}^2, l_{\mathbb{S}}^1 = l_{\mathbb{S}}^2, x_{\mathbb{S}}^1 \neq x_{\mathbb{S}}^2$ is satisfied, we have $d_{\bar{\zeta}}^r \neq 0, d_{\bar{\zeta}}^j = 0, d_{\bar{\zeta}}^q = 1, d_{\bar{\zeta}}^i \in [0, 1]$. There are a total of $\varphi_{3,1} = 2(2^B - 1)$ scenarios having this result, so that the HDs for these cases can be expressed as

$$H_{3,1} = \sum_{i=1}^2 \sum_{r=2}^{2^B} (d_{\bar{\zeta}}^i + 1 + d_{\bar{\zeta}}^r) = 3(2^B - 1) + 2 \sum_{b=1}^B C_B^b b. \quad (34)$$

- 2) If $l_{\mathbb{R}}^1 = l_{\mathbb{R}}^2, l_{\mathbb{S}}^1 \neq l_{\mathbb{S}}^2, x_{\mathbb{R}}^1 \neq x_{\mathbb{R}}^2$ is satisfied, we have $d_{\bar{\zeta}}^r = 0, d_{\bar{\zeta}}^j \neq 0, d_{\bar{\zeta}}^i = 1$ and $d_{\bar{\zeta}}^q \in [0, 1]$. There are a total of $\varphi_{3,2} = 2(2^B - 1)$ scenarios having this result, so that the HDs for these cases can be expressed as

$$H_{3,2} = \sum_{q=1}^2 \sum_{j=2}^{2^B} (d_{\bar{\zeta}}^q + 1 + d_{\bar{\zeta}}^j) = 3(2^B - 1) + 2 \sum_{b=1}^B C_B^b b. \quad (35)$$

As a result, the AHD $D(\bar{\zeta}_3)$ between \mathbf{x}^1 and \mathbf{x}^j having $\bar{\zeta}_3 = 3/2\sigma^2$ is formulated as

$$D(\bar{\zeta}_3 = \frac{3}{2\sigma^2}) = \frac{H_3}{R} = \frac{H_{3,1}+H_{3,2}}{2B+2} = \frac{2[3(2^B-1)+2 \sum_{b=1}^B C_B^b b]}{2B+2}, \quad (36)$$

and the total number of conditions having $\bar{\zeta}_3 = 3/2\sigma^2$ is expressed as $\varphi_3 = \varphi_{3,1} + \varphi_{3,2} = 4(2^B - 1)$.

Case 4 $\bar{\zeta}_4 = 4/2\sigma^2$: According to (21), there is only one condition $l_{\mathbb{R}}^1 = l_{\mathbb{R}}^2, l_{\mathbb{S}}^1 = l_{\mathbb{S}}^2, x_{\mathbb{R}}^1 \neq x_{\mathbb{R}}^2, x_{\mathbb{S}}^1 \neq x_{\mathbb{S}}^2$ having the result of $\bar{\zeta}_4 = 4/2\sigma^2$. In this case, we have $d_{\bar{\zeta}}^r = 0, d_{\bar{\zeta}}^j = 0, d_{\bar{\zeta}}^i = d_{\bar{\zeta}}^q = 1$, the AHD between \mathbf{x}^1 and \mathbf{x}^j is expressed as

$$D(\bar{\zeta}_4 = \frac{4}{2\sigma^2}) = \frac{H_4}{R} = \frac{2}{2B+2}. \quad (37)$$

Since $\sum_{m=1}^4 \varphi_m = 2^R - 1$, $\sum_{b=1}^B C_B^b b = B2^{B-1}$ and $H_1 + H_2 + H_3 + H_4 = \sum_{u=1}^R C_R^u u$, the ABEP of the G-QSM for $L = 1$ can be obtained by

$$P_b^{L=1} = \left[\frac{(B+1)2^B - 1}{B+1} \right] F\left(\frac{1}{2\sigma^2}\right) + \left[\frac{(B+1)2^{2B+1} - (2B+4)2^B + 3}{B+1} \right] F\left(\frac{1}{\sigma^2}\right) + \left[\frac{(B+3)2^B - 3}{B+1} \right] F\left(\frac{3}{2\sigma^2}\right) + \left[\frac{1}{B+1} \right] F\left(\frac{2}{\sigma^2}\right). \quad (38)$$

It is shown in (38) that the ABEP upper bound is obtained by only calculating as few as 4 PEP and HD events, while that in [21] has to calculate 2^{4B+4} PEP events and HDs, which becomes impractical for large-scale QSM systems. As a result, the expression derived in (38) becomes helpful for large-scale QSM systems performance evaluation.

ABEP of G-QSM with $L = 2$

For the case of $L = 2$, we have $R = \sum_{l=1}^2 2B_l + 4$ with $B_l = \log_2(N_Q^l)$. Assuming that \mathbf{x}^i and \mathbf{x}^j are two different G-QSM symbols expressed as

$$\mathbf{x}^i = \begin{bmatrix} \mathbf{x}_1^i \\ \mathbf{x}_2^i \end{bmatrix}, \mathbf{x}^j = \begin{bmatrix} \mathbf{x}_1^j \\ \mathbf{x}_2^j \end{bmatrix}, \quad (39)$$

the ED between them can be computed as

$$\|\mathbf{x}^i - \mathbf{x}^j\|^2 = \left\| \begin{bmatrix} \mathbf{x}_1^i - \mathbf{x}_1^j \\ \mathbf{x}_2^i - \mathbf{x}_2^j \end{bmatrix} \right\|^2 = \|\mathbf{x}_1^i - \mathbf{x}_1^j\|^2 + \|\mathbf{x}_2^i - \mathbf{x}_2^j\|^2. \quad (40)$$

Assuming that \mathbf{x}_1^i and \mathbf{x}_2^i are mapped by the information bits \mathbf{b}_1^i with length of $2B_1 + 2$ and \mathbf{b}_2^i with length of $2B_2 + 2$, respectively, according to (21), there are five values $\{0, 1, 2, 3, 4\}$ for $\chi_1 = \|\mathbf{x}_1^i - \mathbf{x}_1^j\|^2$ and $\chi_2 = \|\mathbf{x}_2^i - \mathbf{x}_2^j\|^2$. Let us define $d_{\zeta_{n_l}}^l = \text{Diff}(\mathbf{b}_l^i, \mathbf{b}_l^j)$ by the HD associated with $\chi_l = n_l$, $\varphi_{n_l}^l$ by the number of this case. The total HD associated with $\chi_l = n_l$, $l = 1, 2$, $n_l = 0, 1, 2, 3, 4$ is given by $H_{n_l}^l = \sum d_{\zeta_{n_l}}^l$. According to (29), (33), (36) and (37), we have the following formulations

$$\{\chi_l\}_{l=1}^L = \begin{cases} 0 \rightarrow H_0^l = 0, \\ 1 \rightarrow H_1^l = 2^{B_l+1}(B_l + 1) - 2, \\ 2 \rightarrow H_2^l = [2^{2B_l+2}(B_l + 1) - 2^{B_l+2}(2 + B_l) + 6], \\ 3 \rightarrow H_3^l = [2^{B_l+1}(3 + B_l) - 6], \\ 4 \rightarrow H_4^l = 2, \end{cases} \quad (41)$$

and

$$\begin{aligned} \varphi_0^l &= 1, \varphi_1^l = 4(2^{B_l} - 1), \varphi_2^l = 4(2^{B_l} - 1)^2 + 2, \\ \varphi_3^l &= 4(2^{B_l} - 1), \varphi_4^l = 1, \sum_{n=1}^4 \varphi_n^l = 2^{2B_l+2}. \end{aligned} \quad (42)$$

