EdgeDrones: Co-Scheduling of Drones for Multi-Location Aerial Computing Missions

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ABSTRACT

Low altitude platform (LAP) unmanned aerial vehicles (UAVs), also called *drones*, are currently being exploited by Edge computing (EC) systems to execute complex resource-hungry use cases, such as virtual reality, smart cities, autonomous vehicles, etc., by attaching portable edge devices on them. However, a typical drone has limited flight time, coupled with the resource-constrained attached edge device, which can jeopardize aerial computing missions if they are not holistically taking into consideration. Moreover, the fundamental challenge is how to co-schedule multi-drone among multi-location where EC services are needed, such that drones are scheduled to maximize the utility from the activities while meeting computing resource and flight time constraints. Therefore, for a given fleet of drones and tasks across disjointed target locations in a city, we derive a machine learning (ML) linear regression model that estimates these tasks resource requirement and excution time. Leveraging this estimation values, we jointly consider each drone's flight time availability and its attached edge device resource capacity, and formulate a novel Multi-Location Capacitated Mission Scheduling Problem (MLCMSP) that selects suitable drones and co-schedules their flight routes with the least total distance to visit and execute tasks at the target locations. Then, we show that faster scheduling and execution of complex tasks at each location, while considering the inter-task dependencies is important to achieve effective solution for our MLCMSP. Hence, we further propose EdgeDrones, a variant bin-packing optimization approach through gang-scheduling of inter-dependent tasks that *co-schedules* and *co-locates* tasks tightly so as to achieve faster execution time, as well as to fully utilize available resources. Extensive experiments on Alibaba cluster trace with information on task dependencies (about 12,207,703 dependencies) show that EdgeDrones achieves up to 73% higher resource utilization, up to 17.6 times faster executions, and up to 2.87 times faster flight travel time compared to the baseline approaches.

1. Introduction

Edge computing (EC) is a distributed computing model which places cloud computing [1, 2] services closer to data sources so as to achieve faster response times and real-time insights. Many latency-sensitive applications that process data from IoT devices and sensors, rely on heterogeneous edge resources in close proximity for faster response times and to promote rapid development. To this end, several EC devices of various sizes and capacities have emerged. For example, AWS Snowcone¹, Azure Stack Edge mini², etc., are portable EC devices that weighs about 2 \sim 3 kg but are inherently resource-constrained relative to their onpremise counterparts, i.e., AWS Snowball³, etc. Nonetherless, EC systems are currently exploiting attaching these portable edge devices on low altitude platform (LAP) unmanned aerial vehicles (UAVs) or drones to execute complex resource-hungry use cases, such as cyber-physical systems [3], virtual reality [4], smart vehicles [5], face recognition [6], smart cities [7], etc. In addition, an autonomous drone

technology called *Drone-in-a-box*⁴, is currently being exploited for aerial EC missions [5, 8, 9]. A drone-in-a-box system can be deployed autonomously from a box that serves as a landing pad and charging base (i.e., depot) to perform on-demand computation activities in a city. An activity involves visiting points of interest (i.e., target locations), hovering and interacting with end devices to execute tasks on its attached edge device(s). After completing the tasks, the results are immediately and deterministically communicated back to the end device, then it returns to its box or depot. However, a typical drone has a limited flight time due to power factor which can jeopardize the entire mission if it is not taking into consideration [10, 11, 12]. Hence, the critical issue is how to assign missions and optimal routes for multiple drones to visit a set of locations, so that they can complete their tasks, subject to the flight time and attached edge resource constraints of each drone, without jeopardizing application performance.

Therefore, for a given fleet of drones and multi-task across disjointed locations in a city, we propose a novel Multi-Location Capacitated Mission Scheduling Problem (MLCMSP) that co-schedules their optimal flight routing among the locations, such that the drones can visit the locations and complete the tasks, within their flight times and computing resource constraints, while maximizing the

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¹https://aws.amazon.com/snowcone/

²https://azure.microsoft.com/en-us/products/azure-

stack/edge/#overview

³https://aws.amazon.com/snowball/

⁴https://en.wikipedia.org/wiki/Drone_in_a_Box

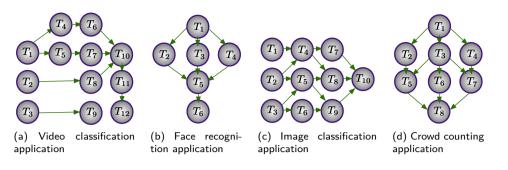


Figure 1: DAG of representative applications.

total utilization. Specifically, our MLCMSP combines elements of Vehicle Routing Problem (VRP), which is a variant of the well known Traveling Salesman Problem (TSP) to find optimal routes for a set of vehicles and customers [13, 14]. Existing works have proposed routing strategies for UAV-enabled task offloading in edge-cloud computing systems [15, 16]. However, they have either focused on coscheduling a fleet of drones or task offloading to meet some specific objectives. It is important to note that a drone routing strategy for EC which do not consider the drones' flight time and resource capacity or charateristics of target tasks, such as dependencies, resource requirement, etc., can lead to loss of job or an incomplete mission [10, 11, 12]. Effective co-scheduling of a given fleet of drones for EC missions across multi-location, requires jointly optimization of the following; (i) update information of each drones' flight time availability and its attached edge device(s)⁵ resource capacity or availability, (ii) locations of end devices requesting EC services interms of flight distance, flight travel time, etc., and (iii) their tasks resource requirement and execution time estimations, so as to select drones which can maximize the utility from the activities. A disjointed approach which interacts individually with each drone, would exhibit high computation complexity and is far from trivial to realize [11, 12]. For this reason, we wish to consider an approach which seamlessly integrate all end devices, service entities and edge resources running across multiple drones in a single pool, such that these information can be holistically obtained and monitored from a single control plane (CP), where it can be used for decision making on efficient mission planning and assignment. This approach is called *Edge Federation* (EF) [17, 18]. For example, recently introduced edge computing frameworks, i.e., KubeEdge, MicroK8s, etc, have the capabilities of integrating service entities and edge resources running across multiple drones, run containerized tasks and eliminate provider lock-in situations. EF can enable effective co-scheduling of multiple drones, by selecting a minimum number of drones which can maximize the utility for any given activities. Hence, a drone can be assigned multiple disjointed locations as part of its mission.

An important challenge is developing an efficient scheduling strategy that can place and execute complex applications on the attached edge devices in a timely manner, while efficiently managing available resources as drones visit their assigned locations. For example, modern applications (i.e., face recognition [6], image classification [19], crowd counting [20], etc.,), as shown in Fig (1) are becoming more complex in nature, structured on microservices architectural style, consisting of a large number of inter-dependent applications and often latency-sensitive [21, 22, 23]. It is naturally important to intelligently schedule such interdependent applications in a best possible way, such that they are quickly executed and immediately sent back to the IoT and end devices. Existing scheduling approaches which do not consider such task dependencies, co-location strategy or which randomly deploy tasks to any available resources can easily result in delay, fragmentation and over-allocation of resources, hence jeopardizing the application performance, given the drones' flight time and resource constraints. To address this challenge, we first estimate tasks resource requirement and execution time at target locations, using linear regression machine learning (ML) model. These estimation values, as well as the drones' flight time, flight distance to target locations, and attached edge resource availability are used as inputs to plan missions for a captive set of drones to accomplish any given activities. One drawback of this concept is that inaccurate estimations of tasks resource requirement and execution time at target locations, could also jeopardize the entire missions for selected drones. Similarly, if the tasks are scheduled naively, e.g., in an edge deployment which can only execute one task or job at any time and where each task is scheduled individually [24, 25], the system might not yield an optimal performance. Therefore, we first investigate the accuracy of our trained linear regression ML model for estimating multi-task resource requirement and execution time, using the normalized absolute estimate error (NAEE) method. This serves as the estimation accuracy measure for the trained linear regression ML model. We further propose EdgeDrones, which extends the state-ofthe-arts by providing an intelligent dependency-aware multitask scheduling and co-location scheme to achieve high resource utilization and faster execution of tasks.

⁵A typical drone-based edge deployment can attach one or more or different combination of portable edge devices, depending on the drones' load capacity.

In particular, we show that EF [17, 18] and ML techniques [26, 27] can help aerial edge systems to achieve effective co-scheduling and optimal route planning for a fleet of autonomous drones to accomplish stochastic service requests from end devices across target locations in a city. With limited edge resources and drones' flight time, it is necessary to consider task dependencies in drone-based EC task offloading, by jointly optimizing the drones' flight time and resource availability, such that all the tasks can be intelligently scheduled and fast executed with minimum resources before the drone returns for recharging or departs for another assigned location. Hence, our aim is to plan optimal flight routes for drones to execute all the tasks by considering dependencies and resource demands, such that the actual scheduling and execution time is minimized, and is much less than the drones' flight time. In summary, to achieve our EdgeDrones implementation, we address the following critical areas:

- We propose an integrated system with global information, through the joint optimization of all service entities, and formally define the novel Multi-Location Capacitated Mission Scheduling Problem (MLCMSP) that selects and co-schedules optimal flight routes for a fleet of drones for any given aerial computing missions in a city.
- Specifically, we derive a multi-task machine learning (ML) execution time and resource requirement estimation, to aid MLCMSP to select drones with requisite resource availability which can maximize the utility from the aerial activities.
- To guarantee optimal usage of cluster resources and faster execution of tasks, we further propose a variant bin-packing optimization approach through gang-scheduling of multi-dependent tasks, which co-schedules and co-locate tasks firmly on available nodes, so as to avoid resource wastage.
- We show that EdgeDrones is capable of minimizing the actual completion time of multi-dependent tasks using minimum resources, and we conduct extensive experiments to compare the performance of our Edge-Drones with several existing approaches using realworld data-trace from Alibaba cluster trace⁶, which provides information on task dependencies.

The rest of the paper is organized as follows. In Section 2, we discuss the related work. In Section 3, we present some preliminaries on task dependency-awareness and discuss our motivation. In Secton 4, we detail our proposed EdgeDrones for achieving high resource utilization and minimizing the execution times of applications deployed on federated aerial edge resources. In Section 5, we compare the performance of our proposed EdgeDrones against that of several state-of-the-art approaches through extensive experiments. Finally, we conclude the paper in Section 6.

2. Related Works

UAV or drone-based edge computing deployments are gaining increasingly popularity due to their autonomous navigation, low cost, mobility, flexibility and adaptive altitude to deliver faster execution closer to data sources. They are currently being exploited for several use cases, i.e., task offloading, data caching, data streaming, etc. For example, the authors of [28] proposed a UAV-enabled edge network to minimize the system-wide computation cost by efficient task offloading and deployment. They have formulated and solved this problem as a stochastic game. The work [3] proposed a joint trajectory, task offloading and caching optimization in a UAV-enabled edge for Cyber-Physical System. They have proposed this to realize energy-efficient performance of the UAV. In [4], the authors explored the UAV-assisted edge and streaming for virtual reality. They formulated this problem as a joint joint UAV placement, edge resource allocation, and 360-degree video content laver assignment. They aim to select the allocation of computing and communications resources, such that the delivered quality of experience (QoE) is maximized.

UAV or drone-enabled multi-task offloading schemes in edge systems can benefit from joint optimization of drones' flight time, resource availability status, multi-task's resource requirement and execution time, such that drone with sufficient resource availability can be deloved to conduct efficient execution request. A-priori information about task execution time is mostly important for drone-based edge deployments [11, 12]. This is because a typical drone has limited flight time, and could possibly lead to a delayed task execution if it is not taken into consideration [10]. In particular, optimal route can be planned and assinged to ensure that drones can complete their tasks, given their flight time constriants. There exist works that explore routing and trajectory scheduling for drone-enable edge systems. For example, the work [15] proposed a mssion scheduling problem (MSP) that co-schedules the flight routes of drones to visit target locations and record videos. Similarly, the work [16] proposed an online algorithm for UAV swarms to jointly optimize the task offloading and multi-hop routing scheduling. In [3], the authors proposed a joint optimization of drone's 3D trajectory scheduling and the task-cache strategies to minimize its total energy consumption. The work [29] formulated an optimization problem to minimize the total energy consumption of a UAV through joint partitioning and UAV trajectory scheduling. In [30], the authors presented a distributed task offloading and path planning algorithm to provide computational support to large-scale IoT nodes. However, these schemes do not consider drones' flight time and assume a drone can fly for unlimited amount of time, which can lead to delay or loss of job due to drones' limited flight time [10, 11, 12].