According to (23) and (40), there are a total of 8 values of $\|\mathbf{x}^i - \mathbf{x}^j\|^2 / (2\sigma^2) = (n_1 + n_2) / (2\sigma^2)$ as

$$\begin{aligned} \bar{\zeta}_1 &= \frac{1}{2\sigma^2} \leftrightarrow n_1=1, n_2=0; n_1=0, n_2=1; \\ \bar{\zeta}_2 &= \frac{2}{2\sigma^2} \leftrightarrow n_1=2, n_2=0; n_1=0, n_2=2; n_1=n_2=1; \\ \bar{\zeta}_3 &= \frac{3}{2\sigma^2} \leftrightarrow n_1=3, n_2=0; n_1=0, n_2=3; \\ &\quad n_1=1, n_2=2; n_1=2, n_2=1; \\ \bar{\zeta}_4 &= \frac{4}{2\sigma^2} \leftrightarrow n_1=4, n_2=0; n_1=0, n_2=4; n_1=n_2=2; \\ &\quad n_1=1, n_2=3; n_1=3, n_2=1; \\ \bar{\zeta}_5 &= \frac{5}{2\sigma^2} \leftrightarrow n_1=4, n_2=1; n_1=1, n_2=4; \\ &\quad n_1=2, n_2=3; n_1=3, n_2=2; \\ \bar{\zeta}_6 &= \frac{6}{2\sigma^2} \leftrightarrow n_1=4, n_2=2; n_1=2, n_2=4; n_1=n_2=3; \\ \bar{\zeta}_7 &= \frac{7}{2\sigma^2} \leftrightarrow n_1=4, n_2=3; n_1=3, n_2=4; \\ \bar{\zeta}_8 &= \frac{8}{2\sigma^2} \leftrightarrow n_1=4, n_2=4. \end{aligned} \quad (43)$$

The total HD between \mathbf{x}^i and \mathbf{x}^j associated with $\bar{\zeta}_m = m/(2\sigma^2)$ $m = (1, \dots, 8)$ is expressed as

$$H_m = \sum_{(n_1, n_2)} (H_{n_1}^1 + H_{n_2}^2) = \sum \sum d_{\zeta_{n_1}}^1 + d_{\zeta_{n_2}}^2, n_1 + n_2 = m. \quad (44)$$

According to (41)-(44), the values of H_m $m = (1, \dots, 8)$ are calculated as

$$H_1 = (\sum \sum d_{\zeta_0}^1 + d_{\zeta_1}^2) + (\sum \sum d_{\zeta_1}^1 + d_{\zeta_0}^2) \\ = H_1^2 + H_1^1 = \sum_{l=1}^2 [2^{B_l+1}(B_l+1)] - 4, \quad (45)$$

$$H_2 = (\sum \sum d_{\zeta_0}^1 + d_{\zeta_2}^2) + (\sum \sum d_{\zeta_2}^1 + d_{\zeta_0}^2) + (\sum \sum d_{\zeta_1}^1 + d_{\zeta_1}^2) \\ = H_2^2 + H_2^1 + \sum_{s=1}^{\varphi_1^1} \sum_{t=1}^{\varphi_1^2} d_{\zeta_1}^{1,s} + d_{\zeta_1}^{2,t} \\ = H_2^2 + H_2^1 + \varphi_1^1 H_1^2 + \varphi_1^2 H_1^1, \quad (46)$$

$$H_3 = \sum \sum (d_{\zeta_0}^1 + d_{\zeta_3}^2) + \sum \sum (d_{\zeta_3}^1 + d_{\zeta_0}^2) \\ + \sum \sum (d_{\zeta_1}^1 + d_{\zeta_2}^2) + \sum \sum (d_{\zeta_2}^1 + d_{\zeta_1}^2) \\ = H_3^2 + H_3^1 + \varphi_2^2 H_1^1 + \varphi_1^1 H_2^2 + \varphi_1^2 H_2^1 + \varphi_2^1 H_1^2, \quad (47)$$

$$H_4 = \sum \sum (d_{\zeta_0}^1 + d_{\zeta_4}^2) + \sum \sum (d_{\zeta_4}^1 + d_{\zeta_0}^2) + \sum \sum (d_{\zeta_1}^1 + d_{\zeta_3}^2) \\ + \sum \sum (d_{\zeta_3}^1 + d_{\zeta_1}^2) + \sum \sum (d_{\zeta_2}^1 + d_{\zeta_2}^2) \\ = H_4^2 + H_4^1 + \varphi_3^2 H_1^1 + \varphi_3^1 H_1^2 + \varphi_1^1 H_3^2 + \varphi_1^2 H_3^1 + \varphi_2^2 H_2^1 + \varphi_2^1 H_2^2, \quad (48)$$

$$H_5 = \sum \sum (d_{\zeta_1}^1 + d_{\zeta_4}^2) + \sum \sum (d_{\zeta_4}^1 + d_{\zeta_1}^2) \\ + \sum \sum (d_{\zeta_2}^1 + d_{\zeta_3}^2) + \sum \sum (d_{\zeta_3}^1 + d_{\zeta_2}^2) \\ = \varphi_4^2 H_1^1 + \varphi_1^1 H_4^2 + \varphi_1^2 H_4^1 + \varphi_4^1 H_1^2 + \\ \varphi_3^2 H_2^1 + \varphi_2^1 H_3^2 + \varphi_2^2 H_3^1 + \varphi_3^1 H_2^2, \quad (49)$$

$$H_6 = \sum \sum (d_{\zeta_2}^1 + d_{\zeta_4}^2) + \sum \sum (d_{\zeta_4}^1 + d_{\zeta_2}^2) + \sum \sum (d_{\zeta_3}^1 + d_{\zeta_3}^2) \\ = \varphi_4^2 H_2^1 + \varphi_2^1 H_4^2 + \varphi_2^2 H_4^1 + \varphi_4^1 H_2^2 + \varphi_3^2 H_3^1 + \varphi_3^1 H_3^2, \quad (50)$$

$$H_7 = \sum \sum (d_{\zeta_3}^1 + d_{\zeta_4}^2) + \sum \sum (d_{\zeta_4}^1 + d_{\zeta_3}^2) \\ = \varphi_4^2 H_3^1 + \varphi_3^1 H_4^2 + \varphi_3^2 H_4^1 + \varphi_4^1 H_3^2, \quad (51)$$

$$H_8 = \sum \sum (d_{\zeta_4}^1 + d_{\zeta_4}^2) = \varphi_4^2 H_4^1 + \varphi_4^1 H_4^2 = 4. \quad (52)$$

Since we have

$$H_{\text{all}} = H_1 + H_2 + H_3 + H_4 + H_5 + H_6 + H_7 + H_8 \\ = \sum_{b=1}^{2(B_1+B_2)+4} C_{2(B_1+B_2)+4}^b b, \quad (53)$$

which can be seen in Proof 2 in the Appendix, the ABEP of the G-QSM for $L = 2$ is given by

$$P_b^{L=2} = D(\bar{\zeta}_1)F(\bar{\zeta}_1) + D(\bar{\zeta}_2)F(\bar{\zeta}_2) + \dots + D(\bar{\zeta}_8)F(\bar{\zeta}_8), \quad (54)$$

where $D(\bar{\zeta}_m) = H_m/R$ and can be obtained from (45)-(52).

ABEP of G-QSM for $L > 2$

For case of $L > 2$, we denote the two different G-QSM symbols by \mathbf{x}^i and \mathbf{x}^j :

$$\mathbf{x}^i = \begin{bmatrix} \mathbf{x}_1^i \\ \vdots \\ \mathbf{x}_L^i \end{bmatrix}, \mathbf{x}^j = \begin{bmatrix} \mathbf{x}_1^j \\ \vdots \\ \mathbf{x}_L^j \end{bmatrix}. \quad (55)$$

Algorithm 1 ABEP of the proposed G-QSM system with $L > 2$

Input: $R = \sum_{l=1}^L 2B_l + 2$, $\varphi_{n_l}^l$, $l = (1, 2, \dots, L)$, $n_l = (0, 1, 2, 3, 4)$, $t = 0$, $P_b^{L>2} = 0$;
Output: $P_b^{L>2}$;
1: **for** $n_1 \in (0, 4)$ **do**
2: $\chi_1 = n_1$;
3: **for** $n_2 \in (0, 4)$ **do**
4: $\chi_2 = n_2$;
5: ...
6: **for** $n_L \in (0, 4)$ **do**
7: $\chi_L = n_L$;
8: $t = t + 1$;
9: $\chi = \chi_1 + \chi_2 + \dots + \chi_L$;
10: $\bar{\zeta} = \chi/2\sigma^2$;
11: $H_{(n_1, \dots, n_L)} = \sum \sum \dots \sum (d_{\zeta_{n_1}}^1 + d_{\zeta_{n_2}}^2 + \dots + d_{\zeta_{n_L}}^L)$;
12: $P_b^{L>2} = P_b^{L>2} + H_{(n_1, \dots, n_L)} F(\bar{\zeta})$;
13: **end for**
14: ...
15: **end for**
16: **end for**
17: $P_b^{L>2} = P_b^{L>2}/R$.

Then the ED between them can be expressed as

$$\chi = \|\mathbf{x}^i - \mathbf{x}^j\|^2 = \sum_{l=1}^L \|\mathbf{x}_l^i - \mathbf{x}_l^j\|^2 = \sum_{l=1}^L \chi_l, \quad (56)$$

where $\chi_l = \|\mathbf{x}_l^i - \mathbf{x}_l^j\|^2 = n_l$. The HD between \mathbf{x}^i and \mathbf{x}^j is:

$$d_{\zeta} = \sum_{l=1}^L d_{\zeta}^l, \quad (57)$$

where d_{ζ}^l denotes the HD between \mathbf{x}_l^i and \mathbf{x}_l^j .

According to (23) and (56), there are only $4L$ different values. For each combination (n_1, n_2, \dots, n_L) , $n_l = (0, 1, 2, 3, 4)$, we have $\chi = n_1 + n_2 + \dots + n_L$ and the corresponding total HD is given by

$$H_{(n_1, \dots, n_L)} = \sum \sum \dots \sum (d_{\zeta_{n_1}}^1 + d_{\zeta_{n_2}}^2 + \dots + d_{\zeta_{n_L}}^L) \\ = \prod_{v=2}^L \varphi_{n_v}^v \sum d_{\zeta_{n_1}}^1 + \prod_{v=1, v \neq 2}^L \varphi_{n_v}^v \sum d_{\zeta_{n_2}}^2 + \\ \prod_{v=1, v \neq 3}^L \varphi_{n_v}^v \sum d_{\zeta_{n_3}}^3 + \dots + \prod_{v=1, v \neq L}^L \varphi_{n_v}^v \sum d_{\zeta_{n_L}}^L \quad (58) \\ = \prod_{v=2}^L \varphi_{n_v}^v H_{n_1}^1 + \prod_{v=1, v \neq 2}^L \varphi_{n_v}^v H_{n_2}^2 + \\ \prod_{v=1, v \neq 3}^L \varphi_{n_v}^v H_{n_3}^3 + \dots + \prod_{v=1, v \neq L}^L \varphi_{n_v}^v H_{n_L}^L.$$

where H_n^l and φ_n^l can be obtained by (41) and (42), respectively. Since $\sum_{n_1=0}^4 \dots \sum_{n_L=0}^4 H_{(n_1, \dots, n_L)} = \sum_{r=1}^R C_R^r r$ (which can be seen in Proof 3 of the Appendix), the ABEP calculation of $P_b^{L>2}$ can be described in Algorithm 1.

Algorithm 2 The proposed M-OMP detector for the G-QSM system

Input: $\hat{\mathbf{y}}, \hat{\mathbf{H}}, m, \Lambda_0 = \phi, N_Q^l (l = 1, \dots, L), \lambda_0 = \phi, \mathbf{r}_0 = \hat{\mathbf{y}}, \text{flag}=0, \mathbb{I} = \phi, \mathbb{S} = \phi$