For multi-task at target locations, an accurate execution time and resource demand estimation is mostly needed to schedule routes for drone(s) and a-priori to conduct efficient multi-task scheduling. Consequently, existing researches

⁶https://github.com/alibaba/clusterdata

have proposed a huge number of learning methods to estimate task's resource requirements and execution time, based on collaborative learning (CL) [11, 18, 31], machine learning (ML) [12, 26, 27, 32], incremental learning (IL) [33], scheduling [34, 35, 36, 37] and statistical models [38]. Our previous works [11, 12] focused on of multidependent tasks orchestration in autonomous drone-enabled edge computing system, while considering the drones' flight time, so as to avoid loss of jobs [10]. Specifically, in [11], we have proposed a multi-output linear regression model based on CL to estimate multi-dependent task's resource requirement and execution time, to select the closest drone deployment having matching resource availability and flight time to execute ready tasks at a given time. In [12], we have proposed a ML based multi-dependent tasks dispatching over a federated autonomous drone-enabled edge computing platform, using the total estimated value of the multidependent tasks' execution time to select a suitable drone.

With limited edge resources, is it also important to avoid any form of resource wastage, i.e., resource under utilization. Efficiently managing edge resources directly dictates service quality and performance [39]. As a result, task co-location has gained attention both in academia and industry as an optimal solution for improving resource utilization and system throughput in distributed systems. However, effective task co-location is a non-trivial task, as it requires an understanding of the computing resource requirement of the corunning tasks, in order to determine how many of them can be co-located. To this end, a tasks co-location mechanism was proposed in [40], where it was showed that by accurately estimating the resource level needed, a co-location scheme can effectively determine how many tasks can be co-located on the same host to improve the system throughput, by taking into consideration the memory and CPU requirements of co-running tasks. With the aim to maximize the resource utilization, the authors of [41] utilized reinforcement learning to co-locate interactive services with batched ML workloads. Our previous works [42, 43] focused on workload co-location in cloud environment, rather than edge systems. To further improve edge resource management, a resource management scheme was proposed in [17, 18] which unifies distributed edge resources, such that they are holistically managed. Our previous work [17] proposed a dependencyaware task scheduling in such unified system. Modern applications are usually structured with inter-task dependencies, whereby a task depends on an input from other task(s). A huge number of existing works, i.e., [21, 22, 23, 44] have tackled the problem scheduling such inter-dependent tasks or multi-dependent tasks, and their common goal is to formulate a scheduling decision that minimizes the average completion time of such tasks.

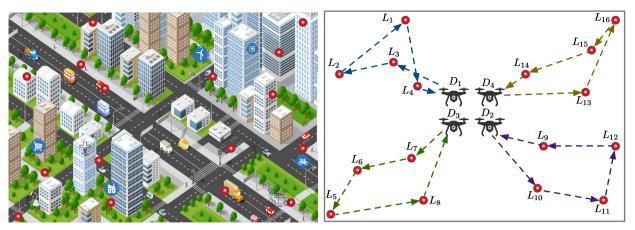
Existing works on UAV-enabled approaches for task offloading and execution in multi-edge deployments do not jointly consider tasks dependencies, do not unify service entities and distributed edge resources, such that they are holistically managed and monitored from a single control plan (CP), where such information can be utilized to co-schedules multi-drone, co-locate multi-task effectively. This motivates our research to extend existing schemes by proposing Edge-Drones. Specifically, we propose a learning-based multidrone route scheduling through a unified system, which include all service entities and resources running across the multi-drone, location of end devices and their applications. We further propose a variant bin-packing optimization approach through gang-scheduling of multi-dependent tasks, which quickly co-schedules and co-locate tasks firmly on available nodes, so as to avoid delay and resource wastage. We finally show that EdgeDrones is capable of minimizing the actual completion time of multi-dependent tasks using minimum resources through extensive experiments and comparison.

3. A Case Study on a Smart City

We consider a smart city senario, where multiple IoT and other end devices are deployed across the city to improve life standards of its citizens. For example, Toyota Motor Corporation has recently embarked on a new smart city project called Woven City⁷, where new technologies such as smart construction and manufacturing, smart homes, robotics, connectivity through AI, autonomy, etc are being deployed. EC provides a promising way of enabling these technologies by offering computing resources with low latency. Smart city solutions are increasingly integrating UAVs or droneenabled EC for enhanced performance [7]. Suppose at time t, there are updates from devices at multiple locations in the city, as shown in Fig. 2(a), drones equipped with EC devices can fly to these locations to render needed services in a timely manner. However, how to select, assign and route flight paths for drones to accomplish these tasks effectively is a major challenge. To address this, each task's resource requirement in terms of CPU and memory at each location L_i is estimated, i.e., $\langle \widetilde{c_i}, \widetilde{m_i} \rangle$, as well as its execution time E_{ex} . These values, as well as the location coordinates are utilized to select suitable drones (i.e., drones with sufficient flight time f_i and resource capacity (c, m), such that their optimal routes to visit the locations are scheduled, as shown in Fig. 2(b). Nevertheless, a routing problem with many locations can take a long time to solve. Therefore, for such problems, it is better to set a search limit which terminates the search after a specific length of time or number of solutions is returned.

Most importantly, as the drones embark on their missions, efficient scheduling strategy for complex applications at each location L_i is absolutely necessary to guarantee high application performance and successful completion of each drone's mission. For example, in Fig. 2(b), Drone 1 is assigned 4 locations (L_1 , L_2 , L_3 and L_4). Suppose at these locations, we have complex applications in the form of directed ayclic graphs (DAGs), as shown in Figs. 2(a), (b), (c) and (d), respectively, where each job typically consists of several tasks whose dependencies are expressed by DAG, i.e., the job in Fig. 1(a) consists of inter-task dependency depth γ of

⁷https://www.woven-city.global/



(a) Locations of end devices needing EC services in a smart city (b) Drones assignement, their optimal routes and flight paths

Figure 2: (a)Locations of end devices needing EC services, (b) drones assignement together with their optimal routes.

12, i.e., $(T_1, T_2, \dots, T_{12})$. Such complex inter-task dependencies with multi-dimensional resource demands, i.e., various amounts of CPU and memory resources, and communication requirements, make resource management in such a droneenabled EC system very challenging. Knowledge about these tasks characteristics, i.e., resource demands and dependencies, is necessary to pack or co-locate them effectively in available resources, ultimately to minimize their collective response times and improve the resource utilization [11, 12]. Clearly, tasks T_1 , T_2 and T_3 are independent task, i.e., no dependency, and they can be started without waiting for any other task(s), while tasks T_4 and T_5 depend on the completion of task T_1 . Similarly, task T_{10} depends on the completion of tasks T_6 , T_7 and T_8 . Hence a key objective of our EdgeDrones is to reduce the collective execution time of such tasks and improve resource utilization by considering the inter-task dependency and resource demands. Therefore, given the *n* multi-dependent tasks T_1, T_2, \dots, T_n , as shown in Fig. 2, EdgeDrones adopts the gang-scheduling [11, 12, 45] strategy and a variant bin-packing optimization to efficiently co-schedule and co-locate them in the attached edge nodes. We consider EdgeDrones as a Full Dependency and Full Packing (FDFP) scheduling approach. Thus, the scheduling time can be expressed as:

$$\sum_{z=1}^{m} \sum_{i=1}^{k_z} S_{ch_{z_i}} / k_z, \tag{1}$$

where *m* is the number of scheduling units, k_z is the number of tasks within the *z*-th scheduling unit having the tasks $\{T_{z_1}, T_{z_2}, \dots, T_{z_{k_z}}\}$.

We illustrate the advantage of the scheduling approach in EdgeDrones over three other existing schemes as follows; (i) a scheduling approach which does not consider tasks' dependencies, but schedules 50% of any given multi-dependent tasks by mainly focusing on task co-location, we refer to this scheduling approach as No Dependency and Full Packing

(NDFP), and it is similar to the approach in [46]; (ii) a scheduling approach which schedules up to 15% of any given multi-dependent tasks at a time, but does not consider task co-location, we refer to this scheduling approach as Partial Dependency and No Packing (PDNP), and it is similar to to the approach in [47]; (iii) a scheduling approach which schedules up to 40% of any given multi-dependent tasks with task co-location, we consider this scheduling approach as a Partial Dependency and Full Packing (PDFP), and it is similar to to the approach in [48]; and (iv) the Random approach, which does not consider both tasks' dependencies and task co-location, we refer to this scheduling approach as No Dependency and No Packing (NDNP). It is important to note that delay in scheduling inter-dependent tasks directly impacts their collective execution time. For the multidependent tasks of Fig. 1 with n = 12 tasks, Table 1 lists the scheduling orders and scheduling units for the schemes compared. EdgeDrones only needs one scheduling unit (m =1) which has $k_1 = 12$ tasks, and it also achieves the lowest execution time of $\frac{1}{12} \sum_{i=1}^{12} E_{ex_i}$. By contrast, Random has m = 12 scheduling units, each having a single task. Hence it has the highest execution time of $\sum_{i=1}^{12} E_{ex_i}$. Thus, Edge-Drones achieves the lowest scheduling and execution time. PDNP, PDFP and NDFP deploy individual or subsets of the tasks at a time. Generally, delay in scheduling dependent tasks directly impacts job completion time, and utilizing gang scheduling is beneficial for overall performance.

4. System Model, Problem Formulation and Algorithm Framework

In this section, we detail our proposed EdgeDrones for achieving optimal routing scheduling for drones aerial missions, high resource utilization and minimizing the execution times of applications deployed on an integrated edge computing system. The system model is shown in Fig. 3.

Scheme	Scheduling Order	Scheduling Units	
EdgeDrones	$\{T_1, T_2, T_3, T_4, T_5, T_6, T_7, T_8, T_9, T_{10}, T_{11}, T_{12}\}$	1	
PDNP	$T_3 \to T_2 \to \{T_1, T_4\} \to \{T_6, T_8\} \to \{T_5, T_7\} \to \{T_{10}, T_{11}\} \to \{T_{12}, T_9\}$	7	
PDFP	$T_1 \to T_2 \to T_3 \to \{T_4, T_5\} \to \{T_6, T_7, T_8, T_9\} \to \{T_{10}, T_{11}, T_{12}\}$	6	
NDFP	$\{T_1, T_2, T_3, T_4, T_5, T_6\} \to \{T_7, T_8, T_9, T_{10}, T_{11}, T_{12}\}$	2	
Random	$T_1 \rightarrow T_2 \rightarrow T_3 \rightarrow T_4 \rightarrow T_5 \rightarrow T_6 \rightarrow T_7 \rightarrow T_8 \rightarrow T_9 \rightarrow T_{10} \rightarrow T_{11} \rightarrow T_{12}$	12	

Table 1

Scheduling orders and units of various schemes

4.1. System Model

At the city center, we have a depot L_0 with coordinate $\{x_0, y_0\}$, where a fleet of drones $\mathbb{DR} = \{DR_1, \dots, DR_N\}$ are stationed and ready for aerial computing missions. Each drone has flight time availability f_i^{aval} and resource capacity $\mathcal{RC}_i^{(c,m)}$, in terms of CPU and memory resources. Then, at time t > 0, we have a set of end devices $\mathbb{D} = \{D_1, \dots, D_M\}$ requesting EC services across multi-location $\mathbb{L} = \{L_1, \dots, L_M\}$ with their coordinates $L_1 = \{x_1, y_1\}, \dots, L_M = \{x_M, y_M\}$ within the city. The edge federated system \mathbb{EF} consists of all service entities in the N drones DR_i , $1 \le i \le N$, and M end devices D_i , $1 \le i \le M$, i.e.,

$$\mathbb{EF} = \mathbb{DR} \bigcup \mathbb{D}.$$
 (2)

A set of multi-task $\mathbb{T} = \{T_1, \dots, T_N\}$ from the end device(s) at each location $L_i \in \mathbb{L}$ require an amount of CPU and memory resources $\mathcal{RR}_i^{\langle c,m \rangle}$ for execution. These resource requirements along with execution times are first estimated using linear regression ML model. The multi-task features $f_{\text{mt}}(\omega, \epsilon, \gamma)$, where ω is the number of instances, ϵ is type of tasks, γ is dependency depth, are fed into the model Θ^* to estimate the values of the resource requirement and execution times according to

$$\boldsymbol{f}_{\mathrm{mt}} \cdot \boldsymbol{\Theta}^{\star} = \left[\widetilde{E}_{ex_1} \widetilde{T}_1^{\langle c, m \rangle} \ \widetilde{E}_{ex_2} \widetilde{T}_2^{\langle c, m \rangle} \cdots \widetilde{E}_{ex_N} \widetilde{T}_N^{\langle c, m \rangle} \right], \ (3)$$

where $\widetilde{T}_i^{\langle c,m\rangle}$ and \widetilde{E}_{ex_i} are the estimated resource requirement (in terms of CPU and memory $\langle c,m\rangle$) and estimated execution time for task *i*, respectively. We show that with these estimated values, suitable drones can be assigned and multi-dependent tasks can be intelligently scheduled with the aim of minimizing their actual completion time, while maximizing available resources. Assuming that $f_{\rm mt} \in \mathbb{R}^{1 \times d}$