Output: $(l_{\mathcal{R}}^1, \dots, l_{\mathcal{R}}^L, l_{\mathcal{S}}^1, \dots, l_{\mathcal{S}}^L), (x_{\mathcal{R}}^1, \dots, x_{\mathcal{R}}^L, x_{\mathcal{S}}^1, \dots, x_{\mathcal{S}}^L)$.

```

1:  $\mathbf{l}_1 = \arg \text{sort}(|\mathbf{H}^H \mathbf{r}_0|)$ ;
2:  $\mathbf{l}_1^m = \mathbf{l}_1(1:m)$ 
3: for  $t_1 \in (1, m)$  do
4:    $l_{t_1} = \mathbf{l}_1^m(t_1)$ 
5:    $\Lambda_1 = [l_{t_1}]$ ;
6:    $\tilde{\mathbf{s}}_1 = \mathbb{D}((\mathbf{H}_{\Lambda_1}^H \mathbf{H}_{\Lambda_1})^{-1} \mathbf{H}_{\Lambda_1}^H \mathbf{y})$ ;
7:    $\mathbf{r}_1 = \mathbf{y} - \mathbf{H}_{\Lambda_1} \tilde{\mathbf{s}}_1$ ;
8:   Get the QSM symbol index  $T_1$  including all the
   indices of the  $t_1$ -th QSM symbol;
9:    $\mathbf{l}_2 = \arg \text{sort}(|\mathbf{H}^H \mathbf{r}_1|)$ ;
10:   $\mathbf{l}_2 = \text{setdiff}(\mathbf{l}_2, T_1)$ ;
11:   $\mathbf{l}_2^m = \mathbf{l}_2(1:m)$ 
12:  for  $t_2 \in (1, m)$  do
13:     $l_{t_2} = \mathbf{l}_2^m(t_2)$ 
14:     $\Lambda_2 = [l_{t_1}, l_{t_2}]$ ;
15:     $\tilde{\mathbf{s}}_2 = \mathbb{D}((\mathbf{H}_{\Lambda_2}^H \mathbf{H}_{\Lambda_2})^{-1} \mathbf{H}_{\Lambda_2}^H \mathbf{y})$ ;
16:     $\mathbf{r}_2 = \mathbf{y} - \mathbf{H}_{\Lambda_2} \tilde{\mathbf{s}}_2$ ;
17:    Get the QSM symbol indices  $T_2$  including all the
    indices of the  $t_1$ -th and  $t_2$ -th QSM symbols;
18:     $\mathbf{l}_3 = \arg \text{sort}(|\mathbf{H}^H \mathbf{r}_2|)$ ;
19:     $\mathbf{l}_3 = \text{setdiff}(\mathbf{l}_3, T_2)$ ;
20:     $\mathbf{l}_3^m = \mathbf{l}_3(1:m)$ 
21:    ...
22:  for  $t_{2L} \in (1, m)$  do
23:     $l_{t_{2L}} = \mathbf{l}_L^m(t_{2L})$ 
24:     $\Lambda_{2L} = [l_{t_1}, l_{t_2}, \dots, l_{t_{2L}}]$ ;
25:     $\tilde{\mathbf{s}}_{2L} = \mathbb{D}((\mathbf{H}_{\Lambda_{2L}}^H \mathbf{H}_{\Lambda_{2L}})^{-1} \mathbf{H}_{\Lambda_{2L}}^H \mathbf{y})$ 
26:     $\mathbb{I} = [\mathbb{I}; \Lambda_{2L}]; \mathbb{S} = [\mathbb{S}; \tilde{\mathbf{s}}_{2L}]$ ;
27:     $\text{flag} = \text{flag} + 1$ ;
28:    if  $\|\hat{\mathbf{y}} - \hat{\mathbf{H}}_{\Lambda_{2L}} \tilde{\mathbf{s}}_{2L}\|_F^2 \leq V_{\text{th}}$  then
29:       $(l_{\mathcal{R}}^1, \dots, l_{\mathcal{R}}^L, l_{\mathcal{S}}^1, \dots, l_{\mathcal{S}}^L) = \Lambda_{2L}$ ,
30:       $(x_{\mathcal{R}}^1, \dots, x_{\mathcal{R}}^L, x_{\mathcal{S}}^1, \dots, x_{\mathcal{S}}^L) = \tilde{\mathbf{s}}_{2L}$ ;
31:      return;
32:    else
33:       $\text{Dist}(\text{flag}) = \|\hat{\mathbf{y}} - \hat{\mathbf{H}}_{\Lambda_{2L}} \tilde{\mathbf{s}}_{2L}\|_F^2$ ;
34:    end if
35:  end for
36: end for
37: end for
38:  $[L_{\min}, s_{\min}] = \min(\text{Dist})$ .
39:  $(l_{\mathcal{R}}^1, \dots, l_{\mathcal{R}}^L, l_{\mathcal{S}}^1, \dots, l_{\mathcal{S}}^L) = \mathbb{I}_{L_{\min}}$ ,
40:  $(x_{\mathcal{R}}^1, \dots, x_{\mathcal{R}}^L, x_{\mathcal{S}}^1, \dots, x_{\mathcal{S}}^L) = \mathbb{S}_{L_{\min}}$ ;

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COMPRESSIVE SENSING DETECTOR FOR THE PROPOSED G-QSM SYSTEM

Proposed EM-OMP detector

In order to reduce the complexity of the ML detector, in this section CS aided detection is investigated in the context of our proposed G-QSM system. Firstly, the received signal of Eq. (2) is rewritten in the real-valued form as

$$\underbrace{\begin{bmatrix} \Re(\mathbf{y}) \\ \Im(\mathbf{y}) \end{bmatrix}}_{\hat{\mathbf{y}}} = \underbrace{\begin{bmatrix} \Re(\mathbf{H}) & -\Im(\mathbf{H}) \\ \Im(\mathbf{H}) & \Re(\mathbf{H}) \end{bmatrix}}_{\hat{\mathbf{H}}} \underbrace{\begin{bmatrix} \Re(\mathbf{x}) \\ \Im(\mathbf{x}) \end{bmatrix}}_{\hat{\mathbf{x}}} + \underbrace{\begin{bmatrix} \Re(\mathbf{n}) \\ \Im(\mathbf{n}) \end{bmatrix}}_{\hat{\mathbf{n}}}, \quad (59)$$

where $\hat{\mathbf{y}} \in \mathbb{R}^{2N_r \times 1}$, $\mathbf{H} \in \mathbb{R}^{2N_r \times 2N_t}$, $\hat{\mathbf{x}} \in \mathbb{R}^{2N_t \times 1}$ and $\hat{\mathbf{n}} \in \mathbb{R}^{2N_r \times 1}$. Therefore, the sparsity of $\hat{\mathbf{x}}$ is determined by $2L$, and the indices of the nonzero elements are distributed rather specifically, because the activated TA indices conveying real-valued APM symbols are located at $(1, \dots, N_Q^1), \dots, (\sum_{i=1}^{L-1} N_Q^i + 1, \dots, N_t)$ in the sequence, while the activated TA indices conveying the imaginary APM symbols are located at $(N_t + 1, \dots, N_t + N_Q^1), \dots, (\sum_{i=1}^{L-1} N_Q^i + N_t + 1, \dots, 2N_t)$ in the sequence. This fact can be beneficially exploited for CS detection.

In the G-QSM detection, the key issue is that of estimating the accurate activated indices by our CS algorithm. In our proposed EM-OMP algorithm, the Multipath Matching Pursuit (MMP) algorithm of [37] is invoked for the detection of the activated indices, which are introduced as follows.

Step 1: Estimate the first activated index as

$$\mathbf{l}_1 = \arg \text{sort}(|\mathbf{H}^H \mathbf{r}_0|). \quad (60)$$

Then we get the m candidates as $\mathbf{l}_1^m = \mathbf{l}_1(1:m)$.