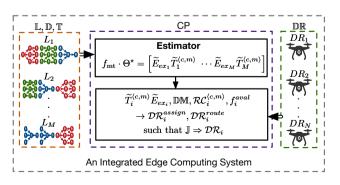


Figure 3: Orchestration overview of EdgeDrones.

is a *d*-dimensional vector (tensor), then Θ is a $(d \times \epsilon)$ dimensional parameter matrix. To build this predictor Θ , we train it using historical data from previously executed tasks/jobs based on Keras⁸. Keras is a library which wraps TensorFlow⁹ complexity into simple and user-friendly API. The dataset $DS = \{(\mathbf{x}_i, \mathbf{y}_i)\}_{i=1}^n$ contain *d*-dimensional tensors of data features $\mathbf{x}_i \in \mathbb{R}^{1 \times d}$ and ϵ -dimensional tensors of labels (actual execution times) $\mathbf{y}_i \in \mathbb{R}^{1 \times \epsilon}$. The learning problem is to solve the following optimization:

$$\Theta^{\star} = \arg\min_{\Theta \in \mathbb{R}^{d \times \epsilon}} \frac{1}{2n} \sum_{i=1}^{n} \|\boldsymbol{x}_{i} \Theta - \boldsymbol{y}_{i}\|_{2}^{2} + \frac{\lambda}{2} \|\Theta\|_{F}^{2}, \quad (4)$$

where λ is the regularization parameter and $\|\cdot\|_{F}$ denotes the Frobenius norm. The optimization (4) is solved using gradient-descent, where the model is updated iteratively until convergence, i.e., $\Theta^{t+1} = \Theta^t - \eta \left(\frac{1}{n} g(\Theta^t) + \lambda \Theta^l \right)$, in which η is the learning rate, $g(\Theta) = \frac{1}{n} X^{T} (X \Theta - Y)$ denotes the gradient of the loss function, $\boldsymbol{X} = \begin{bmatrix} \boldsymbol{x}_1^{\mathrm{T}} \cdots \boldsymbol{x}_n^{\mathrm{T}} \end{bmatrix}^{\mathrm{T}}$ and $Y = [y_1^T \cdots y_n^T]^T$ are the feature set and label set, respectively. These estimation values $\widetilde{T}^{\langle c,m\rangle}$ and \widetilde{E}_{ex} are important information for selecting suitable drones from the depot for the missions. The target location coordinates $L_i = \{x_i, y_i\}$ and depot coordinates $L_0 = \{x_0, y_0\}$ are used to compute the distance matrix to also aid with the selection of suitable drones. The distance matrix is an array of distances between these locations. These distances can be obtained using the Manhattan Distance¹⁰, in which the distance between two locations $L_i = \{x_i, y_i\}$ and $L_i = \{x_i, y_i\}$ is given as:

$$d^{L_{i,j}} = |x_i - x_j| + |y_i - y_j|.$$
(5)

These distances can also be obtained using Google Distance Matrix API¹¹. Hence, the distance matrix is given as;

$$\mathbb{DM} = \begin{bmatrix} d^{L_{0,0}} & d^{L_{0,1}} & d^{L_{0,2}} & \dots & d^{L_{0,n}} \\ d^{L_{1,0}} & d^{L_{1,1}} & d^{L_{1,2}} & \dots & d^{L_{1,n}} \\ \vdots & \vdots & \vdots & \ddots & \vdots \\ d^{L_{m,0}} & d^{L_{m,1}} & d^{L_{m,2}} & \dots & d^{L_{m,n}} \end{bmatrix}.$$
 (6)

The number of rows of the $\mathbb{D}M$ indicates the number of target locations and the depot inclusive. Autonomous drone

⁸https://keras.io/

9https://www.tensorflow.org/

10 https://en.wikipedia.org/wiki/Taxicab_geometry

¹¹https://cloud.google.com/blog/products/maps-platform/how-usedistance-matrix-api

systems have tools to help them estimate flight travel time for distances between locations [11, 12]. A drone's power has a capacity P, which is used for the flights (i.e., traveling and hovering) at a constant energy-efficient speed s. Note that the attached edge device(s) does not depend on the drone's power, instead, it is powered by its on-board power supply. For example, AWS Snowcone can be powered by any standard USB-C power bank and can deliver up to 14 trillion operations per seconds (TOPs) with as little as 10W of power. Therefore, the flight travel time for $d^{L_{i,j}}$ is given as $f^{L_{i,j}}$, and the flight hovering time at L_i is given as f^{hover} . As the actual execution time is unknown at this point, we assume that the flight hovering time is equivalent to the multi-task execution time estimation, i.e., $f^{hover} \approx$ $\sum_{i=1}^{N} \widetilde{E}_{ex_i}$. Suppose that a drone travels from L_i to L_j in one step of its route, and that:

- The drone's cumulative flight travel time upon arrival at L_i , given as f^{L_i} is 60s.
- The drone's cumulative flight travel time upon arrival at L_i , given as f^{L_j} is 130s.
- The drone's flight travel time $f^{L_{i,j}}$ is 50s.

Obviously, the drone can not depart location L_i immediately upon arrival, otherwise its cumulative flight travel time f^{L_j} upon arrival at L_i would be 110s. Instead, the drone must hover and execute the tasks at L_i for 20s before departing for L_i . In other words, the execution time also constitues a drone's flight travel time. On the other hand, note that the a drone's resource capacity $\mathcal{RC}^{(c,m)}$ does not accumulate as it travels along its route. This is because, after executing its tasks at L_i , the results are immediately and deterministically communicated back to the end device at L_i , and its resources becomes available for its next task execution at location L_i . Also, a drone can be assigned multiple disjointed locations as part of its mission, giving that it has sufficient resource availability, i.e., compute resources and flight time. Hence, a drone's total distance and flight travel time for its entire mission is given as;

$$d^{total} = \sum_{i=0}^{n} \sum_{j=i+1}^{m+1} d^{L_{i,j}}$$
(7)

and

$$f^{total} = \sum_{i=0}^{n} \sum_{j=i+1}^{m+1} f^{L_{i,j}} + f^{hover}_{j},$$
(8)

respectively. Since, a drone can be assigned multiple disjointed *m* locations, it must start and end its missions at the depot L_0 . For uniformity, we denote the starting and the ending depot location in its route as L_0 and L_{m+1} , respectively. Therefore, given a federated system \mathbb{EF} consisting of drone-enabled EC deployments \mathbb{DR} , where each participating drone \mathcal{DR}_i is attached with container-optimized nodes, and a set end devices \mathbb{D} requesting EC services across multiple locations in a city, an update state information from the CP which include each drones' fight time availability f_i^{aval} , its total resource capacity $\mathcal{RC}_i^{\langle c,m\rangle}$, each end device \mathcal{D}_i inter-dependent tasks execution and resource demand estimation $\widetilde{E}_{ex_i} \widetilde{T}_i^{\langle c,m\rangle}$, location coordinates $L_i = \{x_i, y_i\}$ and distance matrix \mathbb{DM} , is needed to select, assign and schedule optimal routes \mathcal{DR}_i^{route} for drones to visit these locations and execute the tasks, such that

$$\mathcal{DR}_{i}^{route} = \arg\min_{\mathcal{DR}_{i} \in \mathbb{DR}} \left\{ d_{i}^{total} : f_{i}^{total} < f_{i}^{aval}, \mathcal{RC}_{i}^{\langle c, m \rangle} \text{ sufficient} \right\}$$
(9)

A drone must execute a set of inter-dependent tasks \mathbb{T} at each of its assinged location before it returns to its depot. *EdgeDrones* utilizes the gang scheduling [45] strategy to co-schedule all tasks $T_i \in \mathbb{T}$ at a time, i.e.,

$$\mathbb{T} \Rightarrow \mathcal{DR}_i. \tag{10}$$

For a task $T_i \in \mathbb{T}$, its actual starting time and completion time are denoted as E_{st} and E_{cp} , respectively. Thus, its actual execution time is given as:

$$E_{ex} = E_{cp} - E_{st}.$$
 (11)

Hence the collective actual execution time of a multi (n)task T is given as $\sum_{i=1}^{n} \frac{E_{ex_i}}{n}$. Given a cluster of container-instances or nodes I_p in the edge device(s) attached to \mathcal{DR}_i , let $I_p^{\langle c,m \rangle}$ denote the *p*-th node's resource capacity. The estimated resource demands of k-dependent tasks to be orchestrated $\sum_{q=1}^{k} \widetilde{T}_{q}^{\langle c,m \rangle}$ and the resource capacity of each node is important information needed in order to make an efficient scheduling and co-location decision on I_n at time t. Our system extends to handle bulk requests from multiple end devices at the same location. Suppose at t, there are n service requests from multiple end devices at the same location L_i , where each device \mathcal{D}_i is offloading \mathbb{T} . The collective n requests from the end devices can be scheduled as a multi-Job \mathbb{J} , where $\mathbb{J} = \sum_{i=1}^{n} \mathbb{T}_{i}$, with collective resource demand estimation of each job denoted as $\sum_{q=1}^{k} \widetilde{T}_{q}^{\langle c,m \rangle} = \widetilde{T}^{\langle c,m \rangle'}$, and the aggregate execution time estimation of each job as $\sum_{q=1}^{k} \widetilde{E}_{ex_q} = \widetilde{E}_{ex_l}$. We can gang-schedule and co-locate \mathbb{J} effectively on DR_i :

$$\mathbb{J} \Rightarrow \mathcal{DR}_i,\tag{12}$$

by considering the estimated total resource demand of J:

$$\sum_{J \in \mathbb{J}} \widetilde{T}^{\langle c, m \rangle \prime} = \widetilde{T}^{\langle c, m \rangle \prime}_{\text{total}},$$
(13)

and DR_i resource capability $RC_i^{(c,m)}$. Hence, the total estimated execution time of \mathbb{J} at L_i is given as:

$$\sum_{J \in \mathbb{J}} \widetilde{E}_{ext} = \widetilde{E}_{ext}^{\text{total}}.$$
 (14)

Therefore, estimated resource utilization of DR_i for multijob \mathbb{J} at L_i is given by

$$\widetilde{\rho}_{D\mathcal{R}_{i}}^{\langle c,m\rangle} = \frac{\widetilde{T}_{\text{total}}^{\langle c,m\rangle'}}{\mathcal{R}C_{i}^{\langle c,m\rangle}}.$$
(15)

For a drone DR_i , let the aggregate of the actual execution time of multi-job J at L_i be

$$\sum_{J \in \mathbb{J}} \sum_{q=1}^{k} \frac{E_{ex_q}}{k} = \sum_{J \in \mathbb{J}} E_{ex'} = E_{ex'}^{\text{total}}, \quad (16)$$

and the total resources actually assigned for multi-job \mathbb{J} at L_i be $\mathcal{RA}_{DR_iU}^{\langle c,m\rangle}$. Under the condition that estimated total resource demand $\widetilde{T}_{\text{total}}^{\langle c,m\rangle'}$ is accurate, i.e., $\widetilde{T}_{\text{total}}^{\langle c,m\rangle'} \approx \mathcal{RA}_{DR_iU}^{\langle c,m\rangle}$, then $\mathcal{RA}_{DR_iU}^{\langle c,m\rangle}$ will not exceed $\mathcal{RC}_i^{\langle c,m\rangle}$. Similarly, under the condition that estimated total estimated execution time $\widetilde{E}_{exr}^{\text{total}}$ is accurate, i.e., $\widetilde{E}_{exr}^{\text{total}} \approx \mathcal{E}_{exr}^{\text{total}}$, then the drone \mathcal{DR}_i will have sufficient flight time availability f_i^{aval} for its entire mission.