Step 2: Estimate the second activated index. For the t_1 -th candidate in \mathbf{l}_1^m , m candidates are estimated as

$$\mathbf{l}_2 = \arg \text{sort}(|\mathbf{H}^H \mathbf{r}_1|), \mathbf{l}_2^m = \mathbf{l}_2(1:m), \quad (61)$$

where $\mathbf{r}_1 = \mathbf{y} - \mathbf{H}_{\Lambda_1} \tilde{\mathbf{s}}_1$ with $\Lambda_1 = [l_{t_1}]$ and $\tilde{\mathbf{s}}_1 = \mathbb{D}((\mathbf{H}_{\Lambda_1}^H \mathbf{H}_{\Lambda_1})^{-1} \mathbf{H}_{\Lambda_1}^H \mathbf{y})$.

Step 3: Repeat Step 2, until we get the $2L$ -th activated index. As a result, there may be m^{2L} possible index combinations. Since the locations of the activated indices are restricted, the estimated candidate indices in each step are from different level. Considering $N_t = 4$, $L = 2$, $m = 2$ and $\mathbf{l}_1^m = (1, 4)$ for example, if $\mathbf{l}_2 = (2, 3, 4, 5, 1, 6, 7, 8)$ is estimated in Step 2 based on the index '1', then we have $\mathbf{l}_2^m = (3, 4)$; If we get $\mathbf{l}_2 = (5, 4, 1, 2, 3, 4, 6, 7, 8)$ based on the index '4', we have $\mathbf{l}_2^m = (1, 2)$. Then we arrive at the following 4 index combinations $\{(1, 3), (1, 4), (4, 1), (4, 2)\}$. Finally, we remove the identical combinations and arrive at the final index combinations of $\{(1, 3), (1, 4), (2, 4)\}$.

Step 4: Obtain the accurate one from the m^{2L} possible candidates. For each combination $\Lambda_{2L} = [l_{t_1}, l_{t_2}, \dots, l_{t_{2L}}]$, the corresponding symbol vector is estimated as $\tilde{\mathbf{s}}_{2L} = \mathbb{D}((\mathbf{H}_{\Lambda_{2L}}^H \mathbf{H}_{\Lambda_{2L}})^{-1} \mathbf{H}_{\Lambda_{2L}}^H \mathbf{y})$. If $(\Lambda_{2L}, \tilde{\mathbf{s}}_{2L})$ satisfies

$$\|\hat{\mathbf{y}} - \hat{\mathbf{H}}_{\Lambda_{2L}} \tilde{\mathbf{s}}_{2L}\|_F^2 \leq V_{\text{th}}, \quad (62)$$

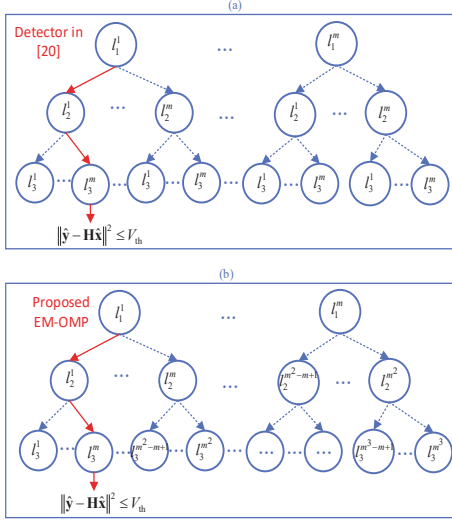


Fig. 4. Differences between the proposed EM-OMP detector and the detector in [20] at $L = 3$.

where $V_{th} = \beta N_r \sigma^2$ is a threshold defined in [36], $(\Lambda_{2L}, \tilde{s}_{2L})$ is taken as the final result; Otherwise, the specific combination associated with the smallest Euclidean distance will be considered as the final result. More specifically, the proposed EM-OMP detector is summarized in Algorithm 2.

Moreover, the difference between the proposed EM-OMP detector and the detector in [20] is shown in Fig. 4, where l_i^j represents the j -th candidate in the i -th step. Observe from Fig. 4 (a) that in [20], $(l_i^1, l_i^2, \dots, l_i^m)$ for the $(i-1)$ -th candidate are always the same and obtained by (60), while those in the EM-OMP detector may be different. Moreover, $m = N_Q$ ordered indices are used for constructing $(\prod_{i=1}^L N_Q^i)^2$ legitimate TACs for approaching the optimal performance in [20], while m is a small number in our proposed EM-OMP detector, imposing a reduced complexity.

Complexity Analysis

In this section, the complexity of the proposed EM-OMP detector is analyzed in terms of real-valued Floating point (Flops) operations, including real-valued multiplications and real-valued additions, since according to (56), the proposed EM-OMP detector operates in the real-valued domain. For specific real matrices $A \in M \times N$, $B \in N \times N$, $D \in M \times 1$, the complexity of the matrix operations AB , $A^H A$, B^{-1} and $\|D\|_F^2$ are expressed as

$$\begin{aligned} C_{AB} &= MN(N + N - 1) \\ C_{A^H A} &= MN^2 + N(M - \frac{N}{2}) - \frac{N}{2} \\ C_{B^{-1}} &= \frac{1}{3}N^3 + \frac{2}{3}N \\ C_{\|D\|^2} &= M + M - 1 \end{aligned} \quad (63)$$

According to (56) and Algorithm 2, the complexity of

the proposed EM-OMP detector is expressed as

$$\begin{aligned} C_{Pro} &= N_{avg} \sum_{t=1}^{2L} \underbrace{[(4N_r^2 + 2N_r)(t-1/2)]}_{C_{(H^H H)}} + \underbrace{(\frac{1}{3}t^3 + \frac{2t}{3})}_{C_{(\cdot)-1}} \\ &+ \underbrace{4N_r t^2 - 2N_r + 4N_r t - t}_{C_{(\cdot H_{\Lambda_t} y)}} + \underbrace{4N_r t}_{C_{(y-H_{\Lambda_t} s_t)}} + \underbrace{8N_t N_r}_{C_{(H^H r_i)}} \cdot 2L \\ &+ N_{avg} \underbrace{4N_r - 1}_{\|\cdot\|^2}. \end{aligned} \quad (64)$$

where N_{avg} denotes the average search candidates of the proposed detector.

For comparison, according to [20], the complexity of the ML, of the Bayesian CS (BCS) [35], ECS [20], and of the Compressive Sampling Matching Pursuit (CoSaMP) [34] detectors are expressed as

$$C_{ML} = (8N_r L + 4N_r - 1)2^R, \quad (65)$$

$$\begin{aligned} C_{ECS} &= \sum_{t=1}^{2L-1} \underbrace{[(4N_r^2 + 2N_r)(t-1/2)]}_{C_{(H^H H)}} + \underbrace{(\frac{1}{3}t^3 + \frac{2t}{3})}_{C_{(\cdot)-1}} \\ &+ \underbrace{4N_r t^2 - 2N_r + 4N_r t - t}_{C_{(\cdot H_{\Lambda_t} y)}} + \underbrace{4N_r t}_{C_{(y-H_{\Lambda_t} s_t)}} + \underbrace{8N_t N_r}_{C_{(H^H r_i)}} \cdot 2L + \\ &N_{avg} \underbrace{[(4N_r^2 + 2N_r)(2L-1/2)]}_{C_{(H^H H)}} + \underbrace{(\frac{1}{3}(2L)^3 + \frac{4L}{3})}_{C_{(\cdot)-1}} \\ &+ \underbrace{4N_r (2L)^2 - 2N_r + 4N_r 2L - 2L}_{C_{(\cdot H_{\Lambda_{2L}} y)}} + \underbrace{4N_r 2L}_{C_{(y-H_{\Lambda_{2L}} s_{2L})}} + \underbrace{4N_r - 1}_{\|\cdot\|^2}, \end{aligned} \quad (66)$$