Our learning-based approach has significant advantages over non-learning based counterparts. By accurately estimating the resource requirement and execution times of multi-tasks/multi-jobs, our scheme can intelligently select suitable drones having requisite resource and flight time availability for the missions, and co-locate multi-dependent tasks in their attached edge nodes, such that the dependent tasks can communicate and execute faster, ultimately to minimize the response times and improve resource utilization, hence guarantees the entire mission completion. The accuracy of the estimated resource requirement and execution times can be ensured by constructing multiple training datasets for different classes of multi-tasks/multijobs from historical data to learn multiple models, one for a class of multi-tasks/multi-jobs. Given the multi-tasks/multijobs to be deployed, the model that is most similar to them is employed to estimate the resource requirement and execution times. Since the estimated total resource demand $\widetilde{T}_{\text{total}}^{(c,m)\prime}$ and execution time $\widetilde{E}_{ex\prime}^{\text{total}}$ are accurate estimates of the actual total resource need to be allocated $\mathcal{RA}_{DR_{i}U}^{\langle c,m\rangle}$ and actual execution time $E_{ex\prime}^{\text{total}}$, it is unlikely that the selected drone edge DR_i will not have sufficient resources. In other words, it is very unlikely that

$$f_i^{\text{total}} > f_i^{aval} \text{ and/or } \mathcal{RA}_{D\mathcal{R}_i U}^{\langle c, m \rangle} > \mathcal{RC}_i^{\langle c, m \rangle},$$
(17)

which would lead to loss of job and jeopadize the entire mission. By contrast, standard non-learning based schemes have no means to intelligently select appropriate drones for ensuring that they will have sufficient resources, and the probability of (17) occurring can be much higher than our intelligent learning approach. There also exists simple and effective measure to guard against estimation error. It is obvious that loss of job may only occur in under estimation scenario. Instead of using the estimates of resource demand and execution time for selecting drones, we can add the two

Table 2						
Common not	ations					
Notation	Description					
EF	Federated edge deployments					
DR	A fleet of autonomous drones					
\mathbb{D}	A set of end devices					
L	A set of disjointed locations					
Т	A set of containerized inter-dependent applications					
DM	Distance matrix among the depot and target locations					
Т	Individual application or task					
$\widetilde{T}^{\langle c,m angle}$	Task resource requirements estimation					
$\widetilde{T}_{ ext{total}}^{\langle c,m angle \prime}$	Estimated total resource requirements for jobs					
I_i	Container-instance or node attached to a drone					
$I_i^{\langle c,m\rangle}$	Resource capacity or availability of a node					
$D\mathcal{R}_i$	Individual drone-enabled EC deployment					
$\mathcal{RC}_{i}^{\langle c,m\rangle}$	Resource capacity in a drone attached edge devices					
$\mathcal{RA}_{\mathcal{DR},U}^{\langle c,m \rangle}$	The total resources actually assigned for jobs					
DR_i^{assign}	A drone's set of assigned locations					
DR_{\cdot}^{route}	The route for a drone's mission					
$\mathcal{DR}_{iU}^{\langle c,m\rangle}$	Actual resources used for execution of jobs					
$\frac{D\mathcal{R}_{iU}^{\langle c,m\rangle}}{D\mathcal{R}_{iARU}^{\langle c,m\rangle}}$	The actual resource usage of a cluster					
$\frac{DR_{iARU}^{\langle c,m\rangle}}{\rho_{DR_i}^{\langle c,m\rangle}}$ $\frac{\overline{\rho_{DR_i}^{\langle c,m\rangle}}}{\rho_{DR_i}^{\langle c\rangle}, \rho_{DR_i}^{\langle m\rangle}}$	Actual resource utilization of jobs					
$\widetilde{\rho}_{DR_i}^{\langle c,m\rangle}$	Estimated resource utilization of jobs					
$\rho_{DR_i}^{\langle c \rangle}, \rho_{DR_i}^{\langle m \rangle}$	Actual cluster CPU, memory resource utilization					
$E_s t$, $E_c p$	Application/task starting, completion time					
E _{ex}	Application or task execution time					
$E_{ex'}^{\text{total}}$	Actual total execution time for jobs					
$\frac{E_{ext}^{\text{total}}}{\widetilde{E}_{ex}}$	Application or task execution time estimation					
$\widetilde{E}_{ex\prime}^{ m total}$	Estimated total execution time for jobs					
L_0, L_i	Drone's depot and destination location					
f_i^{aval}	Drone's flight time availability					
f_{i}^{hover}	Drone's hovering time at location					
d^{total}	Total distance of a drone's mission/trip					
f_i^{total}	Total flight travel time of a drone's mission/trip					
$d^{L_{i,j}}$	Travel distance of a drone from location i to j					
$f^{L_{i,j}}$	Flight travel time of a drone from location i to j					
ω_J	Number of instances of a Job					
ϵ_J	The type of job					
γ_J	Dependency depth of a job					
$m{f}_{ m mt}$	Set of multi-task runtime parameters					
Θ	Multi-output linear regression model					
J , $\mathbb J$	A Job, A set of Jobs					

standard deviations of the estimation to the corresponding estimates and use these 'modified' or 'overly' estimated values to select the drones. This will reduce the probability of (17) occurring to almost zero. It is straightforward to provide both the estimate and estimation standard deviation by dividing the training data into multiple subsets and running the estimation procedure multiple times to provide the average estimate and estimation standard deviation.

4.2. Problem Formulation

The notations adopted are listed in Table 2. Our Multi-Location Capacitated Mission Scheduling Problem (ML-CMSP) is summarized as: *Given a federated system* \mathbb{EF} *consisting of a fleet of autonomous drones* \mathcal{DR}_i , $1 \le i \le N$, and a set of end devices D_i , $1 \le i \le M$ with their computing activities to be performed across multi-location $\mathbb{L} = \{L_1, \dots, L_M\}$ within a city, the objectives are to co-schedule suitable drones, i.e., drones with sufficient flight time and computing resource availability, onto optimal mission routes among the locations, then co-locate the corresponding interdependent tasks on each drone's attached edge nodes at each location, such that all the activities are successful. Particularly, we present EdgeDrones, which intelligently co-schedules and co-locates multi-dependent tasks firmly on nodes, while considering task dependencies in oder to minimize the overall actual execution time and maximize the actual resource utilization, subject to certain constraints.

4.2.1. Constraints

At time t > 0, a set of end devices at different locations in the city require EC services, then a set of suitable drones are selected and co-schedule for EC missions. A drone can be assigned multiple disjointed locations as part of its mission, i.e., $D\mathcal{R}_i^{assign} = \{L_1, \dots, L_n\} \subseteq \mathbb{L}$. For a drone $D\mathcal{R}_i$, its trip route is given as:

$$\mathcal{DR}_{i}^{route} = \left(L_{0} \to L_{1} \to, \cdots, \to L_{n}, \to L_{0} \right), \quad (18)$$

and no location is assigned twice within a mission, i.e.,

$$\vartheta \left[\mathcal{DR}_{i}^{assign}, L_{i} \right] = \begin{cases} 1, & \text{if } L_{i} \text{ is assigned to } \mathcal{DR}_{i}, \\ 0, & \text{otherwise,} \end{cases}$$
(19)

such that;

$$\mathcal{DR}_{i}^{assign} \cap \mathcal{DR}_{k}^{assign} = \emptyset, \forall i, k.$$
(20)

This is to ensure that a location is mapped to just one drone, and that no location is assigned twice within the missions, where the indicator $\vartheta \left[\mathcal{DR}_i^{assign}, L_i \right] = 1$ indicates that location L_i is assigned to the drone \mathcal{DR}_i ; otherwise $\vartheta \left[\mathcal{DR}_i^{assign}, L_i \right] = 0$, all selected drones must start and end their missions at the depot L_0 , i.e.,

$$\sum_{i=0}^{n} \sum_{j=i+1}^{m+1} d^{L_{i,j}} - \sum_{j=0}^{m} \sum_{i=j+1}^{n+1} d^{L_{j,i}} = 0, \forall \mathcal{DR}_i \in \mathbb{DR}, \quad (21)$$

and they must complete all their tasks at any location L_i before departing for another location L_i , i.e.,

$$\forall \mathcal{DR}_i \in \mathbb{DR}, \ \varphi\left(\mathcal{DR}_i^{L_i \to L_j}\right) = \begin{cases} 1, \text{ if } \mathbb{J} \in [E_{cp}], \\ 0, \text{ if } \mathbb{J} \notin [E_{cp}], \end{cases} (22)$$

where the indicator $\varphi\left(\mathcal{DR}_{i}^{L_{i} \to L_{j}}\right) = 1$ indicates that the drone \mathcal{DR}_{i} has completed the execution of \mathbb{J} , and sent the results back to the devices at L_{i} . Hence it can depart L_{i} for L_{j} ; otherwise $\varphi\left(\mathcal{DR}_{i}^{L_{i} \to L_{j}}\right) = 0$. The collective resource

demand estimation of \mathbb{J} at any of the locations assigned to each drone, cannot exceed its resource capacity. Recall that a drone's resource capacity $\mathcal{RC}^{\langle c,m\rangle}$ does not accumulate as it travels along its route. This is because, after executing its tasks at any location L_i , the results are sent back to the device(s) at L_i , and its resources becomes available again for the jobs at its next location L_j . Since the actual total resources that needs to be assigned to the multi-job at L_i is unknown at the scheduling stage, we use the estimated total resource demand $\widetilde{T}_{total}^{\langle c,m\rangle'}$ to replace it:

$$\widetilde{T}_{\text{total}}^{\langle c,m\rangle'} \leq \mathcal{R}C_i^{\langle c,m\rangle}, \quad \forall \mathcal{D}\mathcal{R}_i \in \mathbb{D}\mathbb{R}.$$
(23)

Durng the multi-task scheduling, unused or inactive attached nodes $I_i \in DR_i$ in a selected drone would be shut down. All the nodes can be expressed in one of these two states: *Active* and *Inactive*. An *Active* node is a node that is running and is currently considered for allocation or has at least a job being started, executing or completing. An *Inactive* node is a node that is not running and and is not currently considered for allocation and not having at least a job that is being started, executing or completing. These two states can be expressed as follows:

$$\forall c, m \beta (I_i) = \begin{cases} 1, & Active \text{ if } J_i \in [E_{st}, E_{cp}, E_{ex}], \\ 0, & Inactive \text{ if } J_i \notin [E_{st}, E_{cp}, E_{ex}], \end{cases}$$
(24)

where the indicator $\beta(I_i) = 1$ indicates that the node I_i is ready to accept new jobs, and at least a job J_i is being started, executing or completing, i.e., $J_i \in [E_{st}, E_{cp}, E_{ex}]$, on I_i ; otherwise $\beta(I_i) = 0$.

The aggregate actual execution time of \mathbb{J} at all locations assinged to each drone $\sum_{L=i}^{n} E_{ex_i'}^{\text{total}}$ and the total flight travel time $\sum_{i=0}^{n} \sum_{j=i+1}^{m+1} f^{L_{i,j}}$ cannot exceed its flight time availability f_i^{aval} . Since the aggregate actual execution time $\sum_{L=i}^{n} E_{ex_i'}^{\text{total}}$ is unavailable at the scheduling stage, we replace it with the estimate $\sum_{L=i}^{n} \widetilde{E}_{ex_i'}^{\text{total}} \approx f^{hover'}$:

$$f_i^{total} = \sum_{i=0}^n \sum_{j=i+1}^{m+1} f^{L_{i,j}} + f_j^{hover'},$$
(25)

therefore,

$$f_i^{total} \le f_i^{aval}, \quad \forall f_i^{aval} \in \mathbb{DR}.$$
 (26)

4.2.2. Optimization formulation

Suitable drones are selected and co-schedule with the least total distance to visit and execute tasks at target locations:

Minimize
$$d_i^{total} \quad \forall \mathcal{DR}_i \in \mathbb{DR},$$
 (27)

subject to
$$\vartheta \left[\mathcal{DR}_{i}^{assign}, L_{i} \right] \in \{0, 1\}, \exists, (28)$$

$$D\mathcal{R}_{i}^{assign} \cap D\mathcal{R}_{k}^{assign} = \emptyset, \forall i, k,$$
 (29)

EdgeDrones: Co-Scheduling of Drones for Multi-Location Aerial Computing Missions

5

$$\sum_{i=0}^{n} \sum_{j=i+1}^{m+1} d^{L_{i,j}} - \sum_{j=0}^{m} \sum_{i=j+1}^{n+1} d^{L_{j,i}} = 0, \quad (30)$$

$$\varphi\left(\mathcal{DR}_{i}^{L_{i} \to L_{j}}\right) \in \{0, 1\}, \forall \mathcal{DR}_{i} \in \mathbb{DR}, \exists, (31)$$

$$\widetilde{T}_{\text{total}}^{\langle c,m\rangle\prime} \leq \mathcal{R}C_i^{\langle c,m\rangle}, \,\forall \mathcal{D}\mathcal{R}_i \in \mathbb{D}\mathbb{R}, \quad (32)$$

$$f_i^{\text{total}} \le f_i^{aval}, \,\forall f_i^{aval} \in \mathbb{DR}.$$
(33)

The objective function (27) is to minimize the total distance $d^{total} = \sum_{i=0}^{n} \sum_{j=i+1}^{m+1} d^{L_{i,j}}$ of each drone's route to visit its assigned locations. Constraints (28) to (30) and condition (31) ensure that no location is assigned twice within the missions, all drones must start and finish their trip at the depot, and each drone must execute all its task at any location before it departs for another location. Constraint (32) guarantees that $\widetilde{T}_{total}^{(c,m)\prime}$ of \mathbb{J} at any assigned location for DR_i would not exceed its resource capacity $\mathcal{RC}_i^{(c,m)}$. Constraint (33) also guarantees that each drones' total flight travel time f_i^{total} for any mission would not exceed its flight time availability f_i^{aval} . The details of our optimal mission and route planning is given in Section 4.3 and in Algorithm 2.