$$\begin{aligned} C_{BCS} &= \underbrace{(4N_r^2 + 2N_r)(2N_t - 1/2)}_{C_{(H^H H)}} + 2N_r + \underbrace{(\frac{1}{3}(2N_r)^3 + \frac{4N_r}{3})}_{C_{(\cdot)-1}} \\ &+ 2N_t 4N_r + 2N_t 2N_r(4N_r - 1) + 4N_t, \end{aligned} \quad (67)$$

$$\begin{aligned} C_{Cosamp} &= w_{co} \underbrace{[(4N_r^2 + 2N_r)(K-1/2)]}_{C_{(H_{\Lambda_t}^H H_{\Lambda_t})}} + \underbrace{(\frac{1}{3}K^3 + \frac{2K}{3})}_{C_{(\cdot)-1}} \\ &+ \underbrace{4N_r K^2 - 2N_r + 4N_r K - K}_{C_{(\cdot H_{\Lambda_t} y)}} + \underbrace{4N_r K}_{C_{(y-H_{\Lambda_t} s_t)}} + \underbrace{4N_r - 1}_{\|\cdot\|^2}. \end{aligned} \quad (68)$$

In a word, the complexity order of these detectors are presented in Table I.

SIMULATION RESULTS

In this subsection, the performance results of the proposed G-QSM associated with different antenna configurations are presented and compared. In all the simulation results, Rayleigh channel with perfect channel state information is assumed. For simplicity, (N_t, N_r, L, M) represent the antenna setups.

Fig. 5 and Fig. 6 compare the performance of the proposed G-QSM system to that of the SM-VBLAST and VBLAST systems at 12 bpcu and 16 bpcu, respec-

TABLE I
COMPLEXITY ORDER OF DIFFERENT DETECTORS FOR G-QSM SYSTEM

Detector	Complexity order (Flops)
ML	$O(N_r L 2^R) + O(N_r 2^R)$
CoSaMP	$O(N_r^2 K) + O(N_r K) + O(K^3)$ $+ O(N_r K^2) + O(N_r K) + O(N_r)$
BCS	$O(N_r^3) + O(N_r^2 N_t) + O(N_r N_t)$ $+ O(N_r) + O(N_t)$
ECS	$O(N_{avg} N_r^2 L) + O(N_{avg} L^3) + O(N_{avg} N_r L^2)$ $+ O(N_{avg} N_t N_r L) + O(N_{avg} N_r)$
Proposed	$O(N_{avg} N_r^2 L) + O(N_{avg} L^3) + O(N_{avg} N_r L^2)$ $+ O(N_{avg} N_t N_r L) + O(N_{avg} N_r)$

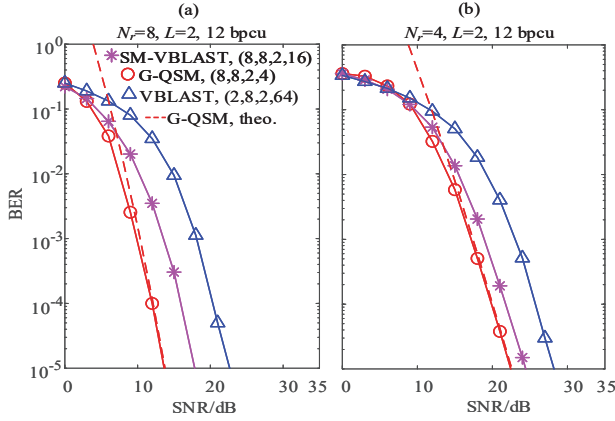


Fig. 5. Performance comparison of the proposed G-QSM system and of the SM-VBLAST and VBLAST systems having $L = 2$ at 12 bpcu; (a) $N_r = 8$; (b) $N_r = 4$.

tively³. Specifically, in Fig. 5 (a), the parameters of (8, 8, 2, 4), (8, 8, 2, 16) and (2, 8, 2, 64) are employed for the proposed G-QSM system, for SM-VBLAST and for the VBLAST systems to obtain the same normalized throughput of 12 bpcu. In Fig. 5 (b), the parameters of (8, 4, 2, 4), (8, 4, 2, 16) and (2, 4, 2, 64) are employed for the proposed G-QSM, SM-VBLAST and VBLAST systems to obtain the same normalized throughput of 12 bpcu. In Fig. 6 (a), the parameters of (8, 8, 2, 16), (8, 8, 2, 64) and (2, 8, 2, 256) are employed for the proposed G-QSM system, SM-VBLAST and VBLAST systems to obtain the same normalized throughput of 16 bpcu. In Fig. 6 (b), the parameters of (8, 4, 2, 16), (8, 4, 2, 64) and (2, 4, 2, 256) are employed for the proposed G-QSM, SM-VBLAST and VBLAST systems to obtain the same normalized throughput of 16 bpcu. ML detectors are employed in the above-mentioned setups. Moreover, the analytical ABEP performances of the proposed G-QSM system are added as benchmarks, which become quite accurate upon increasing the SNR values. Observe from Fig. 5 that the proposed G-QSM system using QPSK is capable of providing a 2.5 dB and 7.5 dB gain over the SM-VBLAST and VBLAST systems at $\text{BER}=10^{-5}$ for both over-determined

³Our comparison is limited to V-BLAST and SM-VBLAST associated with the same value of L because (1) they have the same power consumption; (2) they have the same inter-antenna interference; (3) they have the same inter-antenna synchronization complexity.

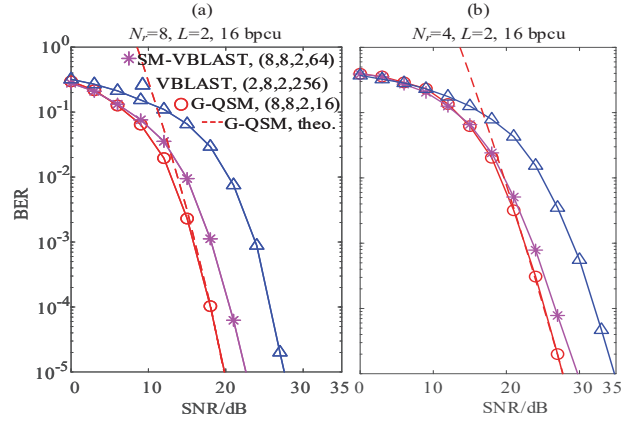


Fig. 6. Performance comparison of the proposed G-QSM system and of the SM-VBLAST and VBLAST systems having $L = 2$ at 16 bpcu; (a) $N_r = 8$; (b) $N_r = 4$.

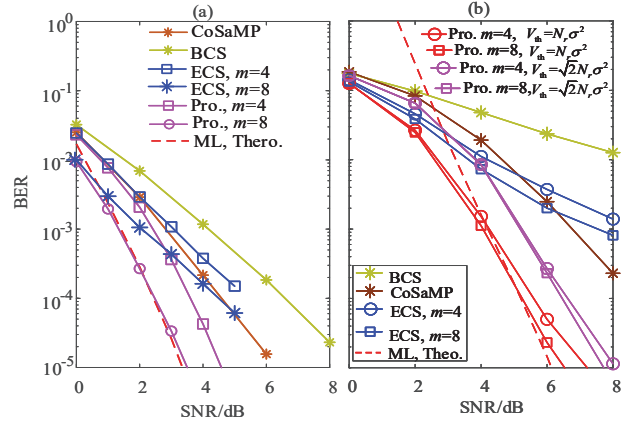


Fig. 7. Performance comparison of the proposed G-QSM system with different detectors. a) $N_t = 128, N_r = 32, L = 1, M = 4$ at 16 bpcu; b) $N_t = 128, N_r = 32, L = 2, M = 4$ at 28 bpcu.

and under-determined setups. This trend can also be observed for the G-QSM system using 16-QAM in Fig. 6.

Figs. 7-8 compare the performance vs complexity of the different detectors used by the proposed G-QSM system having $N_t = 128, N_r = 32, M = 4$ for different values of L , respectively. In massive MIMO setups, the ML detector becomes impractical, hence only the theoretical curves of the ML detector are added as benchmarks. Furthermore, $V_{th} = N_r \sigma^2$ is employed in the proposed EM-OMP detector in Fig. 7 (a) and Fig. 8(a). For the sake of comparing scenarios of similar complexity, the same values of m are selected for the ECS detector of [20]. Since m is much smaller than N_Q , the ECS of [20] exhibits an error floor in the high-SNR region. Observe from Fig. 7 that the performance of the proposed EM-OMP detector improves as m increases and it provides significant performance gains over the existing CoSaMP, BCS, ECS detectors at a lower complexity for both of the above-mentioned setups. For the case of $L = 1$, the transmission rate is 16 bpcu, and the complexity of the ML detector is around 2.5×10^7 Flops/bpcu. By contrast, the complexity becomes around 1.7×10^{11} Flops/bpcu for the case of $L = 2$. Observe from Figs. 7-8 that he