As the actual resource utilization of a cluster/edge is unknown, we maximize the estimated resource utilization:

Maximize
$$\tilde{\rho}_{DR_i}^{(c,m)}$$
, (34)

subject to
$$\mathbb{J} \Rightarrow D\mathcal{R}_i, \exists,$$
 (35)

$$\widetilde{T}_{\text{total}}^{\langle c,m\rangle\prime} \leq \mathcal{RC}_{i}^{\langle c,m\rangle}, \ \forall \mathcal{DR}_{i} \in \mathbb{DR}, \quad (36)$$

$$f_i^{\text{total}} \le f_i^{aval}, \,\forall f_i^{aval} \in \mathbb{DR}, \tag{37}$$

$$\varphi\left(\mathcal{DR}_{i}^{L_{i} \to L_{j}}\right) \in \{0, 1\}, \forall \mathcal{DR}_{i} \in \mathbb{DR}, \exists$$
(38)

$$\beta\left(I_{i}\right) \in \{0,1\}, \ \forall c, m, \ \exists. \tag{39}$$

Provided that the estimated resource utilization $\tilde{\rho}_{DR_i}^{(c,m)}$ is accurate, little optimality will be lost. The constraints (35) to (37) indicate the dispatching of multi-job J at each assigned location to DR_i , given that its resources and flight time availability is sufficient. More specifically, (35) is the multijob J deployment constraint, guaranteeing that J is gangscheduled onto \mathcal{DR}_i attached resources, such that dependent tasks within each $J \in \mathbb{J}$ can communicate and execute faster. The constraint (36) guarantees that $\widetilde{T}_{\text{total}}^{(c,m)'}$ of \mathbb{J} would not exceed $\mathcal{RC}_i^{\langle c,m\rangle}$ of \mathcal{DR}_i , and constraint (37) guarantees that f_i^{total} would not exceed f_i^{aval} of any $\mathcal{DR}_i \in \mathbb{DR}$. The condition (38) guarantees that active nodes $(\beta (I_i) = 1)$ would be used for execution, and inactive nodes $(\beta(I_i)=0)$ would be shut down. Hence, our aim is to minimize the number of active nodes used for execution by co-locating jobs tightly on each node in order to maximize resource utilization. We shall discuss the details of our multi-job colocation principle in Section 4.3 and Algorithm 3.

Then again, $\widetilde{E}_{ex'}^{\text{total}}$ of \mathbb{J} can be minimized depending on scheduling:

Minimize
$$\widetilde{E}_{ext}^{\text{total}}$$
, (40)

subject to
$$\mathbb{J} \Rightarrow \mathcal{DR}_{i^{\star}}, \exists,$$
 (41)

$$\widetilde{T}_{\text{total}}^{\langle c,m\rangle\prime} \le \mathcal{RC}_{i}^{\langle c,m\rangle}, \ \forall \mathcal{DR}_{i} \in \mathbb{DR}, \quad (42)$$

$$f_i^{\text{total}} \le f_i^{aval}, \,\forall f_i^{aval} \in \mathbb{DR},$$
 (43)

$$\varphi\left(\mathcal{DR}_{i}^{L_{i} \to L_{j}}\right) \in \{0, 1\}, \forall \mathcal{DR}_{i} \in \mathbb{DR}, \exists$$

$$(44)$$

Note that the actual overall execution time E_{ext}^{total} is unknown at this stage, and we use the estimated overall execution $\widetilde{E}_{ext}^{\text{total}}$ to replace it in the optimization. Again, provided that the estimate $\widetilde{E}_{ext}^{\text{total}}$ is accurate, little optimality will be lost. The constraint (41) guarantees that \mathbb{J} at any location is dispatched to \mathcal{DR}_i , such that dependent tasks within each $J \in \mathbb{J}$ can communicate and execute faster. The constraint (42) guarantees that $\widetilde{T}_{\text{total}}^{(c,m)'}$ of \mathbb{J} would not exceed $\mathcal{RC}_i^{(c,m)}$ of \mathcal{DR}_i , and constraint (43) guarantees that f_i^{total} would not exceed f_i^{aval} of any $\mathcal{DR}_i \in \mathbb{DR}$. The condition (44) ensures that each drone must execute all its task at any location before it departs for another location.

4.3. EdgeDrones Algorithm Framework

Our EdgeDrones approach consists of linear regression estimation, optimal route planning for multi-drone, and gang scheduling of tasks. These three components aim at providing optimal solution for our Multi-Location Capacitated Mission Scheduling Problem (MLCMSP). Particularly, the optimization (27), (34) and (40) aim at ensuring least travel distance for drones to visit their assigned locations, and every task at each location is fast executed given the available resources, such that the missions are accomplished. The values of the linear regression estimations are required by the router, as well as the update state of the drones' flight time availabity for effective route planning and assignment, while our gang-scheduling approach involves co-scheduling and co-locating tasks firmly on available resources. We detail the procedures of the three components of EdgeDrones as follows:

4.3.1. Resource and execution time estimation

Algorithm 1 describes the resource and execution time estimations for multi-job. As \mathbb{J} are released, their collective resource requirement $\widetilde{T}_{total}^{\langle c,m\rangle\prime}$ and execution time $\widetilde{E}_{ex\prime}^{total}$ are estimated. The set of runtime parameters $f_{mt}(\omega, \epsilon, \gamma)$, where ω is the number of instances, ϵ is type of tasks , γ is dependency depth, are fed into the model Θ^* to to produce the estimation values (line $2 \sim 9$). Once the estimation values are produced, they are used in the assignment and route planning.

4.3.2. Mission planning

Given the update state from the CP, i.e., all ready drones $DR_i \in D\mathbb{R}$ at the depot L_0 (which include each drones'

Algorithm 1 Linear Regression Estimation				
Input : At $t > 0$, \mathbb{J} at L_i ; f_{mt} are fed into Θ^*				
Output: $\widetilde{T}_{\text{total}}^{\langle c,m \rangle \prime}$ and $\widetilde{E}_{ex\prime}^{\text{total}}$				
1: for $J_i \in \mathbb{J}$ do				
2: Type of Job $J = \epsilon_J$				
3: Number of instances of Job $J = \omega_J$				
4: Dependency depth of Job $J = \gamma_J$				
5: for $T_i \in J_i$ do				
6: $f_{\mathrm{mt}}(\omega,\epsilon,\gamma)\cdot\Theta^{\star} = \left[\widetilde{T}_{i}^{\langle c,m\rangle} \widetilde{E}_{ex_{i}}\right]$				
7: end for				
8: $\widetilde{T}_{i}^{\langle c,m\rangle\prime} = \widetilde{T}_{i}^{\langle c,m\rangle\prime} + \widetilde{T}_{i}^{\langle c,m\rangle}$				
8: $\widetilde{T}_{i}^{\langle c,m\rangle\prime} = \widetilde{T}_{i}^{\langle c,m\rangle\prime} + \widetilde{T}_{i}^{\langle c,m\rangle}$ 9: $\widetilde{E}_{ex_{i}\prime} = \widetilde{E}_{ex_{i}\prime} + \widetilde{E}_{ex_{i}}$				
10: end for				

Algorithm 2 Mission Planning

Input: At t > 0; $L_i \in \mathbb{L}$; $\widetilde{T}_{total}^{\langle c,m \rangle \prime} \in L_i$; $\widetilde{E}_{ex\prime}^{total} \in L_i$; $\mathcal{DR}_i \in \mathbb{DR}$; $f^{aval} \in DR_i; \mathcal{R}C^{\langle c,m \rangle} \in DR_i; \text{ and } \mathbb{DM}$ **Output**: $\mathcal{DR}_{i}^{assign}$ and \mathcal{DR}_{i}^{route} 1: for $\mathcal{DR}_i \in \mathbb{DR}$ do Initialize total flight travel time f_i^{total} for \mathcal{DR}_i 2: Initialize total travel distance d_i^{total} for \mathcal{DR}_i 3: f_i^{aval} = Flight travel time availability of \mathcal{DR}_i 4: $\mathcal{RC}_{i}^{\langle c,m\rangle}$ = Resource capacity of \mathcal{DR}_{i} 5: \mathbb{L} = Target locations 6: L_i = route start for \mathcal{DR}_i 7: while $L_i \neq$ end of route do $\widetilde{T}_{\text{total}}^{\langle c,m \rangle \prime} =$ Resource demand at L_i 8. 9: $\begin{aligned} \widetilde{F}_{ext}^{total} &= \text{Execution time estimation at } L_i \\ \widetilde{F}_{ext}^{total} &= d_i^{total} + \widetilde{E}_{ext}^{total} \\ \text{if } f_i^{total} &\leq f_i^{aval}; \ \vartheta \left[\mathcal{DR}_i^{assign}, L_i \right] = 0; \text{ and} \end{aligned}$ 10: 11: 12: $\widetilde{T}_{\text{total}}^{\langle c,m\rangle\prime} \leq \mathcal{R}C_{i}^{\langle c,m\rangle} \text{ then}$ $\vartheta \left[\mathcal{D}\mathcal{R}_{i}^{assign}, L_{i} \right] = 1$ $L_{j} = \text{next feasible neighbor}$ $L_{i} = L_{j}$ 13: 14: 15: else 16: $L_j = \text{next feasible neighbor}$ 17: $L_i = L_i$ 18: 19: end if $d_i^{total} = d_i^{total} + d^{L_{i,j}}$ 20: 21: end while 22: end for

flight time availability f_i^{aval} and resource capacity $\mathcal{R}C_i^{\langle c,m\rangle}$, all end devices at target locations $L_i \in \mathbb{L}$ (which include each device set of inter-dependent tasks resource requirement $\widetilde{T}_{total}^{\langle c,m\rangle'}$ and execution time $\widetilde{E}_{ex\prime}^{total}$ estimation), and the distance matrix DM, such that starting from L_0 , a route is iteratively built and assigned to a drone $\mathcal{D}\mathcal{R}_i$, by selecting from among the nearest locations which meet the constraints, i.e., L_1 whose its end device(s) $\widetilde{T}_{total}^{\langle c,m\rangle'}$ and $\widetilde{E}_{ex\prime}^{total}$ does not exceed the drone's $\mathcal{R}C_i^{\langle c,m\rangle}$ and f_i^{aval} , respectively, and

Algorithm 3 Multi-job Co-location

Input: \mathbb{J} gang-scheduled onto $\mathcal{DR}_{i_{\perp}}$, resource demand estimation $\sum_{I \in \mathbb{J}} \widetilde{T}_{I}^{\langle c,m \rangle'}$, resource availability $I_{i}^{\langle c,m \rangle}$ of all nodes $I_i \in \mathcal{DR}_i$ **Output**: \mathbb{J} is co-located, such that **Minimize** $\sum_{I_i \in D\mathcal{R}_i} I_i$ 1: for $I_i \in \mathcal{DR}_{i_{\perp}}$ do if $\beta(I_i) = 1$ then 2: $I_i^{\langle c,m\rangle} = \langle c,m\rangle$, i.e., initial resource available 3: for $J \in \mathbb{J}$ do 4: if $\Gamma[J, I_i] = 0$ and $\widetilde{T}_J^{\langle c, m \rangle \prime} \leq I_i^{\langle c, m \rangle}$ then 5:
$$\begin{split} & I_{i} \rightarrow I_{i} \\ & J \rightarrow I_{i} \\ & \Gamma \left[J, I_{i} \right] = 1 \\ & I_{i}^{\langle c, m \rangle} = I_{i}^{\langle c, m \rangle} - \widetilde{T}_{J}^{\langle c, m \rangle \prime} \end{split}$$
6: 7. 8: end if if $I_i^{\langle c,m\rangle}$ close to zero then 9: 10: break 11: 12: end if end for 13: 14: end if 15: end for

whose flight travel time f_1^{total} does not exceed the drone's f_i^{aval} . This process resumes from L_1 to find L_2 , and so on until there is no feasible neighbor (line $6 \sim 17$). L_0 is finally added to conclude the route. This procedure is repeated to find other routes until all the possible target locations are chosen. This mission planning also aims to optimized and minimize the distance travelled by each drone, which lets us find an optimal solution from the given Algorithm 2. To solve this mission scheduling problem, we have adopted the CP-SAT solver and the MPSolver wrapper.

4.3.3. Co-location

In the edge resources attached to each drone DR_i , our co-location algorithm uses the $I_i^{\langle c,m\rangle}$ and $\widetilde{T}_i^{\langle c,m\rangle\prime}$ of each $J_i \in \mathbb{J}$ to provide efficient co-location, such that fewer nodes are used for execution at each location. Specifically, the gang scheduling approach is adopted alongside our bin-packing optimization to co-schedule and co-locate J at a time. Binpacking is one the of the most popular packing problems. The goal is to minimize the number of nodes used as given in optimization (45). Unlike other approaches, such as first fit bin packing problem (FFBPP) [49], it requires the next J_i to be placed on the active node, otherwise, it is placed on a new node. Our approach scans all $J \in \mathbb{J}$ and maps J_i to active nodes in full utilization. All $J \in \mathbb{J}$ are co-located firmly on active nodes, so that resource wastage is avoided and fewer nodes are used to execute all jobs concurrently. Hence our co-location strategy is to find the solution to the problem:

$$\mathbf{Minimize} \sum_{I_i \in \mathcal{DR}_i} I_i, \tag{45}$$

subject to
$$\mathbb{J} \Rightarrow \mathcal{DR}_{i_{+}}, \exists,$$
 (46)

Table 3

Multi-job across multiple disjointed locations in a city, where the actual resource consumed for multi-job execution $T_{\text{total}}^{(c,m)'}$ and the actual execution time E_{ext}^{total} are taken from the original Alibaba data, while the estimated resource demand $\tilde{T}_{\text{total}}^{(c,m)'}$ and execution time $\tilde{E}_{ext}^{\text{total}}$ are calculated by Algorithm 1

L	J	T	$\widetilde{T}_{ ext{total}}^{\langle c,m angle \prime}$	$T_{\text{total}}^{\langle c,m angle\prime}$	NAEE	$\widetilde{E}_{ex'}^{\text{total}}(s)$	$E_{ex'}^{\text{total}}(s)$	NAEE
L_1	2	10	⟨570.18, 1.98⟩	⟨595, 1.92⟩	(0.04, 0.03)	173.46	148	0.17
L_2	3 12 (625.06, 2.37) (540, 1.85)		(0.15, 0.28)	189.03	164	0.15		
L_3	2	9	(478.02, 1.73)	⟨340, 0.96⟩	$\langle 0.4, 0.8 \rangle$	167.8	142	0.18
L_4	$\begin{array}{c c c c c c c c c c c c c c c c c c c $		(0.1, 0.2)	56.69	44	0.28		
L_5	5	5 21 (1135.11, 4.13) (1035, 3.38)		(1035, 3.38)	(0.09, 0.22)	355.68	309	0.15
L_6	5	23	(1228.72, 4.56)	(1080, 3.4)	(0.13, 0.3)	370.27	311	0.19
L_7	5	19	(1005.85, 3.89)	(1070, 3.39)	$\langle 0.05, 0.14 \rangle$	236.95	198	0.19
L_8	3	15	⟨773.7, 3.02⟩	(655, 2.13)	$\langle 0.18, 0.4 \rangle$	211.06	178	0.18
L_9	2	10	⟨570.18, 1.98⟩	⟨595, 1.92⟩	(0.04, 0.03)	173.46	148	0.17
L_{10}	$3 14 \langle 727.81, 2.8 \rangle$		(670, 2.1)	(0.08, 0.3)	202.45	172	0.17	
L_{11}	3	15	⟨773.7, 3.02⟩	(655, 2.13)	$\langle 0.18, 0.4 \rangle$	211.06	178	0.18
L_{12}	4	17	(876.45, 3.44)	(785, 2.38)	(0.11, 0.4)	224.49	188	0.19
L_{13}	5	19	(1025.99, 3.7)	(885, 2.88)	(0.15, 0.28)	341.79	298	0.14
L_{14}	3	17	(925.32, 3.48)	(990, 3.14)	(0.06, 0.1)	341.79	182	0.87
L_{15}	3	14	(727.81, 2.8)	(670, 2.1)	(0.08, 0.33)	202.45	172	0.17
L_{16}	4	15	⟨773.7, 3.02⟩	(655, 2.13)	$\langle 0.18, 0.4 \rangle$	202.45	178	0.13

$$\sum_{J \in \mathbb{J}} \Gamma \left[J, \ I_i \right] \cdot \widetilde{T}_J^{\langle c, m \rangle \prime} \le I_i^{\langle c, m \rangle}, \quad \forall c, m,$$

$$\tag{47}$$

where

$$\Gamma \begin{bmatrix} J, & I_i \end{bmatrix} = \begin{cases} 1, & \text{if } J \Rightarrow I_i, \\ 0, & \text{otherwise.} \end{cases}$$
(48)

The constraint (46) is the multi-job J deployment constraint, guaranteeing that \mathbb{J} is gang-scheduled to $D\mathcal{R}_{i}$, such that dependent tasks within each $J \in \mathbb{J}$ can communicate and execute faster. The constraint (47) indicates that the total estimated resource requirements of co-located jobs $\sum_{i=1}^{N} \widetilde{T}_{i}^{\langle c,m \rangle i}$ cannot exceed $I_i^{\langle c,m\rangle}$, the node resource availability. The condition (48) means that if job J_i is placed on the node I_i , then $\Gamma[J_i, I_i] = 1$; otherwise, $\Gamma[J_i, I_i] = 0$. This is to guarantee that each $J \in \mathbb{J}$ is placed in exactly one node. To solve this multi-job packing problem, we have adopted the solving Constraint Integer Programs (SCIP) solver, which is currently one of the fastest mathematical programming (MP) solvers for this problem [12]. Algorithm 3 describes the co-location strategy which co-locates multi-dependent tasks firmly on nodes, such that for any given jobs, resource wastage is avoided and fewer nodes are used for execution. It takes the resource demand estimation of multi-task/job and resource availability of nodes as input, then scans all $J \in \mathbb{J}$ and maps them to active nodes in full utilization (line $2 \sim 7$).

4.3.4. Connection with optimization objectives

As stated previously, our objectives are to minimize the number of selected drones and total distance travelled by each drone, maximize the actual edge cluster resource utilization, and to minimize the overall actual execution time of the task-dependent multi-jobs. Algorithms 1, 2 and 3 together achieve these objectives. By gang-dispatching the task-dependent multi-jobs to an edge having the sufficient resource for the jobs and flight time availability, Algorithm 2 ensures that drones assigned for missions allocates the sufficient actual resources needed for jobs execution $\mathcal{DR}_{iU}^{\langle c,m\rangle}$, such that the dependent tasks can be executed faster, ultimately leading to a smaller actual aggregate execution time E_{ext}^{total} and better actual cluster resource utilization. By intelligently packing dependent tasks tightly on nodes, Algorithm 3 is capable of fully utilizing available resources at edge clusters, ultimately leading to the actual resource assigned to the execution of jobs $\mathcal{DR}_{iU}^{\langle c,m\rangle}$ as small as possible while guaranteeing it is sufficient for the multi-jobs. More specifically, the actual resource usage (ARU) of the cluster for multi-job J deployment is given by

$$D\mathcal{R}_{iARU}^{\langle c,m\rangle} = \frac{D\mathcal{R}_{iU}^{\langle c,m\rangle}}{\mathcal{R}C_i^{\langle c,m\rangle}}.$$
(49)

It can be seen that solving the optimization (45) is directly linked to minimize the ARU (49). Let the actual CPU resource and the actual memory resource assigned for \mathbb{J} be $\mathcal{DR}_{iU}^{\langle c \rangle}$ and $\mathcal{DR}_{iU}^{\langle m \rangle}$, respectively. Further denote the actual CPU consumed and the actual memory consumed in executing \mathbb{J} as $\sum_{J \in \mathbb{J}} T^{\langle c \rangle'}$ and $\sum_{J \in \mathbb{J}} T^{\langle m \rangle'}$, respectively. Then the actual CPU utilization $\rho_{DR_i}^{\langle c \rangle}$ and the actual memory utilization $\rho_{DR_i}^{\langle m \rangle}$ are defined respectively by

$$\rho_{DR_{i}}^{\langle c \rangle} = \frac{\sum_{J \in \mathbb{J}} T^{\langle c \rangle'}}{DR_{iU}^{\langle c \rangle}},\tag{50}$$

$$\rho_{DR_{i}}^{\langle m \rangle} = \frac{\sum_{J \in \mathbb{J}} \mathcal{T}^{\langle m \rangle \prime}}{DR_{iU}^{\langle m \rangle}}.$$
(51)

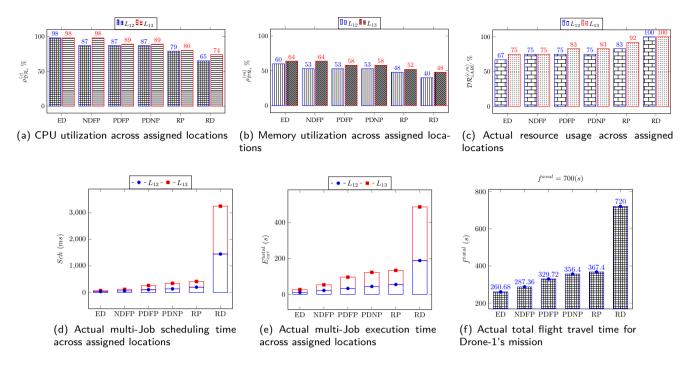


Figure 4: Activities and utilities of Drone-1's mission.

Algorithms 2 and 3 are directly connected with minimizing d^{total} , minimizing E_{ext}^{total} as well as maximizing $\rho_{DR}^{\langle c \rangle}$ and maximizing $\rho_{DR}^{\langle m \rangle}$.

5. Performance Evaluation

In this section, we described our experimental setup including cluster resource configuration, the Alibaba cluster data traces used, and the comparison baselines. We perform extensive experiments to compare EdgeDrones against some existing schemes. We will also compare the performance of EdgeDrones against exisiting schemes in individual drones. We show that EdgeDrones can achieve minimized actual execution time of multi-dependent tasks, achieve high resource utilization, achieve load balancing, use fewer cluster resources and avoid loss of job in an aerial edge computing system.

5.1. Experimental Setup

Drone's Mission Scheduling and Resources: Our ML-CMSP is implemented using Google OR Tools¹². It uses MPSolver wrapper for solving LP and MIP problems. We perform experiments for a set of drones with computing resource and flight time constraints, among a set of 16 target locations in a city. At each of these locations are end devices, with multi-dependent tasks/jobs needed to be executed. Our linear regression ML model, as given in Algorithm 1, estimates the resource demands and execution time of tasks at each location (as shown in Table 3).

Each of the drones has different payload capacity. The payload is the *weight* a drone can carry in the air. The total weight of the payload has a great impact on the flight time of the drone. For example, the Aurelia X8 Standard drone¹³ with a payload of 8 kg, has a net flight time of 25 min. However, with a maximum payload capacity, the maximum flight time availability of the drone will be 12 min. It has a maximum flight speed *s* of 15 m/s and a maximum wind resistance speed of 9 m/s. Also, for the Aurelia X6 Pro drone, its net flight time is 55 min, however, with its full payload of 5 kg, it will have a 30 min flight time availability. Hence, the flight time availability of the selected drones, as well as their resource capacities is given in Table 4.

The optimal routes of the selected drones is also given in Table 4. These optimal routes are derived using Algorithm 2. The entire missions covered a distance of 6552 meters, with total flight travel time of 3659 seconds, given that each drone is traveling at a constant flight speed *s* of 13 m/s.

Multi-dependent Tasks: We employ the v-2018 version of Alibaba cluster trace¹⁴, which records the activities of about 4000 machines in a perids of 8 days. The entire trace contains more than 14 million tasks with more than 12 million dependencies, and more than 4 million jobs. Among which we have deployed 54 jobs with total of 238 tasks (including dependencies) for our experiments. The number of tasks within each job ranges from (1, 5], while the task dependency depth among the jobs ranges from (1, 4]. The multi-task dependencies in in the data trace is valuable for our investigation. Researchers have thoroughly investigated

¹² https://developers.google.com/optimization

¹³https://aurelia-aerospace.com/our-drones/ ¹⁴https://github.com/alibaba/clusterdata

U. Awada et al.: Preprint submitted to Elsevier

DR_i	\mathcal{DR}_{i}^{route}	$d_i^{total}(m)$	$f_i^{hover'}(s)$	Attached Edge Devices and total weight	$\mathcal{RC}_{i}^{\langle c \rangle}$	$\mathcal{RC}_{i}^{\langle m \rangle}$	$f_i^{aval}(s)$				
DR_1	$\{L_0 \rightarrow L_{12} \rightarrow L_{13} \rightarrow L_0\}$	936	566	Huawei AR502H Series x3 = 3.3 kg	12 Cores	6 GiB	600				
DR_2	$\{L_0 \to L_9 \to L_{14} \to L_{16} \to L_8 \to L_7 \to L_0\}$	1712	1165	HIVECELL $x_2 = 2.72 \text{kg}$	12 Cores	16 GiB	1200				
DR_3	$\{L_0 \rightarrow L_1 \rightarrow L_4 \rightarrow L_3 \rightarrow L_{15} \rightarrow L_{11} \rightarrow L_0\}$	2192	811	HIVECELL + Huawei AR502H Series $= 2.46 \mathrm{kg}$	10 Cores	8 GiB	900				
DR_4	$\{L_0 \rightarrow L_{10} \rightarrow L_2 \rightarrow L_6 \rightarrow L_5 \rightarrow L_0\}$	1712	1117	Azure Stack Edge mini $= 3.17 \mathrm{kg}$	16 Cores	48 GiB	1200				

Drones assigned locations, optimal routes and resources

Table 4

v-2018 version of Alibaba cluster trace and used it for various task scheduling problems [11, 12, 50, 51].