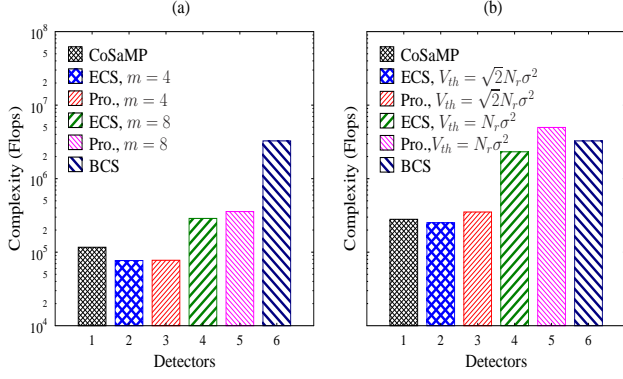


Fig. 8. Complexity comparison of the proposed G-QSM system with different detectors. a) $N_t = 128, N_r = 32, L = 1, M = 4$ at 16 bpcu; b) $N_t = 128, N_r = 32, L = 2, M = 4$ at 28 bpcu.

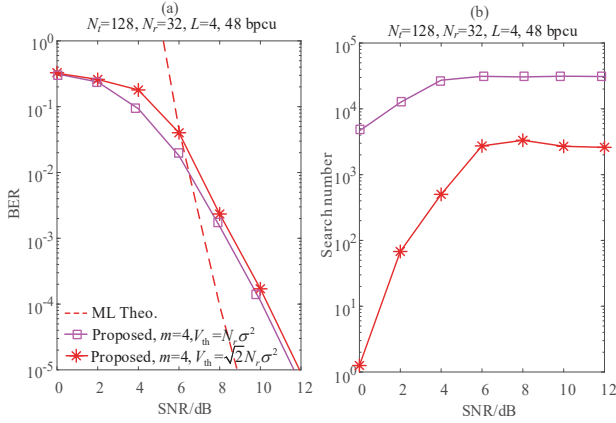


Fig. 9. Performance and complexity comparison of the proposed G-QSM system with different detectors having $N_t = 128, N_r = 32, L = 4, M = 4$ at 48 bpcu.

proposed EM-OMP detector using $m = 8, V_{th} = N_r \sigma^2$ is capable of approaching the performance of the ML detector, despite its more than 99% complexity reduction. Furthermore, observe from Figs. 7-8 (b) that the proposed EM-OMP detector using $m = 4, V_{th} = \sqrt{2} N_r \sigma^2$ is capable of achieving a significant complexity reduction over its counterpart using $m = 4, V_{th} = N_r \sigma^2$ at a modest performance loss.

Fig. 9 characterizes the performance of the G-QSM system with proposed detector for $N_t = 128, N_r = 32, M = 4, L = 4$. For the case of $L = 4$, the transmission rate is 48 bpcu, where the number of search candidates for the ML detector is 2^{48} . Observe from Fig. 9 that the number of search steps carried out by the proposed EM-OMP detector having $m = 4, V_{th} = \sqrt{2} N_r \sigma^2$ is reduced to as few as 10^3 from 2^{48} candidates at the cost of around 3 dB performance loss.

In order to provide further insights, Figs. 10-11 compare the performance and complexity of the different detectors for the G-QSM system having $N_t = 1024, N_r = 32, M = 4$ at different values of L . For the case of $L = 1$, the transmission rate is 22 bpcu, while the number of search

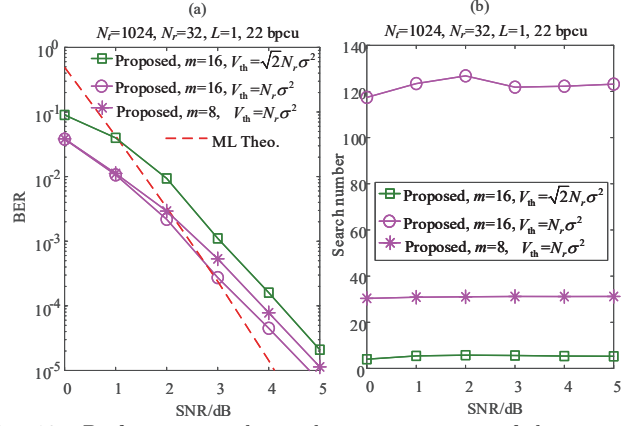


Fig. 10. Performance and complexity comparison of the proposed G-QSM system with different detectors having $N_t = 1024, N_r = 32, L = 1, M = 4$ at 22 bpcu.

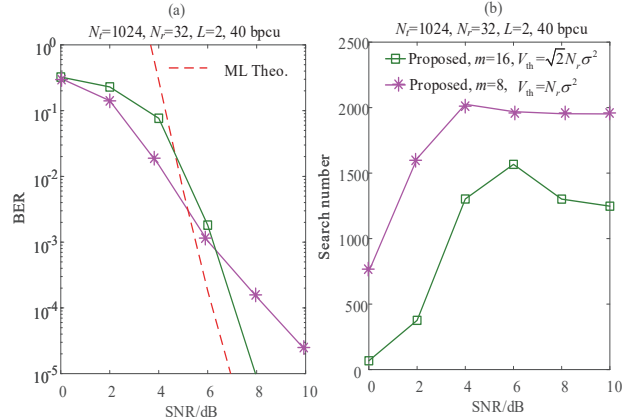


Fig. 11. Performance and complexity comparison of the proposed G-QSM system with different detectors having $N_t = 1024, N_r = 32, L = 2, M = 4$ at 40 bpcu.

candidates of the ML detector is 2^{22} . Observe from Fig. 10 that the number of search steps carried out by the proposed EM-OMP detectors having $m = 16, V_{th} = \sqrt{2} N_r \sigma^2$, $m = 8, V_{th} = N_r \sigma^2$ and $m = 16, V_{th} = N_r \sigma^2$ is reduced to as few as 5, 30 and 120 from 2^{22} candidates at the cost of around 1.2 dB, 1 dB, 0.8 dB performance loss. For the case of $L = 2$, the transmission rate is 40 bpcu, while the number of search candidates of the ML detector is 2^{40} . Observe from Fig. 11 that the number of search steps to be carried out by the proposed EM-OMP detectors having $m = 16, V_{th} = \sqrt{2} N_r \sigma^2$, $m = 8, V_{th} = N_r \sigma^2$ is reduced to about 10^3 and 2×10^3 from 2^{40} candidates at the cost of around 3 dB and 1 dB performance loss. In conclusion, the proposed EM-OMP detector is eminently suitable for our proposed G-QSM system, since it strikes a compelling performance vs complexity trade-off.

Finally, Fig. 12 compares the performance of the proposed G-QSM system to that of the VBLAST system having the same number of RF chains at high transmission rates. Specifically, $N_t = 128, N_r = 32, L = 1, M = 4$, $N_t = 1024, N_r = 32, L = 1, M = 4$ are employed for the proposed G-QSM system, while $N_t = 2, N_r = 32, L = 2, M = 256$ and $N_t = 2, N_r = 32, L = 2, M = 2048$ are

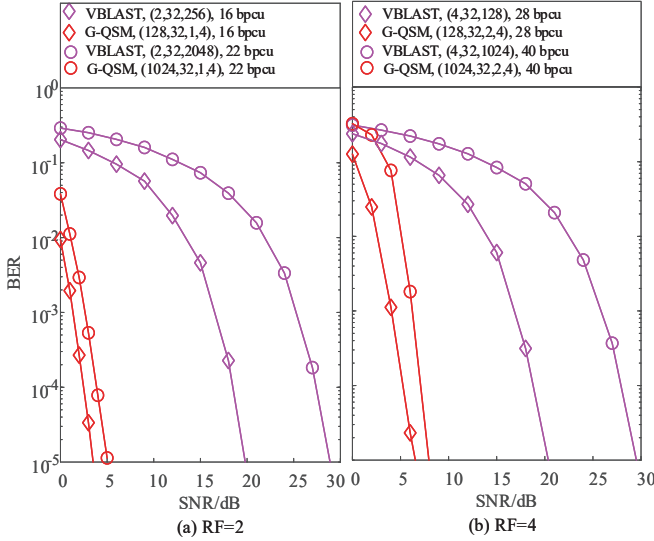


Fig. 12. Performance comparison of the proposed G-QSM system and of the SM-VBLAST and VBLAST systems having the same RF chains and transmission rate; a) RF=2; b) RF=4.

invoked for the VBLAST system to obtain the normalized throughput of 16 bpcu and 22 bpcu having 2 RF chains. Furthermore, $N_t = 128, N_r = 32, L = 2, M = 4$, $N_t = 1024, N_r = 32, L = 2, M = 4$ are employed for the proposed G-QSM system, while $N_t = 4, N_r = 32, L = 4, M = 128$ and $N_t = 4, N_r = 32, L = 4, M = 1024$ are used for the VBLAST system to obtain the normalized throughput of 28 bpcu and 40 bpcu having 4 RF chains. In order to compare systems of the same complexity, the efficient EM-OMP detectors are employed in the proposed G-QSM and MMSE detectors are employed in the VBLAST systems. For the case of 2 RF chains, observe from Fig. 12 (a) that the proposed G-QSM system is capable of outperforming the VBLAST system⁴ by 17 dB and 24 dB at 16 bpcu and 22 bpcu respectively at $\text{BER}=10^{-5}$ at the cost of an extra channel estimation action, which will be eliminated in our future research by conceiving non-coherent detection at 3 dB performance loss [41]. A similar trend can be observed for the case of 4 RF chains in Fig. 12 (b).

CONCLUSIONS