Comparison Baselines: We compare the scheduling approach of EdgeDrones **(ED)** with the following three existing schemes and the random approach, fixing each drone's routes to that of EdgeDrones, as follows:

- 1. An approach which does not consider tasks' dependencies, but schedules 50% of any given multidependent tasks by mainly focusing on task co-location. We refer to this approach as No Dependency and Full Packing (**NDFP**), and it is similar to the approach in [46].
- 2. An approach which schedules up to 40% of any given multi-dependent tasks with task co-location. We consider this approach as a Partial Dependency and Full Packing (**PDFP**), and it is similar to to the approach in [48].
- 3. An approach which schedules up to 15% of any given multi-dependent tasks at a time, but does not consider task co-location. We refer to this approach as Partial Dependency and No Packing (**PDNP**), and it is similar to to the approach in [47].
- 4. An approach which schedules tasks according to their resource requests for execution, i.e., the more resource demand, the higher priority for the task to be scheduled and allocate resources. We refer to this strategy as Resource Priority (**RP**), and it is similar to the approach in [52].
- 5. Random (**RD**) approach schedules a single task individually and assumes a node can only execute a task at a time.

5.2. Deployment Results and Performance Comparison

Our investigation focuses on CPU and memory usage/utilization, task deployment, scheduling time, execution time and successful mission completion. The results obtained by ED, NDFP, PDFP, PDNP, RP and RD are compared.

5.2.1. Resource and execution time estimation accuracy

As detailed in the previous section, to implement the proposed learning based intelligent drone routing and multitask co-location strategy, we train a linear regression model from a training dataset. In the real-time application experiments, the trained model is used to estimate the resource requirement and execution time (Algorithm 1). The estimated resource requirement and execution time are then employed to aid our optimal route planning and intelligent multi-task scheduling strategy (Algorithms 2 and 3). Clearly, the accuracy of Algorithm 1 impacts the achievable performance of our EdgeDrones. Therefore, we first investigate the accuracy of our trained linear regression model.

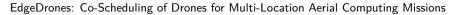
The multi-job execution information across the federated deployments, obtained according to Alibaba data, are listed in Table 3, where the estimated resource demand $\widetilde{T}_{total}^{\langle c,m\rangle'}$ and the estimated execution time $\widetilde{E}_{ext}^{total}$ are calculated using Algorithm 1, while the actual resource consumed for the multi-job execution $T_{total}^{\langle c,m\rangle'}$ and the actual execution time E_{ext}^{total} are taken from the original data. The normalized absolute estimate error (NAEE), defined as

NAEE =
$$\frac{\left|\text{estimated value} - \text{actual value}\right|}{\text{actual value}}$$
, (52)

is also listed in Table 4 for both resource consumed and execution time, which serves as the estimation accuracy measure for the trained multi-output linear regression model. The average NAEE across 16 locations is 0.13 for CPU resource, 0.28 for memory resource, and 0.22 for execution time. From Tables 3 and 4, it can be seen that $\widetilde{T}_{total}^{(c,m)\prime} < D\mathcal{R}_i^{(c,m)}$ and $f^{total} < f^{aval} \forall D\mathcal{R}_i$, given that each drone is traveling at a constant flight speed *s* of 13 m/s. In other words, each drone has sufficient resource to execute its multi-jobs assigned to it. This further indicates the suitability or accuracy of the trained ML model to provide the necessary information for our intelligent co-location strategy.

5.2.2. Performance comparison across integrated edge resources

After completing the optimal route planning, as shown in Table 4, we are now ready to co-schedule the drones for their missions, apply our EdgeDrones to orchestrate 54 jobs with 238 tasks among the four drones and compare its performance with those of the benchmark schemes. We first investigate the CPU utilization across the assigned locations of the four drones, depicted in Figs. 4(a), 5(a), 6(a) and 7(a). It can be observed that both EdgeDrones (ED) achieved the highest CPU utilization across the entire missions. Specifically, in Drone-1's mission, as shown in Fig. 4(a), EdgeDrones acheieves an average of 98% CPU utilization across Drones-1's assigned locations. This is followed by the NDFP, PDFP and PDNP schemes, which achieves the same average of 92.5%. The remaining two schemes, RP and RD achieves the lowest CPU utilization across the same assigned locations of Drone-1. RP achieves



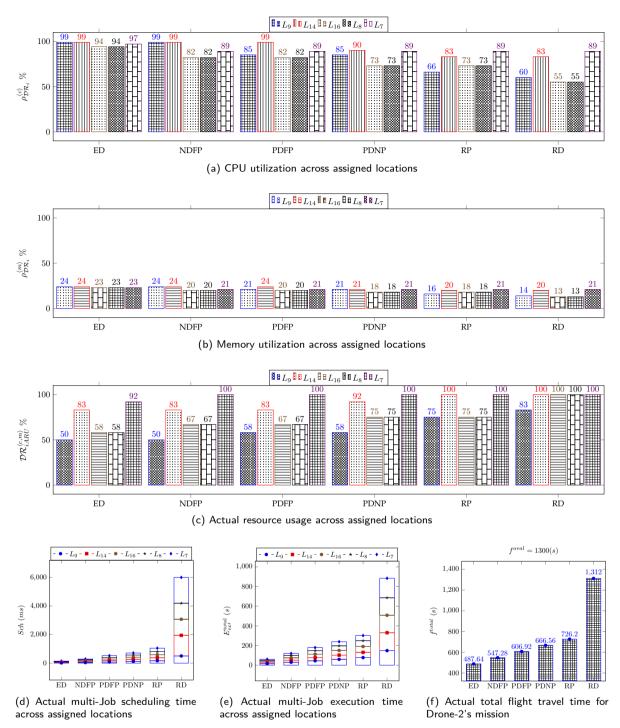


Figure 5: Activities and utilities of Drone-2's mission.

an average of 79.5%, while RD obviously achieves an average of 69.5%. It can also be seen that Edgedrones achieves the highest CPU utilization, according to Figs. 5(a), 6(a) and 7(a) for Drone-2, Drone-3 and Drone-4, respectively. Edgedrones is able to intelligently gang-schedule and co-locate all task tightly on nodes, resulting in higher resource utilization. In Figs. 5(a), 6(a) and 7(a), EdgeDrones achieves the highest average CPU utilization of 96.6, 92.6 and 94.5, respectively, compared to other schemes. In particular, NDFP achieves the second highest average CPU utilization across the assigned locations of the drones, i.e., it achieves an average CPU utilization of 6.4%, 4.8% and 3% less than EdgeDrones across Drone-2. Drone-3 and Drone-4 missions, respectively. PDFP and PDNP schemes performed averagely in terms of CPU utilization compared to EdgeDrones and NDFP, i.e., PDFP and PDNP achieve 9.2%, 14.6%; 14%, 18%; and 5%, 8.25%

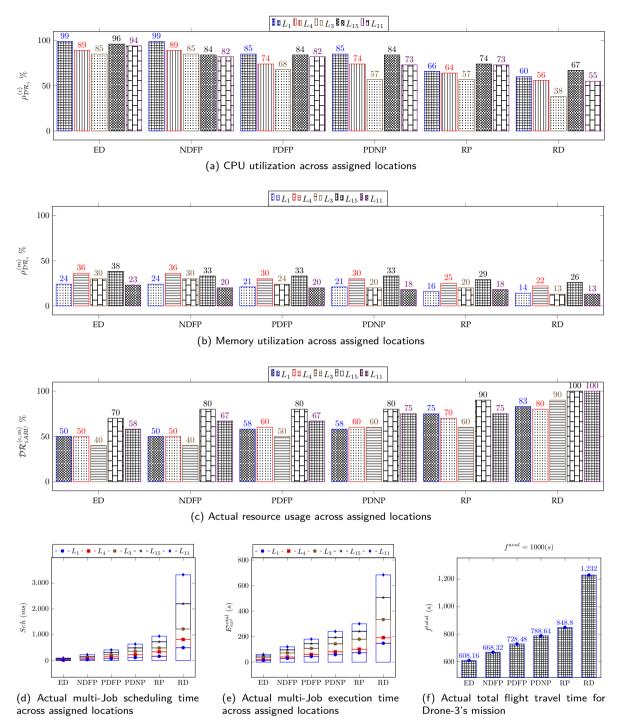


Figure 6: Activities and utilities of Drone-3's mission.

less than EdgeDrones across Drone-2, Drone-3 and Drone-4 missions, respectively. However, RP and RD schemes performed poorly mainly due to their resource under-utilization, i.e., both RP and RD schieve an average CPU utilization of 19.8%, 28.2%; 25.8%, 37.4%; and 18.25%, 44.25% less than EdgeDrones across Drone-2, Drone3 and Drone4 missions, respectively.

Figs. 4(b), 5(b), 6(b) and 7(b) compares the *Memory Utilization* of EdgeDrones with those of the four baseline

schemes and the random approach. Note that all the tasks executed across the assigned locations of the four drones are CPU intensive tasks, hence, the memory utilization across the locations are lower compared to the CPU utilizations. Nevertheless, EdgeDrones outperformes all the benchmark schemes, including the random approach. For example, across Drone-1's activities at its assigned locations, EdgeDrones is superior in achieving higher memory utilization with 3.5%, 6.5%, 6.5%, 12% and 18% more

memory utilizations than NDFP, PDFP, PDNP, RP and RD, respectively. It achieves 23.4% average memory utilization across Drone-2's activities, which surpasses NDFP, PDFP, PDNP, RP and RD by 1%, 2.2%, 3.6%, 4.8% and 7.2%, respectively. EdgDrones also outperforms all the baseline schemes and random approach across Drone-3 and Drone-4 assigned locations. It achieves an average of 30.2% and 17% memeory utilization at Drone-3 and Drone-4 activities, respectively. NDFP came second at both drones memory utilization (with 1.6% and 1.25% less compared to EdgeDrones at Drone-3 and Drone-4, respectively). PDFP and PDNP came third and forth, respectively. PDFP achieves 4.6% and 1.5% less in memory utilization compared to EgeDrones at Drone-3 and Drone-4 activities, respectively, while PDNP achieves 5.8% and 1.75% less compared to EgeDrones at Drone-3 and Drone-4 activities, respectively. RP and RD performance are the worst compared to all other schemes. In particular, the random approach (RD) achieves an average of 17.6% and 11.25% memory utilization across the assigned locations of Drone-3 and Drone-4, respectively, which is 12.6% and 5.75% less than EdgeDrones' achievement across the two drones activitites. RP on the other hand, achieves an average of 8.6% less compared to EdgeDrones across Drone-3 assigned locations, and an average of 3.5% less than EdgeDrones across Drone-4 assigned locations.

Figs. 4(c), 5(c), 6(c) and 7(c) compares the actual resource usage $\mathcal{DR}_{iARU}^{(c,m)}$ of EdgeDrones with those of the four baseline schemes and the random approach. It can be seen that solving the optimization (45) is directly linked to minimize the ARU (49), by packing or co-locating tasks firmly on available resources. Hence, it can be seen that EdgeDrones consumes the fewest resources across the integrated drones activities with NPFP as the very close second best, while Random uses all the resources across almost all the drones activities with RP as the second worst. PDFP ranks in the middle, in terms of resource usage across the drones activities. Again, EdgeDrones and NDFP are superior than PDFP, PDNP, RD and Random, and they achieve the highest and close second highest resource utilization across the integrated drones activities, respectively. For example, across Drone-1's assigned locations, EdgeDrone uses the fewest resources amounting to an average of 71% compared to NDFP, PDFP, PDNP and RP, which use 4%, 8%, 8%, 16.5% more than EdgeDrones, respectively. However, RD uses all available resources, i.e., 100%, due to its inability to co-locate tasks on nodes. Across Drone-2, Drone-3 and Drone-4's assgined locations, it can also be seen that EdgeDrones uses fewer resources, up to an average of 37% less compared to other baseline schemes and the random approach.

Three other key metrics are the actual multi-tasks/job scheduling time $\sum_{J \in \mathbb{J}} \sum_{z=1}^{m} \sum_{i=1}^{k_z} S_{ch_{z_i}}/k_z$, where *m* is the number of scheduling units, k_z is the number of tasks within the *z*-th scheduling unit, and more importantly, the actual multi-tasks/jobs execution time E_{ext}^{total} and total flight travel time f^{total} . Figs. (4(d), 5(d), 6(d), (7(d)); (4(e), 5(e), 6(e), 7(e)); and (4(f), 5(f), 6(f), 7(f)) compares the actual

multi-tasks/jobs scheduling time, multi-tasks/job execution time and total flight travel time of EdgeDrones with those of the four benchmarks and random approach, respectively. The results show that EdgeDrones is the best, NDFP is the second best, and PDFP is the third best, PDNP is forth best, while RP and Random is the worst and RP the second worst, in terms of actual task scheduling times, actual task execution times and drone's total flight time. The superior performance of EdgeDrones over the other benchmarks is overwhelmingly clear.

5.2.3. Performance comparison in individual Drones

Figs. 4 \sim 7 show the performance of the schemes in terms of resource utilization, actual resource usage, actual task scheduling times, actual task execution times and actual flight total travel times across the integrated drones. We now delve into the individual drone to examine the performance of all the schemes.

Drone-1 is attached with three Huawei AR502H Series edge devices, with total resource capacity of 12 Cores and 6 GiB for CPU and memory, respectively. The entire weight of the devices is ≈ 3.3 kg. Its assigned locations $DR_1^{assign} =$ $\{L_{12}, L_{13}\}$; its route $\mathcal{DR}_1^{route} = L_0 \rightarrow L_{12} \rightarrow L_{13} \rightarrow L_0$; and its flight time availability $f^{aval} = 700(s)$. We deploy 9 jobs with a total of 36 tasks, where the job has a task dependency depth γ (1, 5]. Utilizing the gang scheduling strategy, EdgeDrones co-shedules and co-locates all the 9 jobs at a time in the attached edge devices as possible to minimize the overall used nodes. These jobs are tightly co-located, which enables dependent tasks to communicate and share data effectively. As a result, EdgeDrones achieves the fastest scheduling time and execution time compared to NDFP, PDFP, PDNF, RP and the random approach. In addition, EdgeDrones only uses an average 71% of resources to execute the jobs. Using the same resource capacity, NPFP, PDFP, PDNP and RP utilize an average of 75%, 79%, 79% and 87.5% of the resources, respectively, as shown in Fig 4(c). The random approach uses all available resources. It is observed that EdgeDrones is 1.7 times and 2 times faster than the second best NDFP in both the scheduling time and execution time, respectively. EdgeDrones is more than 3 times and more than 3 times faster than PDFP as well as more than 5 times and more than 4 times faster than PDNP in the scheduling time and execution time, as shown in Figs 4(d) and 4(e), respectively. EdgeDrones is 6 times and 5 times faster than the RP, as well as 47 times and 18 times faster than the random approach in the scheduling and execution times, respectively. The most important is the total flight travel time f^{total} of Drone-1, such that if along the drone's mission f^{total} becomes greater than its flight time availability f^{aval} , then it might lead to loss of job or mission failure. Recall that $f^{total} = \sum_{i=0}^{n} \sum_{j=i+1}^{m} f^{L_{i,j}} + f^{hover}_{j}$, where $\sum_{j=0}^{m} f_{j}^{hover} = E_{ext}^{\text{total}}.$ Hence, EdgeDrones is bale to quickly schedule and execute all the 9 jobs at each assigned location, resulting to a successful mission, and the fastest f^{total} (upto 2.7 times faster) compared to the four baseline schems and the random approach, as shown in Fig. 4(f). The random

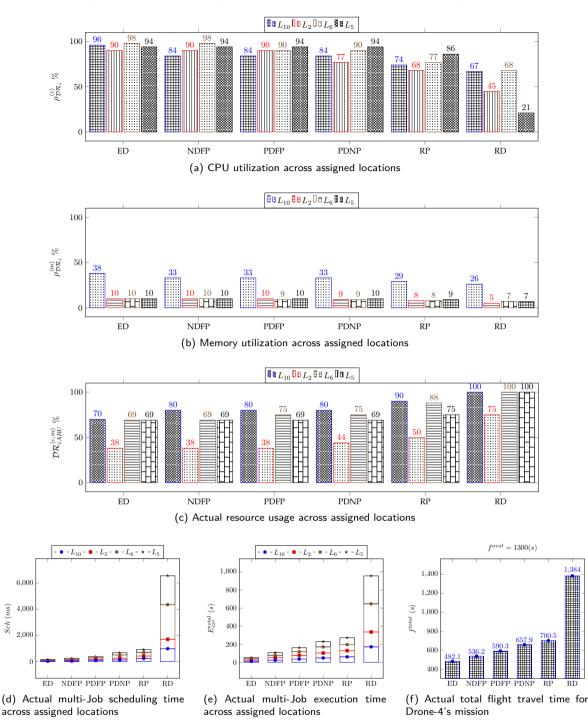


Figure 7: Activities and utilities of Drone-4's mission.

approach could not successfully schedule and execute all the jobs within the drone's f^{aval} , thereby leading to a failed mission.

Like Drone-1, Drone-2 is attached with two HIVECELL portable edge devices with total weight of ≈ 2.72 kg, and total resource capacity of 12 Cores and 16 GiB for CPU and memory, respectively. Its assigned locations $\mathcal{DR}_1^{assign} = \{L_7, L_8, L_9, L_{14}, L_{16}\}$; it has to visit and execute tasks on route: $\mathcal{DR}_1^{route} = L_0 \rightarrow L_9 \rightarrow L_{14} \rightarrow L_{16} \rightarrow L_8 \rightarrow L_7 \rightarrow$

 L_0 ; and its flight time availability $f^{aval} = 1200(s)$. Here, we deploy a total of $\mathbb{J} = 17$, where each $J \in \mathbb{J}$ has a task dependency in the range of (1, 4]. The total number of tasks in $\Sigma \mathbb{J}$ is 76. We ensure that the attached edge resources are fully utilized by co-locating the jobs tightly on them. As discussed earlier, application container provides isolation to co-located tasks, thereby eliminating interference and resource contentions in the cluster. A single node is capable of running several containerized tasks, given that available

resources are sufficient. In this drone's activities, Edgedrones consume an average of 5.2% fewer resources than NDFP, an average of 6.8%, 11.8%, 16.8% and 28.4% fewer resources than PDFP, PDNP, RP and Random, EdgeDrones, also gain an average of upto 28.2% higher CPU utilization over NDFP, PDFP, PDNP, RP and Random (as shown in Fig. 5(a)), as well as an average of upto 7.2% higher memory utilization than NDFP, PDFP, PDNP, RP and Random, as shown in Fig. 5(b). More significantly, EdgeDrones is 2, 3.7, 5, 7.4 and 42.5 times faster in the scheduling time than NDFP, PDFP, PDNP, RP and the random approach respectively, while it is 2, 3, 4, 5 and 14.8 times faster in the execution time than NDFP, PDFP, PDNP, RP and the random approach, respectively across the assigned locations. Although all the schemes, except for the random approach were able to complete their task within the drone's flight time availability, nonetheless, EdgeDrones achieves the fasters mission completion time, which is much more less than the drone's flight time availability. It can be seen in Fig. 5(f) that EdgeDrones is upto 2.7 times faster than other schemes.

Drone-3 is attached with one HIVECELL and one Huawei AR502H Series portable edge devices with total weight of \approx 2.46kg. Its total resource capacity is 10 Cores and 8 GiB of CPU and memory, respectively. Its assigned locations $D\mathcal{R}_1^{assign} = \{L_1, L_3, L_4, L_{11}, L_{15}\};$ flight travel route $\mathcal{DR}_{1}^{loute} = \{L_{0} \rightarrow L_{1} \rightarrow L_{4} \rightarrow L_{3} \rightarrow L_{15} \rightarrow L_{11} \rightarrow L_{0}\};$ and flight time availability $f^{aval} = 900(s)$. In this cluster, we deploy $\mathbb{J} = 12$ in total of 56 tasks, where each $J \in \mathbb{J}$ has a task dependency depth γ range (2, 5]. Across this drone's activities, EdgeDrones achieve reduced $D\mathcal{R}_{iARU}^{\langle c,m\rangle}$ by 3.8%, 9.4%, 13%, 20.4% and 37% compared with NDFP, PDFP, PDNP, RD and Random, respectively (as shown in Fig. 6(c)). EdgeDrones achieve 4.8%, 14%, 18%, 25.8% and 37.4% higher CPU utilization as well as 1.6%, 4.6%, 5.8%, 8.6% and 12.6% higher memory utilization compared to NDFP, PDFP, PDNP, RD and Random, respectively. In terms of scheduling. EdgeDrones is about 2.3 times, 4.2 times, 6.5 times, 9.6 times and 34 times faster than NDFP, PDFP, PDNP, RD and Random, respectively (as shown in Fig. 6(d)). It achieves approximately 2 times, 3 times, 4 times, 5 times and 11 times faster execution times than NDFP, PDFP, PDNP, RD and Random, respectively (as shown in Fig. 6(e)). Not surprisingly, Random has the worst scheduling time and execution time performance, resulting to incomplete mission (since its $f^{total} > f^{aval}$). On the other hand, EdgeDrones achieves the fasters mission completion time, which is upto 1.4 times faster than NDFP, PDFP, PDNP and RD, as shown in Fig. 6(f).

Drone-4 is attached with Azure Stack Edge mini memory intensive edge device, with resource capacity of 16 Cores and high memory capacity of 48 GiB. It is four locations $D\mathcal{R}_1^{assign} = \{L_2, L_5, L_6, L_{10}\}$, where there are total of 16 jobs made up of 70 tasks to be executed. Its flight travel route $D\mathcal{R}_1^{route} = \{L_0 \rightarrow L_{10} \rightarrow L_2 \rightarrow L_6 \rightarrow L_5\}$, and flight time availability $f^{aval} = 1200(s)$. It is observed that EdgeDrones consumes the fewest resources at an average of 61.5%, followed by NDFP at 64%. PDFP consumes an average of 65.5%, PDNP consumes an average of 67%, RP consumes an average 75.75% of the resources, while the Random approach uses all the available resources at L_{10} , L_6 and L_5 locations, but consumes an average of 93.75% of resources across the assigned locations. EdgeDrones also achieves 3%, 5%, 8.25%, 18.25% and 44.25% higher CPU utilization over NDFP, PDFP, PDNP, RP and Random, respectively (as shown in Fig. 7(a)). Note the edge device attached to this drone is memory intensive, i.e., it has huge memory capacities compared to the memory resource request of jobs at the assigned locations. Therefore, the jobs can only consume few such capacities, as shown in Fig. 7(b). In terms of scheduling time, EdgeDrones is approximately 1.6 times, 1.5 times, 4.3 times, 5.8 times and 41.7 times faster than NDFP, PDFP, PDNP, RP and random respectively (Fig. 7(d)). In terms of execution time, EdgeDrones is about 2 times, 3 times, 4.2 times, 5 times and 17.6 times faster than NDFP, PDFP, PDNP, RP and the random approach respectively (Fig. 7(e)). Importantly, Drone-4 completed its mission with the fastest time under EdgeDrones strategy, with up to 1.5 times faster compared to the baseline schemes, as shown in Fig. 7(f)).

5.3. Discussion

Overall, EdgeDrones has demonstrated better performance in an integrated edge computing system. It has consistently outperform existing schemes (NDFP, PDFP, PDNP, RP and Random) by achieving faster scheduling times and excution times, while using fewer resources. Most importantly, effective multi-tasks scheduling and execution of EdgeDrones across the locations, enables faster tasks response times and mission completion times. EdgeDrones achievements is attributed to its effective orchestration strategy, gang-deployment and co-location of multi-jobs, which allows inter-dependent tasks within each job to communicate and share data faster. Such fast execution is crucial for modern applications to perform better. The existing schemes do not consider task's dependecies or multi-tasks co-location, leading to limited edge resource wastage through under utilization, as well as causing execution delay.

6. Conclusions

This paper has presented a novel Multi-Location Capacitated Mission Scheduling Problem (MLCMSP) that selects suitable drones and co-schedules their flight routes with the least total distance to visit and execute tasks at the target locations. We proposed an intelligent multi-dependent tasks orchestration scheme called EdgeDrones, a variant binpacking optimization approach through gang-scheduling of multi-dependent tasks, that co-schedules and co-locate tasks firmly on available nodes, so as to avoid resource wastage. Evaluations using real world workloads from Alibaba clusters, shows that EdgeDrones is superiror compared to the baseline schemes. Importantly, EdgeDrones is able to avoid loss of jobs in aerial edge computing missions. In our future research, we can further integrate cost models to MLCMSP by assigning an operational cost per drone's mission, and convert the MLCMSP into a profit maximization problem. In addition, we can also deploy on-premise (fog-based computing) alongside with the drones (aerial-based computing) to form a hybrid deployment.

CRediT authorship contribution statement

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Declaration of competing interest

The authors declare that they have no known competing financial interests or personal relationships that could have appeared to influence the work reported in this paper.

Data availability

Data will be made available on request.

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