The new low-complexity, high-throughput G-QSM scheme has been proposed. The upper bound of the ABEP of the proposed G-QSM system has also been derived. Furthermore, an efficient CS detector termed as EM-OMP has been proposed for our G-QSM system. Both the theoretical and simulation results have shown that the proposed G-QSM scheme employing a large number of TAs is capable of providing a 20 dB gain over the V-BLAST

⁴Explicitly, we are not making the unrealistic claim that our scheme outperforms V-BLAST in general, only under the specific assumption of having a rather limited number of RF chains. As the affordable number of RF chains is increased, the throughput becomes quite high, which would necessitate a more advanced CS detector, which will be designed in our future research.

system having a small identical number of RF chains and the same throughput.

APPENDIX

Proof 1: $\sum_{b=1}^B C_B^b b = B2^{B-1}$.

Specifically, the elements of $C_B^b b = 1, \dots, B$ can be represented in a matrix-form as

$$\{C_B^b b\}_{b=1}^B \rightarrow \begin{bmatrix} C_B^1 & & & & & & \\ C_B^2 & C_B^2 & & & & & \\ \vdots & \vdots & \ddots & & & & \\ C_B^b & C_B^b & \cdots & C_B^b & & & \\ \vdots & \vdots & \vdots & \vdots & \ddots & & \\ C_B^{B-1} & C_B^{B-1} & C_B^{B-1} & C_B^{B-1} & \cdots & C_B^{B-1} & \\ C_B^B & C_B^B & C_B^B & C_B^B & C_B^B & \cdots & C_B^B \end{bmatrix}. \quad (69)$$

Then the elements of $2C_B^b b - BC_B^B$ ($b = 1, \dots, B$) can also be represented in a matrix-form as

$$\begin{bmatrix} C_B^1 & C_B^{B-1} & C_B^{B-1} & C_B^{B-1} & C_B^{B-1} & \cdots & C_B^{B-1} \\ C_B^2 & C_B^2 & C_B^{B-2} & C_B^{B-2} & C_B^{B-2} & \cdots & C_B^{B-2} \\ \vdots & \vdots & \ddots & \vdots & \vdots & \vdots & \vdots \\ C_B^b & C_B^b & \cdots & C_B^b & C_B^{B-b} & \cdots & C_B^{B-b} \\ \vdots & \vdots & \vdots & \vdots & \vdots & \vdots & \vdots \\ C_B^{B-1} & C_B^{B-1} & C_B^{B-1} & C_B^{B-1} & \cdots & C_B^{B-1} & C_B^1 \\ C_B^B & C_B^B & C_B^B & C_B^B & C_B^B & \cdots & C_B^B \end{bmatrix}. \quad (70)$$

Since $\sum_{b=1}^B C_B^b = 2^B - 1$, $C_B^b = C_B^{B-b}$, according to (69) and (70), we have

$$2 \sum_{b=1}^B C_B^b b - C_B^B B = B \sum_{b=1}^B C_B^b = B(2^B - 1). \quad (71)$$

Hence, we have $\sum_{b=1}^B C_B^b b = B2^{B-1}$.

Proof 2: $H_{\text{all}} = \sum_{t=1}^8 H_t = \sum_{r=1}^R C_R^r r$ with $R = \sum_{l=1}^2 (2B_l + 2)$. According to (45)-(52), the value of H_{all} can be expressed as

$$\begin{aligned} H_{\text{all}} &= H_1 + H_2 + H_3 + H_4 + H_5 + H_6 + H_7 + H_8 \\ &= (1 + \varphi_1^2 + \varphi_2^2 + \varphi_3^2 + \varphi_4^2)(H_1^1 + H_2^1 + H_3^1 + H_4^1) + \\ &\quad (1 + \varphi_1^1 + \varphi_2^1 + \varphi_3^1 + \varphi_4^1)(H_1^2 + H_2^2 + H_3^2 + H_4^2) \\ &= 2^{2B_2+2} 2^{2B_1+1} (2B_1 + 2) + 2^{2B_1+2} 2^{2B_2+1} (2B_2 + 2) \\ &= 2^{2B_2+2+2B_1+1} (2B_1 + 2 + 2B_2 + 2) \\ &= \sum_{u=1}^{2(B_1+B_2)+4} C_{2(B_1+B_2)+4}^u u. \end{aligned} \quad (72)$$

Proof 3: $H_{\text{all}} = \sum_{t=1}^m H_t = \sum_{r=1}^R C_R^r r$ with $R = \sum_{l=1}^L (2B_l + 2)$. According to (58), the value of H_{all} can be

expressed as

$$\begin{aligned}
 H_{\text{all}} &= \sum_{n_1=0}^4 \cdots \sum_{n_L=0}^4 H_{(n_1, \dots, n_L)} \\
 &= \sum_{n_1=0}^4 \cdots \sum_{n_L=0}^4 \left[\prod_{v \neq 1} \varphi_{n_v}^v H_{n_1}^1 + \cdots \prod_{v \neq L} \varphi_{n_v}^v H_{n_L}^L \right] \\
 &= \left(\sum_{n_2=0}^4 \cdots \sum_{n_{L-1}=0}^4 \prod_{v=2}^L \varphi_{n_v}^v \right) \sum_{n_1=0}^4 H_{n_1}^1 + \cdots \\
 &\quad + \left(\sum_{n_2=0}^4 \cdots \sum_{n_{L-1}=0}^4 \prod_{v \neq L} \varphi_{n_v}^v \right) \sum_{n_L=0}^4 H_{n_L}^L.
 \end{aligned} \tag{73}$$

Since we have

$$\begin{aligned}
 &\left(\sum_{n_2=0}^4 \cdots \sum_{n_{L-1}=0}^4 \prod_{v=2}^L \varphi_{n_v}^v \right) \sum_{n_1=0}^4 H_{n_1}^1 \\
 &= \left(\sum_{n_2=0}^4 \cdots \sum_{n_{L-1}=0}^4 \prod_{v=2}^{L-1} \varphi_{n_v}^v \right) \sum_{n_L=0}^4 \varphi_{n_L}^L \sum_{n_1=0}^4 H_{n_1}^1 \\
 &= \sum_{n_2=0}^4 \varphi_{n_2}^2 \cdots \sum_{n_{L-1}=0}^4 \varphi_{n_{L-1}}^{L-1} \sum_{n_L=0}^4 \varphi_{n_L}^L \sum_{n_1=0}^4 H_{n_1}^1 \\
 &= 2^{(2B_2+2)} \cdots 2^{(2B_{L-1}+2)} 2^{(2B_L+2)} 2^{(2B_1+1)} (2B_1+2) \\
 &= 2^{\sum_{l=1}^L (2B_l+2)-1} (2B_1+2),
 \end{aligned} \tag{74}$$

$$\begin{aligned}
 &\left(\sum_{n_1=0}^4 \sum_{n_3=0}^4 \cdots \sum_{n_L=0}^4 \prod_{v=1, v \neq 2}^L \varphi_{n_v}^v \right) \sum_{n_2=0}^4 H_{n_2}^2 \\
 &= 2^{\sum_{l=1}^L (2B_l+2)-1} (2B_2+2), \\
 &\vdots \\
 &\left(\sum_{n_2=0}^4 \cdots \sum_{n_{L-1}=0}^4 \prod_{v=1, v \neq L}^L \varphi_{n_v}^v \right) \sum_{n_L=0}^4 H_{n_L}^L \\
 &= 2^{\sum_{l=1}^L (2B_l+2)-1} (2B_L+2).
 \end{aligned} \tag{75}$$

(73) can be finally represented as

$$\begin{aligned}
 H_{\text{all}} &= \sum_{n_1=0}^4 \sum_{n_2=0}^4 \cdots \sum_{n_L=0}^4 [\sum \cdots \sum (d_{\zeta n_1}^1 + d_{\zeta n_2}^2 + \cdots d_{\zeta n_L}^L)] \\
 &= 2^{\sum_{l=1}^L (2B_l+2)-1} \sum_{l=1}^L (2B_l+2) \\
 &= \sum_{r=1}^R C_R^r r.
 \end{aligned} \tag{76}$$

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